

Quillen-Suslin theorem: algorithms and applications

Alban Quadrat

INRIA Sophia Antipolis, APICS Project, 2004 route des lucioles, BP 93,
06902 Sophia Antipolis cedex, France.

`Alban.Quadrat@sophia.inria.fr`

`http://www-sop.inria.fr/members/Alban.Quadrat/index.html`

KIAS Winter School on Algebraic Systems Theory

KIAS, Seoul, 15/12/08

Introduction

- In Talk 1, we have shown how to check whether or not a left module $M = D^{1 \times p} / (D^{1 \times q} R)$ over a certain Ore algebra D was:
 - ① torsion,
 - ② admits non-trivial torsion elements,
 - ③ torsion-free,
 - ④ reflexive,
 - ⑤ projective.
- In this talk, we study stably free and free modules.
- We shall focus on the so-called Quillen-Suslin theorem proving that projective modules over $D = k[x_1, \dots, x_n]$ (k a field) are free.
 - ① Constructive computation of bases of free D -modules:
package `QUILLEN``SUSLIN`.
 - ② Applications in mathematical systems theory.

Shorter free resolutions

- **Theorem:** Let us consider a **finite free resolution** of M :

$$0 \longrightarrow D^{1 \times p_m} \xrightarrow{\cdot R_m} D^{1 \times p_{m-1}} \xrightarrow{\cdot R_{m-1}} \dots \xrightarrow{\cdot R_2} D^{1 \times p_1} \xrightarrow{\cdot R_1} D^{1 \times p_0} \xrightarrow{\pi} M \longrightarrow 0.$$

1. If $m \geq 3$ and there exists $S_m \in D^{p_{m-1} \times p_m}$ such that $R_m S_m = I_{p_m}$, then we have the **finite free resolution of M** :

$$0 \longrightarrow D^{1 \times p_{m-1}} \xrightarrow{\cdot T_{m-1}} D^{1 \times (p_{m-2} + p_m)} \xrightarrow{\cdot T_{m-2}} D^{1 \times p_{m-3}} \xrightarrow{\cdot R_{m-3}} \dots \xrightarrow{\pi} M \longrightarrow 0, \quad (1)$$

$$\text{where } T_{m-1} = \begin{pmatrix} R_{m-1} & S_m \end{pmatrix}, \quad T_{m-2} = \begin{pmatrix} R_{m-2} \\ 0 \end{pmatrix}.$$

2. If $m = 2$ and there exists $S_2 \in D^{p_1 \times p_2}$ such that $R_2 S_2 = I_{p_2}$, then we have the **finite free resolution**

$$0 \longrightarrow D^{1 \times p_1} \xrightarrow{\cdot T_1} D^{1 \times (p_0 + p_2)} \xrightarrow{\tau} M \longrightarrow 0, \quad (2)$$

$$\text{where } T_1 = \begin{pmatrix} R_1 & S_2 \end{pmatrix} \text{ and } \tau = \begin{pmatrix} \pi \\ 0 \end{pmatrix}.$$

Example: annihilator of δ

- δ satisfies the system: $t^2 y(t) = 0$, $t \dot{y}(t) + 2y(t) = 0$.
- We consider $D = \mathbb{Q}[t] \left[\partial; \text{id}, \frac{d}{dt} \right]$ and the left D -module:

$$M = D / (D t^2 + D (t \partial + 2)).$$

- M admits the following **finite free resolution** of M :

$$0 \longrightarrow D \xrightarrow{\cdot R_2} D^{1 \times 2} \xrightarrow{\cdot R_1} D \xrightarrow{\pi} M \longrightarrow 0,$$

$$R_1 = \begin{pmatrix} t^2 & t \partial + 2 \end{pmatrix}^T, \quad R_2 = \begin{pmatrix} \partial & -t \end{pmatrix}.$$

- $S_2 = \begin{pmatrix} t & \partial \end{pmatrix}^T$ is a **right-inverse** of R_2 , and thus, we get:

$$0 \longrightarrow D^{1 \times 2} \xrightarrow{\cdot T_1} D^{1 \times 2} \xrightarrow{\tau} M \longrightarrow 0, \quad T_1 = \begin{pmatrix} t^2 & t \\ t \partial + 2 & \partial \end{pmatrix}.$$

