

# Introduction to constructive algebraic analysis

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# Mathematical systems theory

- Systems theory aims at studying **general systems of equations**
  - ① ordinary/partial differential equations, difference equations, time-delay equations... = **functional equations**,
  - ② determined, overdetermined, underdetermined (open, close),
  - ③ linear, time-varying, nonlinear systems...

coming from engineering sciences (e.g., electromagnetism, elasticity, hydrodynamics), mathematical physics, control theory...

- The purpose of these talks is to give a short introduction to **algebraic systems theory** which studies linear functional systems by means of **constructive algebra**, **module theory**, **homological algebra**.
- Algebraic systems theory can be traced back to the development of **algebraic analysis** by Malgrange, Sato, Kashiwara...

# Outline

The purpose of this talk is to develop the following 3 ideas:

- ① A large class of linear functional systems can be studied by means of a **non-commutative polynomial approach** over skew polynomial rings and Ore algebras of functional operators.  
Non-commutative Gröbner bases  $\Rightarrow$  **constructive approach**.
- ② **Algebraic analysis** is a natural mathematical framework for the intrinsic study of linear systems theory (**module theory**).
- ③ **Constructive homological algebra** allows us to develop algorithms and symbolic packages dedicated to the study of the **structural properties** of linear functional systems.

# Matrices of differential operators

- **Newton:** Fluxion calculus (1666) (“dot-age”)

$$\begin{cases} \ddot{x}_1(t) + \alpha x_1(t) - \alpha u(t) = 0, \\ \ddot{x}_2(t) + \alpha x_2(t) - \alpha u(t) = 0, \end{cases} \quad \alpha = g/l.$$

- **Leibniz:** Infinitesimal calculus (1676) (“d-ism”)

$$\begin{cases} \frac{d^2 x_1(t)}{dt^2} + \alpha x_1(t) - \alpha u(t) = 0, \\ \frac{d^2 x_2(t)}{dt^2} + \alpha x_2(t) - \alpha u(t) = 0. \end{cases}$$

- **Boole:** Operational calculus (1859-60)

$$\begin{pmatrix} \frac{d^2}{dt^2} + \alpha & 0 & -\alpha \\ 0 & \frac{d^2}{dt^2} + \alpha & -\alpha \end{pmatrix} \begin{pmatrix} x_1(t) \\ x_2(t) \\ u(t) \end{pmatrix} = 0.$$

⇒ Ring of differential operators  $D = \mathbb{Q}(\alpha) [\frac{d}{dt}]$ :

$$\sum_{i=0}^n a_i \left( \frac{d}{dt} \right)^i \in D, \quad a_i \in \mathbb{Q}(\alpha), \quad \left( \frac{d}{dt} \right)^i = \frac{d}{dt} \circ \dots \circ \frac{d}{dt} = \frac{d^i}{dt^i}.$$

# Functional operators

- Differential operators:  $\left(\sum_{j=0}^m b_j(t) \frac{d^j}{dt^j}\right) \left(\sum_{i=0}^n a_i(t) \frac{d^i}{dt^i}\right)$

$$\frac{d}{dt}(a y) = a \frac{d}{dt} y + \left(\frac{da}{dt}\right) y \quad \Rightarrow \quad \frac{d}{dt} a \cdot = a \frac{d}{dt} \cdot + \frac{da}{dt}.$$

- Shift operators:  $\delta a(t) = a(t - h)$ ,  $\sigma a_n = a_{n+1}$ .

$$\delta(a(t)y(t)) = a(t - h)y(t - h) = \delta a \delta y \quad \Rightarrow \quad \delta a \cdot = (\delta a) \delta \cdot.$$

$$\sigma(a_n y_n) = a_{n+1} y_{n+1} = \sigma a \sigma y \quad \Rightarrow \quad \sigma a \cdot = (\sigma a) \sigma \cdot.$$

- Difference operators:  $\Delta a(x) = a(x + 1) - a(x)$ .

$$\Delta a(x) \cdot = a(x + 1) \Delta \cdot + (\Delta a) \cdot$$

- Divided difference operators:  $d_{x_0} a(x) = \frac{a(x) - a(x_0)}{x - x_0}$ .

$$d_{x_0} a(x) \cdot = a(x_0) d_{x_0} \cdot + (d_{x_0} a) \cdot$$

- $q$ -difference,  $q$ -shift,  $q$ -dilation, Frobenius, Euler operators

# Skew polynomial rings (Ore, 1933)

- Definition: A **skew polynomial ring**  $A[\partial; \alpha, \beta]$  is a non-commutative polynomial ring in  $\partial$  with coefficients in  $A$  satisfying

$$\forall a \in A, \quad \boxed{\partial a = \alpha(a) \partial + \beta(a)}$$

where  $\alpha : A \longrightarrow A$  and  $\beta : A \longrightarrow A$  are such that:

$$\left\{ \begin{array}{l} \alpha(1) = 1, \\ \alpha(a+b) = \alpha(a) + \alpha(b), \\ \alpha(ab) = \alpha(a)\alpha(b), \end{array} \right. \quad \left\{ \begin{array}{l} \beta(a+b) = \beta(a) + \beta(b), \\ \beta(ab) = \alpha(a)\beta(b) + \beta(a)b. \end{array} \right.$$

- $P \in A[\partial; \alpha, \beta]$  has a unique expression  $P = \sum_{i=0}^n a_i \partial^i$ ,  $a_i \in A$ .
  - Ring of differential operators:  $A[\partial; \text{id}, \frac{d}{dt}]$ .
  - Ring of shift operators:  $A[\partial; \delta, 0]$ ,  $A[\partial; \sigma, 0]$ .
  - Ring of difference operators:  $A[\partial; \tau, \tau - \text{id}]$ ,  $\tau a(x) = a(x+1)$ .

# Ore algebras (Chyzak-Salvy, 1996)

- We can iterate skew polynomial rings to obtain Ore extensions:

$$A[\partial_1; \alpha_1, \beta_1] \dots [\partial_n; \alpha_n, \beta_n]$$

- Definition: An Ore extension  $A[\partial_1; \alpha_1, \beta_1] \dots [\partial_n; \alpha_n, \beta_n]$  is called an Ore algebras if the  $\partial_i$ 's commute  $\partial_i \partial_j = \partial_j \partial_i$ , i.e., if we have

$$1 \leq j < i \leq m, \quad \alpha_i(\partial_j) = \partial_j, \quad \beta_i(\partial_j) = 0,$$

and the  $\alpha_i|_A$ 's and  $\beta_j|_A$ 's commute for  $i \neq j$ .