Example: contact transformations

- $D = A_3(\mathbb{Q}) = \mathbb{Q}[x_1, x_2, x_3] \left[\partial_1; \text{id}, \frac{\partial}{\partial x_1} \right] \left[\partial_2; \text{id}, \frac{\partial}{\partial x_2} \right] \left[\partial_3; \text{id}, \frac{\partial}{\partial x_3} \right],$

$$R_1 = \begin{pmatrix} \frac{1}{2} x_2 \partial_1 & x_2 \partial_2 + 1 & x_2 \partial_3 + \frac{1}{2} \partial_1 \\ -\frac{1}{2} x_2 \partial_2 - \frac{3}{2} & 0 & \frac{1}{2} \partial_2 \\ -\partial_1 - \frac{1}{2} x_2 \partial_3 & -\partial_2 & -\frac{1}{2} \partial_3 \end{pmatrix} \in D^{3 \times 3}.$$

- If $R_2 = (\partial_2 \quad -(\partial_1 + x_3 \partial_3) \quad x_2 \partial_2 + 2)$, then the left D -module $M = D^{1 \times 3} / (D^{1 \times 3} R_1)$ admits the **finite free resolution**:

$$0 \longrightarrow D \xrightarrow{\cdot R_2} D^{1 \times 3} \xrightarrow{\cdot R_1} D^{1 \times 3} \xrightarrow{\pi} M \longrightarrow 0.$$

- $S_2 = (-x_2 \quad 0 \quad 1)^T$ is a **right-inverse** of R_2 and we get:

$$0 \longrightarrow D^{1 \times 3} \xrightarrow{\cdot T_1} D^{1 \times 4} \xrightarrow{\tau} M \longrightarrow 0,$$

$$T_1 = \begin{pmatrix} \frac{1}{2} x_2 \partial_1 & x_2 \partial_2 + 1 & x_2 \partial_3 + \frac{1}{2} \partial_1 & -x_2 \\ -\frac{1}{2} x_2 \partial_2 - \frac{3}{2} & 0 & \frac{1}{2} \partial_2 & 0 \\ -\partial_1 - \frac{1}{2} x_2 \partial_3 & -\partial_2 & -\frac{1}{2} \partial_3 & 1 \end{pmatrix}.$$

Projective dimensions

- **Definition:** A **projective resolution** of a left D -module M is an exact sequence of the form

$$0 \longrightarrow P_n \xrightarrow{\delta_n} P_{n-1} \xrightarrow{\delta_{n-1}} \dots \xrightarrow{\delta_1} P_0 \xrightarrow{\delta_0} M \longrightarrow 0, \quad (\star)$$

where the P_i 's are projective left D -modules.

- **Definition:** We call **left projective dimension** of a left D -module M , denoted by $\text{lpd}_D(M)$, the smallest n such that there exists a **projective resolution** of the form (\star) .
- **Proposition:** $\text{lpd}_D(M) = n$ iff **there exists** a finite projective resolution (\star) of M , where δ_n is **nonsplit**, i.e., there exists no D -morphism $\tau_n : P_{n-1} \longrightarrow P_n$ such that $\tau_n \circ \delta_n = \text{id}_{P_n}$, with the convention $P_{-1} = M$.

Computation of left projective dimensions

- **Algorithm:** 1. Compute a finite free resolution of M .
- 2. Set $j = m$ and $T_j = R_m$.
- 3. Check if R_j admits a right-inverse S_j over D .
- ⇒ If not, then exit and $\text{lpd}_D(M) = j$.
- ⇒ If yes and:
 - (a) If $j = 1$, then exit with $\text{lpd}_D(M) = 0$.
 - (b) If $j = 2$, then compute (2) and return to 3 with $j \leftarrow j - 1$.
 - (c) If $j \geq 3$, then compute (1) and return to 3 with $j \leftarrow j - 1$.
- **Example:** The left $A_1(\mathbb{Q})$ -module M associated with the annihilator of $\dot{\delta}$ has $\text{lpd}_D(M) = 1$.
- **Example:** The left $A_3(\mathbb{Q})$ -module M associated with the contact transformations has $\text{lpd}_D(M) = 0$.

Shortest free resolutions

- If M is a **projective** left D -module, then $\text{lpd}_D(M) = 0$.
- Moreover, if M admits a finite free resolution

$$0 \longrightarrow D^{1 \times p_m} \xrightarrow{\cdot R_m} D^{1 \times p_{m-1}} \xrightarrow{\cdot R_{m-1}} \dots \xrightarrow{\cdot R_2} D^{1 \times p_1} \xrightarrow{\cdot R_1} D^{1 \times p_0} \xrightarrow{\pi} M \longrightarrow 0,$$

then the previous algorithm returns a matrix $R \in D^{q \times p}$ such that

$$0 \longrightarrow D^{1 \times q} \xrightarrow{\cdot R} D^{1 \times p} \xrightarrow{\pi} M \longrightarrow 0 \quad (\star)$$

is a **split finite free resolution of M** , i.e., (\star) is exact and R admits a right-inverse $S \in D^{p \times q}$, i.e., $RS = I_q$.

In particular, we have $D^{1 \times p} \cong M \oplus D^{1 \times q}$, which proves that M is a **stably free left D -module** (Serre's theorem).

- The matrix R will be called a **minimal presentation matrix of M** .

Example: contact transformations

- We consider the left $D = A_3(\mathbb{Q})$ -module $M = D^{1 \times 3} / (D^{1 \times 3} R_1)$:

$$R_1 = \begin{pmatrix} \frac{1}{2} x_2 \partial_1 & x_2 \partial_2 + 1 & x_2 \partial_3 + \frac{1}{2} \partial_1 \\ -\frac{1}{2} x_2 \partial_2 - \frac{3}{2} & 0 & \frac{1}{2} \partial_2 \\ -\partial_1 - \frac{1}{2} x_2 \partial_3 & -\partial_2 & -\frac{1}{2} \partial_3 \end{pmatrix}.$$

- M is a **stably free** left D -module defined by the **minimal presentation matrix** T_1 : $0 \longrightarrow D^{1 \times 3} \xrightarrow{T_1} D^{1 \times 4} \xrightarrow{\tau} M \longrightarrow 0$,

$$T_1 = \begin{pmatrix} \frac{1}{2} x_2 \partial_1 & x_2 \partial_2 + 1 & x_2 \partial_3 + \frac{1}{2} \partial_1 & -x_2 \\ -\frac{1}{2} x_2 \partial_2 - \frac{3}{2} & 0 & \frac{1}{2} \partial_2 & 0 \\ -\partial_1 - \frac{1}{2} x_2 \partial_3 & -\partial_2 & -\frac{1}{2} \partial_3 & 1 \end{pmatrix},$$

$$S_1 = \begin{pmatrix} 0 & -1 & 0 \\ 1 & 0 & x_2 \\ 0 & -x_2 & 0 \\ \partial_2 & -\partial_1 - x_2 \partial_3 & x_2 \partial_2 + 2 \end{pmatrix}, \quad T_1 S_1 = I_3.$$

Existence of finite free resolutions

• **Theorem:** Let A be a left noetherian ring of finite left projective dimension $\text{lgld}(A)$ and whose finitely generated projective left A -module are stably free. Let $D = A[\partial_1; \alpha_1, \beta_1] \dots [\partial_m; \alpha_m, \beta_m]$ be an Ore algebra where the α_i 's are automorphisms. Then, we have:

- 1 Every finitely generated left D -module admits a finite free resolution of length $\text{lpd}(D) + 1$.
 - 2 Every finitely generated projective left D -module is stably free.
- **Example:** $D = A[x_1, \dots, x_n]$, where A is a principal ideal domain (e.g., $A = \mathbb{Z}$, $A = k$ a field).
- **Example:** The Weyl algebras $D = A_n(k)$ and $B_n(k)$.

Characterization of free modules

- Let M be a stably free left D -module defined by a minimal presentation matrix R :

$$0 \longrightarrow D^{1 \times q} \xrightarrow{\cdot R} D^{1 \times p} \xrightarrow{\pi} M \longrightarrow 0, \quad RS = I_q.$$

- $\mathrm{GL}_p(D) = \{U \in D^{p \times p} \mid \exists V \in D^{p \times p} : UV = VU = I_p\}$.
- Theorem:** Let $R \in D^{q \times p}$ be a matrix admitting a right-inverse over D . Then, the left D -module $M = D^{1 \times p} / (D^{1 \times q} R)$ is free of rank $p - q$ iff there exists $U \in \mathrm{GL}_p(D)$ satisfying:

$$RU = (I_q \quad 0).$$

Then, $U^{-1} = \begin{pmatrix} R \\ T \end{pmatrix}$, where $T \in D^{(p-q) \times p}$, and the family $\{\pi(T_i \bullet)\}_{i=1, \dots, p-q}$ forms a basis of the free left D -module M .