- Ring of differential operators:  $A\left[\partial_1; \text{id}, \frac{\partial}{\partial x_1}\right] \dots \left[\partial_n; \text{id}, \frac{\partial}{\partial x_n}\right]$ .
- Ring of differential delay operators:  $A\left[\partial_1; \text{id}, \frac{d}{dt}\right] [\partial_2; \delta, 0]$ .
- Ring of shift operators:  $A[\partial_1; \sigma_1, 0] \dots [\partial_n; \sigma_n, 0]$ .

# Matrix of functional operators

- The wind tunnel model (Manitius, IEEE TAC 84):

$$\begin{cases} \dot{x}_1(t) + a x_1(t) - k a x_2(t-h) = 0, \\ \dot{x}_2(t) - x_3(t) = 0, \\ \dot{x}_3(t) + \omega^2 x_2(t) + 2\zeta\omega x_3(t) - \omega^2 u(t) = 0. \end{cases} \quad (*)$$

- We first introduce the **commutative Ore algebra** of ordinary differential time-delay operators:

$$D = \mathbb{Q}(a, k, \omega, \zeta) \left[ \partial_1; \text{id}, \frac{d}{dt} \right] [\partial_2; \delta, 0].$$

- Then, the system (\*) can be rewritten as:

$$\begin{pmatrix} \partial_1 + a & -k a \partial_2 & 0 & 0 \\ 0 & \partial_1 & -1 & 0 \\ 0 & \omega^2 & \partial_1 + 2\zeta\omega & -\omega^2 \end{pmatrix} \begin{pmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \\ u(t) \end{pmatrix} = 0.$$

# Matrix of functional operators

- Linearization of the Navier-Stokes  $\sim$  a parabolic Poiseuille profile

$$\begin{cases} \partial_t u_1 + 4y(1-y)\partial_x u_1 - 4(2y-1)u_2 - \frac{1}{Re}(\partial_x^2 + \partial_y^2)u_1 + \partial_x p = 0, \\ \partial_t u_2 + 4y(1-y)\partial_x u_2 - \frac{1}{Re}(\partial_x^2 + \partial_y^2)u_2 + \partial_y p = 0, \\ \partial_x u_1 + \partial_y u_2 = 0. \end{cases} \quad (*) \quad (\text{e.g., Vazquez-Krstic, IEEE TAC 07})$$

- Let us introduce the so-called Weyl algebra ( $\partial_x x = x\partial_x + 1$ ):

$$D = \mathbb{Q}(Re)[t, x, y] \left[ \partial_t; \text{id}, \frac{\partial}{\partial t} \right] \left[ \partial_x; \text{id}, \frac{\partial}{\partial x} \right] \left[ \partial_y; \text{id}, \frac{\partial}{\partial y} \right].$$

- The system  $(*)$  is defined by the matrix of PD operators:

$$\begin{pmatrix} \partial_t + 4y(1-y)\partial_x - \frac{1}{Re}(\partial_x^2 + \partial_y^2) & -4(2y-1) & \partial_x \\ 0 & \partial_t + 4y(1-y)\partial_x - \frac{1}{Re}(\partial_x^2 + \partial_y^2) & \partial_y \\ \partial_x & \partial_y & 0 \end{pmatrix}.$$

# Non-commutative Gröbner bases

- Let  $D = A[\partial_1; \alpha_1, \beta_1] \dots [\partial_m; \alpha_m, \beta_m]$  be an Ore algebra.
- **Theorem:** (Kredel, 93) Let  $A = k[x_1, \dots, x_n]$  a commutative polynomial ring ( $k = \mathbb{Q}, \mathbb{F}_p$ ) and  $D$  an Ore algebra satisfying

$$\alpha_i(x_j) = a_{ij}x_j + b_{ij}, \quad \beta_i(x_j) = c_{ij},$$

for certain  $0 \neq a_{ij} \in k$ ,  $b_{ij} \in k$ ,  $c_{ij} \in A$  and  $\deg(c_{ij}) \leq 1$ . Then, a non-commutative version of **Buchberger's algorithm** terminates for any term order and its result is a **Gröbner basis**.

- **Implementation** in the Maple package **Ore\_algebra** (Chyzak).
- Gröbner bases can be used to effectively compute over  $D$ .
- $D$  is a **left noetherian domain**  $\Rightarrow$  **left Ore domain**:

$$\forall a, b \in D \setminus \{0\}, \exists c, d \in D \setminus \{0\} : c a = d b.$$

# Algebraic analysis

- Let  $D$  be an Ore algebra and  $R \in D^{q \times p}$ .
- Let us consider the **left  $D$ -homomorphism**:

$$\begin{array}{ccc} D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} \\ \lambda = (\lambda_1 \dots \lambda_q) & \longmapsto & \lambda R. \end{array}$$

- We introduce the finitely presented left  $D$ -module:

$$M = D^{1 \times p} / \text{im}_D(\cdot R) = D^{1 \times p} / (D^{1 \times q} R).$$

- $M$  is formed by the equivalence classes  $\pi(\mu)$  of  $\mu \in D^{1 \times p}$  for the equivalence relation  $\sim$  on  $D^{1 \times p}$  defined by:

$$\mu_1 \sim \mu_2 \Leftrightarrow \mu_1 - \mu_2 \in D^{1 \times q} R.$$

- ① Number theory:  $\mathbb{C} = \mathbb{R}[x]/(x^2 + 1)$ ,  $\mathbb{Z}[i\sqrt{5}] = \mathbb{Z}[x]/(x^2 + 5)$ .
- ② Algebraic geometry:  $\mathbb{C}[x, y]/(x^2 + y^2 - 1, x - y)$ .