- Let us suppose that there exists $U \in GL_p(D)$ such that:

$$RU = (I_q \ 0).$$

- We obtain the following commutative exact diagram

$$\begin{array}{ccccccc}
 & 0 & & 0 & & & \\
 & \downarrow & & \downarrow & & & \\
 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M \longrightarrow 0 \\
 & & \parallel & & \downarrow \cdot U & & \\
 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot (I_q \ 0)} & D^{1 \times p} & \xrightarrow{\cdot \begin{pmatrix} 0 \\ I_{p-q} \end{pmatrix}} & D^{1 \times (p-q)} \longrightarrow 0, \\
 & & \downarrow & & \downarrow & & \\
 & & 0 & & 0 & &
 \end{array}$$

which proves that $M \cong D^{1 \times (p-q)}$, i.e., M is a free of rank $p - q$.

Proof

- Let M be a **free left D -module**, i.e., $\phi : M \xrightarrow{\cong} D^{1 \times (p-q)}$.
- We have the following **exact commutative diagram** ($\cdot Q = \phi \circ \pi$)

$$\begin{array}{ccccccc}
 & & \xleftarrow{\cdot S} & & \xleftarrow{\cdot T} & & \\
 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\cdot Q} & D^{1 \times (p-q)} \longrightarrow 0 \\
 & & \parallel & & \parallel & & \uparrow \phi \\
 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M \longrightarrow 0,
 \end{array}$$

where the **first horizontal exact sequence splits**, namely:

$$RQ = 0, \quad RS = I_q, \quad TQ = I_{p-q}, \quad TS = 0, \quad SR + QT = I_p,$$

$$\text{i.e., } \begin{pmatrix} R \\ T \end{pmatrix} (S \quad Q) = \begin{pmatrix} I_q & 0 \\ 0 & I_{p-q} \end{pmatrix} = I_p, \quad (S \quad Q) \begin{pmatrix} R \\ T \end{pmatrix} = I_p.$$

$$\begin{array}{ccccccc}
 & & \xleftarrow{\cdot S} & & \xleftarrow{\cdot T} & & \\
 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\cdot Q} & D^{1 \times (p-q)} \longrightarrow 0 \\
 & & \parallel & & \parallel & & \uparrow \phi \downarrow \phi^{-1} \\
 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M \longrightarrow 0,
 \end{array}$$

The **isomorphism** ϕ is defined by:

$$\begin{array}{ll}
 \phi : M & \longrightarrow D^{1 \times (p-q)} & \phi^{-1} : D^{1 \times (p-q)} & \longrightarrow M \\
 \pi(\lambda) & \longmapsto \lambda Q, & \mu & \longmapsto \pi(\mu T).
 \end{array}$$

- If we denote by $\{h_k\}_{k=1, \dots, p-q}$ the standard basis of $D^{1 \times (p-q)}$, then $\{\phi^{-1}(h_k) = \pi(h_k T) = \pi(T_{k \bullet})\}_{k=1, \dots, (p-q)}$ is a basis of M

\Rightarrow the residue classes of the rows of T in M define a basis of M .

Injective parametrizations

- Let $\{f_j\}_{j=1,\dots,p}$ be the standard basis of $D^{1 \times p}$ and $\{y_j = \pi(f_j)\}_{j=1,\dots,p}$ a family of generators of $M = D^{1 \times p} / (D^{1 \times q} R)$.
- For $j = 1, \dots, p$, we have

$$y_j = \phi^{-1}(\phi(y_j)) = \phi^{-1}(f_j Q) = \phi^{-1} \left(\sum_{k=1}^{p-q} Q_{jk} h_k \right) = \sum_{k=1}^{p-q} Q_{jk} z_k, \quad (*)$$

which shows that Q defines a parametrization of M .

- The elements $z_k = \phi^{-1}(h_k) = \pi(T_{k\bullet})$ of the basis of M satisfy

$$z_k = \pi \left(\sum_{j=1}^p T_{kj} f_j \right) = \sum_{j=1}^p T_{kj} \pi(f_j) = \sum_{j=1}^p T_{kj} y_j,$$

which proves that $(*)$ is an injective parametrization of M .