# Generators and relations of $M = D^{1 \times p} / (D^{1 \times q} R)$

- Let  $\{f_k\}_{k=1,\dots,p}$  the standard basis of  $D^{1 \times p}$  defined by:

$$f_1 = (1 \ 0 \ \dots \ 0), \quad f_k = (0 \ \dots \ 1 \ \dots \ 0), \quad f_p = (0 \ \dots \ 0 \ 1).$$

- Let  $\pi : D^{1 \times p} \longrightarrow M$  be the  $D$ -homomorphism sending  $\mu$  to  $\pi(\mu)$ ,

$$y_k = \pi(f_k), \quad k = 1, \dots, p.$$

- $\{y_k\}_{k=1,\dots,p}$  is a family of generators of  $M$ :

$$\forall m \in M, \exists \mu \in D^{1 \times p} : m = \pi(\mu) = \sum_{k=1}^p \mu_k \pi(f_k) = \sum_{k=1}^p \mu_k y_k.$$

- $\{y_k\}_{k=1,\dots,p}$  satisfies the relations:

$$\pi((R_{I1} \dots R_{Ip})) = \pi \left( \sum_{k=1}^p R_{Ik} f_k \right) = \sum_{k=1}^p R_{Ik} y_k = 0, \quad I = 1, \dots, q,$$

$\Rightarrow y = (y_1 \dots y_p)^T$  satisfies the relation  $R y = 0$ .

# Duality modules – systems

- Let  $\mathcal{F}$  be a left  $D$ -module and  $\text{hom}_D(M, \mathcal{F})$  the abelian group:

$$\text{hom}_D(M, \mathcal{F}) = \{f : M \rightarrow \mathcal{F} \mid f(d_1 m_1 + d_2 m_2) = d_1 f(m_1) + d_2 f(m_2)\}.$$

- Applying the contravariant left exact functor  $\text{hom}_D(\cdot, \mathcal{F})$  to

$$D^{1 \times q} \xrightarrow{\cdot R} D^{1 \times p} \xrightarrow{\pi} M \longrightarrow 0,$$

we obtain the following exact sequence of abelian groups:

$$\mathcal{F}^q \xleftarrow{R \cdot} \mathcal{F}^p \xleftarrow{\iota \circ \pi^*} \text{hom}_D(M, \mathcal{F}) \longrightarrow 0.$$

- Theorem (Malgrange):

$$\boxed{\text{hom}_D(M, \mathcal{F}) \cong \ker_{\mathcal{F}}(R \cdot) = \{\eta \in \mathcal{F}^p \mid R \eta = 0\}}$$

- Remark:  $\text{hom}_D(M, \mathcal{F})$  intrinsically characterizes  $\ker_{\mathcal{F}}(R \cdot)$  as it does not depend on the embedding of  $\ker_{\mathcal{F}}(R \cdot)$  into  $\mathcal{F}^p$ .

# Module theory

- Definition: 1.  $M$  is **free** if  $\exists r \in \mathbb{Z}_+$  such that  $M \cong D^r$ .
- 2.  $M$  is **stably free** if  $\exists r, s \in \mathbb{Z}_+$  such that  $M \oplus D^s \cong D^r$ .
- 3.  $M$  is **projective** if  $\exists r \in \mathbb{Z}_+$  and a  $D$ -module  $P$  such that:

$$M \oplus P \cong D^r.$$

- 4.  $M$  is **reflexive** if  $\varepsilon : M \longrightarrow \text{hom}_D(\text{hom}_D(M, D), D)$  is an isomorphism, where:

$$\varepsilon(m)(f) = f(m), \quad \forall m \in M, \quad \forall f \in \text{hom}_D(M, D).$$

- 5.  $M$  is **torsion-free** if:

$$t(M) = \{m \in M \mid \exists 0 \neq P \in D : P m = 0\} = 0.$$

- 6.  $M$  is **torsion** if  $t(M) = M$ .

$$0 \longrightarrow t(M) \longrightarrow M \longrightarrow M/t(M) \longrightarrow 0 \quad \text{exact sequence.}$$

# Classification of modules

- Theorem: 1. We have the following implications:

free  $\Rightarrow$  stably free  $\Rightarrow$  projective  $\Rightarrow$  reflexive  $\Rightarrow$  torsion-free.

- 2. If  $D$  is a principal domain (e.g.,  $B_1(\mathbb{Q}) = \mathbb{Q}(t) [\partial; \text{id}, \frac{d}{dt}]$ ), then:

torsion-free = free.

- 3. If  $D$  is a hereditary ring (e.g.,  $A_1(\mathbb{Q}) = \mathbb{Q}[t] [\partial; \text{id}, \frac{d}{dt}]$ ), then:

torsion-free = projective.

- 4. If  $D = k[x_1, \dots, x_n]$  and  $k$  a field, then:

projective = free (Quillen-Suslin theorem).

- 5. If  $D = A_n(k)$  or  $B_n(k)$ ,  $k$  is a field of characteristic 0, then

projective = free (Stafford theorem),

for modules of rank at least 2.

# Free resolutions

- Definition: A sequence of  $D$ -morphisms  $M' \xrightarrow{f} M \xrightarrow{g} M''$  is called a **complex** if  $g \circ f = 0$ , i.e.,  $\text{im } f \subseteq \ker g$ .

The **defect of exactness** at  $M$  is  $H(M) = \ker g / \text{im } f$ .

The complex is **exact** at  $M$  if  $\text{im } f = \ker g$ .

- Definition: A **free resolution** of a left  $D$ -module  $M$  is an exact sequence of the form

$$\dots \xrightarrow{.R_3} D^{1 \times I_2} \xrightarrow{.R_2} D^{1 \times I_1} \xrightarrow{.R_1} D^{1 \times I_0} \xrightarrow{\pi} M \longrightarrow 0,$$

where  $R_i \in D^{I_i \times I_{i-1}}$  and:

$$\begin{array}{ccc} D^{1 \times I_i} & \xrightarrow{.R_i} & D^{1 \times I_{i-1}} \\ (P_1 \dots P_{I_i}) & \longmapsto & (P_1 \dots P_{I_i}) R_i. \end{array}$$

- Algorithm: Find a basis of the compatibility conditions of the inhomogeneous system  $R_i y = u$  by **eliminating  $y$**  (e.g., Gb):

$$\forall P \in \ker_D(.R_i), \quad P(R_i y) = P u \Rightarrow P u = 0.$$

# Extension functor $\text{ext}_D^i(\cdot, \mathcal{F})$

- We define the reduced free resolution of  $M$  by:

$$\dots \xrightarrow{R_3} D^{1 \times I_2} \xrightarrow{R_2} D^{1 \times I_1} \xrightarrow{R_1} D^{1 \times I_0} \longrightarrow 0 \quad (\star).$$

- Let  $\mathcal{F}$  be a left  $D$ -module. Applying the functor  $\text{hom}_D(\cdot, \mathcal{F})$  to  $(\star)$ , we obtain the following **complex**:

$$\dots \xleftarrow{R_3 \cdot} \mathcal{F}^{I_2} \xleftarrow{R_2 \cdot} \mathcal{F}^{I_1} \xleftarrow{R_1 \cdot} \mathcal{F}^{I_0} \longleftarrow 0 \quad (\star\star)$$

$$\begin{aligned} \text{where } \mathcal{F}^{I_i} &\xleftarrow{R_i \cdot} \mathcal{F}^{I_{i-1}} \\ R_i \eta &\longleftarrow \eta. \end{aligned}$$

- We denote the **defects of exactness** of  $(\star\star)$  by:

$$\begin{cases} \text{ext}_D^0(M, \mathcal{F}) = \text{hom}_D(M, \mathcal{F}) \cong \ker_{\mathcal{F}}(R_1 \cdot), \\ \text{ext}_D^i(M, \mathcal{F}) \cong \ker_{\mathcal{F}}(R_{i+1} \cdot) / \text{im}_{\mathcal{F}}(R_i \cdot), \quad i \geq 1. \end{cases}$$

- **Theorem:** The abelian group  $\text{ext}_D^i(M, \mathcal{F})$  only depends on  $M$  and  $\mathcal{F}$  and not on the resolution  $(\star)$ .

# Auslander transpose

- **Definition:** Let  $R \in D^{q \times p}$ . If  $M = D^{1 \times p}/(D^{1 \times q} R)$  denotes the left  $D$ -module finitely presented by  $R$ , then its **Auslander transpose** is the right  $D$ -module defined by  $N = D^q/(R D^p)$ .
- **Proposition:** The Auslander transpose  $N = D^q/(R D^p)$  only depends on  $M$  up to a projective equivalence.

Hence, if we have  $M = D^{1 \times p'}/(D^{1 \times q'} R')$  and  $N' = D^{q'}/(R' D^{p'})$  denotes the corresponding Auslander transpose, then there exist two projective right  $D$ -modules  $P$  and  $P'$  such that:

$$N \oplus P \cong N' \oplus P'.$$

- **Proposition:** If  $P$  is a **projective module**, then we have:

$$\text{ext}_D^i(P, \mathcal{F}) = 0, \quad \forall i \geq 1.$$

- **Corollary:** The  $\text{ext}_D^i(N, \mathcal{F})$ 's,  $i \geq 1$ , only depend on  $M$  and  $\mathcal{F}$ .

Module $M$	Homological algebra	$\mathcal{F}$ injective cogenerator
with torsion	$t(M) \cong \text{ext}_D^1(N, D)$	$\emptyset$
torsion-free	$\text{ext}_D^1(N, D) = 0$	$\ker_{\mathcal{F}}(R.) = Q \mathcal{F}^l$
reflexive	$\text{ext}_D^i(N, D) = 0$ $i = 1, 2$	$\ker_{\mathcal{F}}(R.) = Q_1 \mathcal{F}^l$ $\ker_{\mathcal{F}}(Q_1.) = Q_2 \mathcal{F}^l$
projective = stably free	$\text{ext}_D^i(N, D) = 0$ $i = 1, \dots, n = \text{gld}(D)$	$\ker_{\mathcal{F}}(R.) = Q_1 \mathcal{F}^l$ $\ker_{\mathcal{F}}(Q_1.) = Q_2 \mathcal{F}^l$ $\dots$ $\ker_{\mathcal{F}}(Q_{n-1}.) = Q_n \mathcal{F}^l$
free	$\exists Q \in D^{p \times m}, T \in D^{m \times p},$ $\ker_D(.Q) = D^{1 \times q} R, T Q = I_m$	$\ker_{\mathcal{F}}(R.) = Q \mathcal{F}^m,$ $\exists T \in D^{m \times p} : T Q = I_m$

# Parametrizability problem

$$\begin{array}{ccccccc} D^{1 \times q} & \xrightarrow{.R} & D^{1 \times p} & \xrightarrow{\pi^*} & M & \longrightarrow 0 \\ 0 \longleftarrow N \xleftarrow{\kappa} & D^q & \xleftarrow{R.} & D^p & \xleftarrow{\pi^*} & \hom_D(M, D) & \longleftarrow 0 \\ 0 \longleftarrow N \xleftarrow{\kappa} & D^q & \xleftarrow{R.} & D^p & \xleftarrow{Q.} & D^m \\ D^{1 \times q} & \xrightarrow{.R} & \color{red}{D^{1 \times p}} & \xrightarrow{.Q} & D^{1 \times m} & & (\star) \end{array}$$

$t(M) \cong \text{ext}_D^1(N, D) \cong \ker_D(.Q)/(D^{1 \times q} R) = (D^{1 \times q'} R')/(D^{1 \times q} R)$ ,  
where  $R'$  is defined by  $\ker_D(.Q) = D^{1 \times q'} R'$ , and:

$$M/t(M) \cong D^{1 \times p}/(D^{1 \times q'} R').$$

- $t(M) = 0$  iff  $(\star)$  is an exact sequence. If so and if  $\mathcal{F}$  is an injective cogenerator, then  $(\star)$  is exact iff so is the complex:

$$\mathcal{F}^q \xleftarrow{R.} \mathcal{F}^p \xleftarrow{Q.} \mathcal{F}^m, \text{ i.e., } \ker_{\mathcal{F}}(R.) = Q \mathcal{F}^m.$$

# Injective cogenerator modules

- Definition: A left  $D$ -module  $\mathcal{F}$  is injective if

$$\forall q \geq 1, \quad \forall R \in D^q, \quad \forall \zeta \in \ker_{\mathcal{F}}(S.),$$

where  $\ker_D(.R) = D^{1 \times r} S$ , there exists  $\eta \in \mathcal{F}$  satisfying  $R\eta = \zeta$ .

- Definition: If  $\mathcal{F}$  is a injective left  $D$ -module, then we have:

$$\text{ext}_D^i(M, \mathcal{F}) = 0, \quad \forall i \geq 1.$$

- Definition: A left  $D$ -module  $\mathcal{F}$  is cogenerator if:

$$\hom_D(M, \mathcal{F}) = 0 \Rightarrow M = 0.$$

- Proposition: Injective cogenerator left  $D$ -module always exists.
- Example: If  $\Omega$  is an open convex subset of  $\mathbb{R}^n$ , then the  $\mathbb{C} \left[ \partial_1; \text{id}, \frac{\partial}{\partial x_1} \right] \dots \left[ \partial_n; \text{id}, \frac{\partial}{\partial x_n} \right]$ -modules  $C^\infty(\Omega)$ ,  $\mathcal{D}'(\Omega)$ ,  $\mathcal{A}(\Omega)$ ,  $\mathcal{O}(\Omega)$  and  $\mathcal{B}(\Omega)$  are injective cogenerators.

# Involutions and adjoints

- Definition: A  $k$ -linear map  $\theta : D \longrightarrow D$  is an **involution** of  $D$  if:

$$\forall P, Q \in D : \theta(PQ) = \theta(Q)\theta(P), \quad \theta^2 = \text{id}.$$

- Example: 1. If  $D$  is a commutative ring, then  $\theta = \text{id}$ .

- 2. An involution of  $D = A \left[ \partial_1; \text{id}, \frac{\partial}{\partial x_1} \right] \dots \left[ \partial_n; \text{id}, \frac{\partial}{\partial x_n} \right]$  is:

$$\forall a \in A, \quad \theta(a(x)) = a(x), \quad \theta(\partial_i) = -\partial_i, \quad i = 1, \dots, n.$$

- 3. An involution of  $D = A \left[ \partial_1; \text{id}, \frac{d}{dt} \right] \left[ \partial_2; \delta, 0 \right]$  is defined by:

$$\forall a \in A, \quad \theta(a(t)) = a(-t), \quad \theta(\partial_i) = \partial_i, \quad i = 1, 2.$$

- The **adjoint** of  $R \in D^{q \times p}$  is defined by  $\theta(R) = (\theta(R_{ij}))^T \in D^{p \times q}$ .

- $\tilde{N} = D^{1 \times q} / (D^{1 \times p} \theta(R))$  is called the **adjoint** of  $M$ .

# Extension functor $\text{ext}_D^1(\cdot, D)$

- Parametrizability:  $Ry = 0 \iff \exists Q \in D^{p \times m} : y = Qz.$

$$4. \quad \theta(P)z = y \implies Ry = 0 \quad 1.$$

$$\begin{array}{ccc} \uparrow & & \downarrow \\ \text{involution } \theta & & \text{involution } \theta \\ \uparrow & & \downarrow \end{array}$$

$$3. \quad 0 = P\mu \stackrel{\text{Gb}}{\iff} \theta(R)\lambda = \mu \quad 2.$$

$$\begin{aligned} P \circ \theta(R) = 0 \implies \theta(P \circ \theta(R)) &= \theta^2(R) \circ \theta(P) \\ &= R \circ \theta(P) = 0. \end{aligned}$$

$$5. \quad \theta(P)z = y \stackrel{\text{Gb}}{\iff} R'y = 0, \quad R' \in D^{q' \times p}.$$

$$\boxed{\text{ext}_D^1(N, D) \cong (D^{1 \times q'} R') / (D^{1 \times q} R)}$$

- Using Gb, we can test whether or not  $\text{ext}_D^1(N, D) = 0$

# Wind tunnel model (Manitius, IEEE TAC 84)

1. The w.t.m. is defined by the **underdetermined system**:

$$\begin{pmatrix} \partial_1 + a & -k a \partial_2 & 0 & 0 \\ 0 & \partial_1 & -1 & 0 \\ 0 & \omega^2 & \partial_1 + 2\zeta\omega & -\omega^2 \end{pmatrix} \begin{pmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \\ u(t) \end{pmatrix} = 0.$$

2. We compute  $\theta(R) = R^T$  and define  $\theta(R)\lambda = \mu$ :

$$\left\{ \begin{array}{l} (\partial_1 + a) \lambda_1 = \mu_1, \\ -k a \partial_2 \lambda_1 + \partial_1 \lambda_2 + \omega^2 \lambda_3 = \mu_2, \\ -\lambda_2 + (\partial_1 + 2\zeta\omega) \lambda_3 = \mu_3, \\ -\omega^2 \lambda_3 = \mu_4. \end{array} \right. \quad (2)$$

(2) is **overdetermined**  $\xrightarrow{\text{Gb}}$  compatibility conditions  $P\mu = 0$ .

## Wind tunnel model (Manitius, IEEE TAC 84)

3. We obtain the compatibility condition  $P\mu = 0$ :

$$\begin{aligned} & \omega^2 k a \partial_2 \mu_1 + \omega^2 (\partial_1 - a) \mu_2 + \omega^2 (\partial_1^2 + a \partial_1) \mu_3 \\ & + (\partial_1^3 + 2\zeta\omega \partial_1^2 + a \partial_1^2 + \omega^2 \partial_1 + 2a\zeta\omega \partial_1 + a\omega^2) \mu_4 = 0. \end{aligned}$$

4. We consider the overdetermined system  $P^T z = y$ .

$$\left\{ \begin{array}{l} \omega^2 k a \partial_2 z = x_1, \\ \omega^2 (\partial_1 - a) z = x_2, \\ \omega^2 (\partial_1^2 + a \partial_1) z = x_3, \\ (\partial_1^3 + (2\zeta\omega + a) \partial_1^2 + (\omega^2 + 2a\omega\zeta) \partial_1 + a\omega) z = u. \end{array} \right. \quad (4)$$

5. The compatibility conditions of  $P^T z = y$  are exactly generated by  $Ry = 0$  and (4) is a parametrization of the w.t.m.

# Moving tank (Petit, Rouchon, IEEE TAC 02)

1. The model of a moving tank is defined by:

$$\begin{pmatrix} \partial_1 & -\partial_1 \partial_2^2 & a \partial_1^2 \partial_2 \\ \partial_1 \partial_2^2 & -\partial_1 & a \partial_1^2 \partial_2 \end{pmatrix} \begin{pmatrix} y_1(t) \\ y_2(t) \\ y_3(t) \end{pmatrix} = 0.$$

2. We compute  $\theta(R) = R^T$  and define  $\theta(R) \lambda = \mu$ :

$$\begin{cases} \partial_1 \lambda_1 + \partial_1 \partial_2^2 \lambda_2 = \mu_1, \\ -\partial_1 \partial_2^2 \lambda_1 - \partial_1 \lambda_2 = \mu_2, \\ a \partial_1^2 \partial_2 \lambda_1 + a \partial_1^2 \partial_2 \lambda_2 = \mu_3. \end{cases} \quad (2)$$

(2) is overdetermined  $\xrightarrow{\text{Gb}}$  compatibility conditions  $P \mu = 0$ .

# Moving tank (Petit, Rouchon, IEEE TAC 02)

3. We obtain the compatibility condition  $P\mu = 0$ :

$$a \partial_1 \partial_2 \mu_1 - a \partial_1 \partial_2 \mu_2 - (1 + \partial_2^2) \mu_3 = 0.$$

4. We consider the overdetermined system  $P^T z = y$ .

$$\begin{cases} a \partial_1 \partial_2 z = y_1, \\ -a \partial_1 \partial_2 z = y_2, \\ -(1 + \partial_2^2) z = y_3. \end{cases} \quad (4)$$

5. The compatibility conditions of  $P^T z = y$  are  $R'y = 0$ :

$$\begin{pmatrix} 1 & 1 & 0 \\ 0 & (1 + \partial_2^2) & -a \partial_1 \partial_2 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \\ y_3 \end{pmatrix} = 0.$$

# Moving tank (Petit, Rouchon, IEEE TAC 02)

$$t(M) \cong \text{ext}_D^1(N, D) \cong \\ \left( D^{1 \times 2} \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 + \partial_2^2 & -a\partial_1\partial_2 \end{pmatrix} \right) / \left( D^{1 \times 2} \begin{pmatrix} \partial_1 & -\partial_1\partial_2^2 & a\partial_1^2\partial_2 \\ \partial_1\partial_2^2 & -\partial_1 & a\partial_1^2\partial_2 \end{pmatrix} \right)$$

$$\begin{cases} y_1 + y_2 = z_1, \\ \partial_1 y_1 - \partial_1 \partial_2^2 y_2 + a \partial_1^2 \partial_2 y_3 = 0, \\ \partial_1 \partial_2^2 y_1 - \partial_1 y_2 + a \partial_1^2 \partial_2 y_3 = 0, \end{cases} \xrightarrow{\text{Gb}} \partial_1 (\partial_2^2 - 1) z_1 = 0.$$

$$\begin{cases} (1 + \partial_2^2) y_2 - a \partial_1 \partial_2 y_3 = z_2, \\ \partial_1 y_1 - \partial_1 \partial_2^2 y_2 + a \partial_1^2 \partial_2 y_3 = 0, \\ \partial_1 \partial_2^2 y_1 - \partial_1 y_2 + a \partial_1^2 \partial_2 y_3 = 0, \end{cases} \xrightarrow{\text{Gb}} \partial_1 (\partial_2^2 - 1) z_2 = 0.$$

$\Rightarrow z_1$  and  $z_2$  are non-trivial torsion elements.

# Examples

- 2D Stokes equations:

$$\begin{pmatrix} -\nu(\partial_x^2 + \partial_y^2) & 0 & \partial_x \\ 0 & -\nu(\partial_x^2 + \partial_y^2) & \partial_y \\ \partial_x & \partial_y & 0 \end{pmatrix} \begin{pmatrix} u \\ v \\ p \end{pmatrix} = 0$$
$$\Rightarrow \begin{cases} (\partial_x^2 + \partial_y^2)^2 u = 0, \\ (\partial_x^2 + \partial_y^2)^2 v = 0, \\ (\partial_x^2 + \partial_y^2) p = 0. \end{cases} \quad \text{torsion module}$$

- Moving tank (Petit, Rouchon, IEEE TAC 02):

$$\begin{cases} \dot{y}_1(t) - \dot{y}_2(t - 2h) + \alpha \ddot{y}_3(t - h) = 0, \\ \dot{y}_1(t - 2h) - \dot{y}_2(t) + \alpha \ddot{y}_3(t - h) = 0, \\ z_1(t) = y_1(t) + y_2(t), \\ z_2(t) = y_2(t) + y_2(t - 2h) - a \dot{y}_3(t - h), \\ \frac{d}{dt}(1 - \delta^2) z_i(t) = 0, \quad i = 1, 2. \end{cases} \quad \text{module with torsion}$$

## Examples: torsion-free modules

- Wind tunnel model (Manitius, IEEE TAC 84):

$$\begin{cases} \dot{x}_1(t) + a x_1(t) - k a x_2(t-h) = 0, \\ \dot{x}_2(t) - x_3(t) = 0, \\ \dot{x}_3(t) + \omega^2 x_2(t) + 2\zeta\omega x_3(t) - \omega^2 u(t) = 0, \end{cases}$$

$$\Leftrightarrow \begin{cases} x_1(t) = \omega^2 k a z(t-h), \\ x_2(t) = \omega^2 \dot{z}(t) - a \omega^2 z(t), \\ x_3(t) = \omega^2 \ddot{z}(t) + \omega^2 a \dot{z}(t), \\ u(t) = z(t)^{(3)} + (2\zeta\omega + a) \ddot{z}(t) + (\omega^2 + 2a\omega\zeta) \dot{z}(t) + a\omega z(t). \end{cases}$$

$\Rightarrow$  motion planning and tracking (Fliess et al).

- 2D stress tensor (elasticity theory):

$$\begin{cases} \partial_x \sigma^{11} + \partial_y \sigma^{12} = 0, \\ \partial_x \sigma^{12} + \partial_y \sigma^{22} = 0, \end{cases} \Leftrightarrow \begin{cases} \sigma^{11} = \partial_y^2 \lambda, \\ \sigma^{12} = -\partial_x \partial_y \lambda, \text{ Airy function } \lambda. \\ \sigma^{22} = \partial_x^2 \lambda, \end{cases}$$

## Examples: reflexive modules

- div-curl-grad:  $\vec{\nabla} \cdot \vec{B} = 0 \Leftrightarrow \vec{B} = \vec{\nabla} \wedge \vec{A}$ ,  $\vec{\nabla} \wedge \vec{A} = \vec{0} \Leftrightarrow \vec{A} = \vec{\nabla} f$ .
- First group of Maxwell equations:

$$\begin{cases} \frac{\partial \vec{B}}{\partial t} + \vec{\nabla} \wedge \vec{E} = \vec{0}, \\ \vec{\nabla} \cdot \vec{B} = 0, \end{cases} \Leftrightarrow \begin{cases} \vec{E} = -\frac{\partial \vec{A}}{\partial t} - \vec{\nabla} V, \\ \vec{B} = \vec{\nabla} \wedge \vec{A}. \end{cases}$$

$$\begin{cases} -\frac{\partial \vec{A}}{\partial t} - \vec{\nabla} V = \vec{0}, \\ \vec{\nabla} \wedge \vec{A} = \vec{0}, \end{cases} \Leftrightarrow \begin{cases} \vec{A} = \vec{\nabla} \xi, \\ V = -\frac{\partial \xi}{\partial t}. \end{cases}$$

- 3D stress tensor: Maxwell, Morera parametrizations...
- Linearized Einstein equations (system of PDEs  $10 \times 10$ )?

⇒ **OREMODULES** (Chyzak, Q., Robertz)

<http://wwwb.math.rwth-aachen.de/OreModules/>

# Dictionary systems – modules

Module $M$	Structural properties $\ker_{\mathcal{F}}(R.)$	Stabilization problems Optimal control
Torsion	Autonomous system Poles/zeros classifications	
With torsion	Existence of autonomous elements	
Torsion-free	No autonomous elements, Controllability, Parametrizability, $\pi$ -freeness	Variational problem without constraints (Euler-Lagrange equations)
Projective	Bézout identities, Internal stabilizability	Computation of Lagrange parameters without integration Existence of a parametrization all stabilizing controllers
Free	Flatness, Poles placement, Doubly coprime factorization	Youla-Kučera parametrization Optimal controller

# Projectiveness, observability and controllability

- **Theorem:** If  $R \in D^{q \times p}$  has full row rank, then the left  $D$ -module  $M = D^{1 \times p}/(D^{1 \times q} R)$  is projective iff:

$$N = D^{1 \times q}/(D^{1 \times p} \theta(R)) = 0 \Leftrightarrow \exists S \in D^{p \times q} : RS = I_q.$$

- Let  $D = \mathcal{A}(I) [\partial; \text{id}, \frac{d}{dt}]$  and  $R = (\partial I_n - A \quad -B) \in D^{n \times (n+m)}$ .  
 $M = D^{1 \times (n+m)}/(D^{1 \times n} R)$  is projective iff  $\theta(R)\lambda = 0 \Leftrightarrow \lambda = 0$ :

$$\begin{cases} -\partial\lambda - A^T\lambda = 0, \\ -B^T\lambda = 0, \end{cases} \Rightarrow \begin{cases} \partial\lambda = -A^T\lambda, \\ B^T\lambda = 0, \\ B^T\partial\lambda + \dot{B}^T\lambda = (-B^TA^T + \dot{B}^T)\lambda = 0. \end{cases}$$

Hence,  $M$  is projective iff, for all  $t_0 \in I$ , we have:

$$\text{rank}_{\mathbb{R}}(B \mid AB - \dot{B} \mid A^2B + \dots \mid A^{n-1}B + \dots \mid \dots)(t_0) = n.$$

- $D^{1 \times p}/(D^{1 \times q} (P(\partial) - Q(\partial)))$  proj. iff  $P(\partial)X(\partial) + Q(\partial)Y(\partial) = I_q$ .

# Controllability à la Willems and flatness

- We consider two pendula mounted on a car:

$$\begin{cases} m_1 L_1 \ddot{w}_1(t) + m_2 L_2 \ddot{w}_2(t) - w_3(t) + (M + m_1 + m_2) \ddot{w}_4(t) = 0, \\ m_1 L_1^2 \ddot{w}_1(t) - m_1 L_1 g w_1(t) + m_1 L_1 \ddot{w}_4(t) = 0, \\ m_2 L_2^2 \ddot{w}_2(t) - m_2 L_2 g w_2(t) + m_2 L_2 \ddot{w}_4(t) = 0. \end{cases} \quad (*)$$

- (\*) is parametrizable iff  $L_1 \neq L_2$ .
- If  $L_1 \neq L_2$  then a parametrization of (\*) is defined by:

$$\begin{cases} w_1(t) = -L_2 \xi^{(4)}(t) + g \ddot{\xi}(t), \\ w_2(t) = -L_1 \xi^{(4)}(t) + g \ddot{\xi}(t), \\ w_3(t) = L_1 L_2 M \xi^{(6)}(t) - (L_1 m_2 + L_2 m_1 + g(L_1 + L_2) M) \xi^{(4)}(t) \\ \quad + g^2 (m_1 + m_2 + M) \xi^{(2)}(t) \\ w_4(t) = L_1 L_2 \xi^{(4)}(t) - g(L_1 + L_2) \ddot{\xi}(t) + g^2 \xi(t). \end{cases}$$

- The parametrization of (\*) is injective as we have:

$$\xi(t) = \frac{1}{g^2(L_1 - L_2)} (L_1^2 w_1(t) - L_2^2 w_2(t) + (L_1 - L_2) w_4(t)). \quad (**)$$

# Controllability à la Willems and flatness

- Patching problem  $\Leftrightarrow$  controllability:  $T \geq 0$ .

$w^P = (w_1^P, w_2^P, w_3^P, w_4^P)$  a past trajectory of  $(\star)$  on  $]-\infty, 0]$ .

$w^f = (w_1^f, w_2^f, w_3^f, w_4^f)$  a future trajectory of  $(\star)$  on  $[T, +\infty[$ .

$\Rightarrow \exists w = (w_1, w_2, w_3, w_4) \in C^\infty(\mathbb{R})^4$  trajectory of  $(\star)$ :

$$\begin{cases} w_{]-\infty, 0]} = w^P, \\ w_{[T, +\infty[} = w^f. \end{cases}$$

- Using the flat output

$$\xi(t) = \frac{1}{g^2(L_1 - L_2)} (L_1^2 w_1(t) - L_2^2 w_2(t) + (L_1 - L_2) w_4(t))$$

and the parametrization, it is enough to find  $\xi \in C^\infty(\mathbb{R})$  s.t.:

$$\xi_{]-\infty, 0]} = \xi^P \quad \& \quad \xi_{[T, +\infty[} = \xi^f.$$

# Motion planning

- Flexible rod with a torque (Mounier 95):

$$\begin{cases} \sigma^2 \frac{\partial^2 q(\tau, x)}{\partial \tau^2} - \frac{\partial^2 q(\tau, x)}{\partial x^2} = 0, \\ \frac{\partial q}{\partial x}(\tau, 0) = -u(\tau), \\ \frac{\partial q}{\partial x}(\tau, L) = -J \frac{\partial^2 q}{\partial \tau^2}(\tau, L), \\ y(\tau) = q(\tau, L). \end{cases} \quad (*)$$

- $q(\tau, x) = \phi(\tau + \sigma x) + \psi(\tau - \sigma x)$ ,  $t = (\sigma/J)\tau$ ,  $v = (2J/\sigma^2)u$ ,

$$(*) \Rightarrow \ddot{y}(t+1) + \ddot{y}(t-1) + \dot{y}(t+1) - \dot{y}(t-1) = v(t)$$

$$\Leftrightarrow \begin{cases} y(t) = \xi(t-1), \\ v(t) = \ddot{\xi}(t) + \ddot{\xi}(t-2) + \dot{\xi}(t) - \dot{\xi}(t-2). \end{cases}$$

- If  $y_r$  is a desired trajectory then  $\xi_r(t) = y_r(t+1)$  and we obtain the open-loop control law:

$$v_r(t) = \ddot{y}_r(t+1) + \ddot{y}_r(t-1) + \dot{y}_r(t+1) - \dot{y}_r(t-1).$$

# Optimal control

- Let us minimize  $\frac{1}{2} \int_0^T (x(t)^2 + u(t)^2) dt$  (1) under:

$$\dot{x}(t) + x(t) - u(t) = 0, \quad x(0) = x_0. \quad (2)$$

- (2) is parametrized by  $\begin{cases} x(t) = \xi(t), \\ u(t) = \dot{\xi}(t) + \xi(t). \end{cases}$  (3)

- (1) & (3)  $\Rightarrow \min \frac{1}{2} \int_0^T (\xi(t)^2 + (\dot{\xi}(t) + \xi(t))^2) dt,$

$$\Rightarrow \text{Euler-Lagrange equations} \quad \begin{cases} \xi(t) = x(t), \\ \dot{\xi}(t) + \xi(t) = u(t), \\ \ddot{\xi}(t) - 2\xi(t) = 0, \\ \dot{\xi}(T) + \xi(T) = 0, \\ \xi(0) = x_0, \end{cases}$$

$$\Rightarrow u(t) = \frac{-e^{\sqrt{2}(t-T)} + e^{-\sqrt{2}(t-T)}}{(1 - \sqrt{2}) e^{\sqrt{2}(t-T)} - (1 + \sqrt{2}) e^{-\sqrt{2}(t-T)}} x(t).$$

# Variational problems

- Let us extremize the electromagnetic action

$$\int \left( \frac{1}{2\mu_0} \parallel \vec{B} \parallel^2 - \frac{\epsilon_0}{2} \parallel \vec{E} \parallel^2 \right) dx_1 dx_2 dx_3 dt, \quad (1)$$

where  $\vec{B}$  and  $\vec{E}$  satisfy the first group of Maxwell equations:

$$\begin{cases} \frac{\partial \vec{B}}{\partial t} + \vec{\nabla} \wedge \vec{E} = \vec{0}, \\ \vec{\nabla} \cdot \vec{B} = 0, \end{cases} \Leftrightarrow \begin{cases} \vec{E} = -\vec{\nabla} V - \frac{\partial \vec{A}}{\partial t}, \\ \vec{B} = \vec{\nabla} \wedge \vec{A}. \end{cases} \quad (3)$$

- Substituting (3) in (1) and using Lorentz gauge

$$\vec{\nabla} \cdot \vec{A} + \frac{1}{c^2} \frac{\partial V}{\partial t} = 0, \quad c^2 = 1/(\epsilon_0 \mu_0),$$

$$\Rightarrow \begin{cases} \frac{1}{c^2} \frac{\partial^2 \vec{A}}{\partial t^2} - \Delta \vec{A} = 0, \\ \frac{1}{c^2} \frac{\partial^2 V}{\partial t^2} - \Delta V = 0. \end{cases} \quad (\text{electromagnetic waves}).$$

# OREMODULES (Chyzak, Q., Robertz)

- OREMODULES is a tool-box developed in *Maple*.
- OREMODULES uses *Ore-algebra* developed by Chyzak.
- OREMODULES handles linear systems of ODEs, PDEs, difference equations, differential time-delay equations . . .
- OREMODULES computes:
  1. free resolutions,  $\text{ext}_D^i(\cdot, D)$ , projective dim., Hilbert series,
  2. torsion elements, autonomous elements,
  3. parametrizations of underdetermined linear functional systems,
  4. left-/right-/generalized inverses,
  5.  $\pi$ -polynomials, bases, flat outputs,
  6. first integrals of motion, Euler-Lagrange equations . . .

<http://wwwb.math.rwth-aachen.de/OreModules/>

# Conclusion

- Based on algebraic analysis, module theory, constructive homological algebra and Ore algebras, we have developed a general non-commutative polynomial approach to functional linear systems.
- The different results are implemented in OREMODULES.

This new approach allowed us to:

- ① Develop an intrinsic approach (independent of the form).
- ② Develop generic algorithms and generic implementations.
- ③ Constructively study the parametrizability problem.
- ④ Solve conjectures in mathematical systems theory.

F. Chyzak, A. Quadrat, D. Robertz, “Effective algorithms for parametrizing linear control systems over Ore algebras”, *Applicable Algebra in Engineering, Communications and Computing*, 16 (2005), 319-376.