Injective parametrizations

- Let $M = D^{1 \times p} / (D^{1 \times q} R)$ be a torsion-free left D -module.
- From the vanishing of $\text{ext}_D^1(N, D)$, where $N = D^q / (R D^p)$ is the Auslander transpose of M , we obtain $Q \in D^{p \times m}$ such that

$$\begin{array}{ccccccc}
 D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M & \longrightarrow & 0 \\
 0 \longleftarrow N \xleftarrow{\kappa} & D^q & \xleftarrow{R \cdot} & D^p & \xleftarrow{Q \cdot} & D^m & \\
 D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\cdot Q} & D^{1 \times m}, & &
 \end{array}$$

$$\text{ext}_D^1(N, D) = \ker_D(\cdot Q) / (D^{1 \times q} R) = 0 \Leftrightarrow M \cong D^{1 \times p} Q \subseteq D^{1 \times m}.$$

- If $\cdot Q : D^{1 \times p} \longrightarrow D^{1 \times m}$ is **surjective**, i.e., Q admits a left-inverse $T \in D^{m \times p}$, i.e., $T Q = I_m$, then we have

$$M \cong D^{1 \times p} Q = D^{1 \times m},$$

i.e., M is a **free** left D -module of rank m , $\{\pi(T_{k\bullet})\}_{k=1, \dots, m}$ is a **basis** and Q an **injective parametrization of M** .

Example: contact transformations

- We consider the left $D = A_3(\mathbb{Q})$ -module $M = D^{1 \times 3} / (D^{1 \times 3} R)$:

$$R = \begin{pmatrix} \frac{1}{2} x_2 \partial_1 & x_2 \partial_2 + 1 & x_2 \partial_3 + \frac{1}{2} \partial_1 \\ -\frac{1}{2} x_2 \partial_2 - \frac{3}{2} & 0 & \frac{1}{2} \partial_2 \\ -\partial_1 - \frac{1}{2} x_2 \partial_3 & -\partial_2 & -\frac{1}{2} \partial_3 \end{pmatrix}.$$

- Checking the vanishing of $\text{ext}_D^1(\tilde{N}, D) = 0$, we obtain that

$$Q = (-\partial_2 \quad x_2 \partial_3 + \partial_1 \quad -(x_2 \partial_2 + 2))^T$$

defines a **parametrization of M** , i.e., $M \cong D^{1 \times 3} Q \subseteq D$.

- Q admits the **left-inverse $T = \frac{1}{2}(x_2 \quad 0 \quad -1)$** , which proves that $M \cong D^{1 \times 3} Q = D$ and $z = \pi(T)$ is a **basis of M** .
- The generators $\{y_i = \pi(f_j)\}_{j=1,2,3}$ of M satisfying the relations $Ry = 0$, where $y = (y_1, y_2, y_3)^T$, satisfy $y = Qz$ and $z = Ty$:

$$y_1 = -\partial_2 z, \quad y_2 = (x_2 \partial_3 + \partial_1) z, \quad y_3 = -(x_2 \partial_2 + 2) z,$$
$$z = \frac{1}{2}(x_2 y_1 - y_3).$$

Minimal parametrizations

- We generally have $m \geq \text{rank}_D(M)$ (e.g., $\ker_D(Q) \neq 0$).
- **Theorem:** Let $M = D^{1 \times p} / (D^{1 \times q} R)$ be a torsion-free left D -module. Then, there exists a parametrization $P \in D^{p \times l}$ of M such that $l = \text{rank}_D(M)$, i.e., $L = D^{1 \times l} / (D^{1 \times p} P)$ is torsion.
- P is called a **minimal parametrization** of M .
- **Algorithm:** P can be obtained by selecting $\text{rank}_D(M)$ right D -linearly independent columns of a parametrization Q of M .
- **Heuristic method** for computing basis of a free left D -modules:
 - 1 Compute $Q \in D^{p \times m}$ such that $\ker_D(R) = Q D^m$.
 - 2 Define $P \in D^{p \times \text{rank}_D(M)}$ by selecting $\text{rank}_D(M)$ right D -linearly independent columns of Q .
 - 3 Check whether or not P admits a left-inverse over D .

Example

We consider the $D = \mathbb{Q}[x_1, x_2]$ -module $M = D^{1 \times 3} / (D R)$, where:

$$R = (x_1 x_2^2 + 1 \quad 3x_2/2 + x_1 - 1 \quad 2x_1 x_2), \quad \text{rank}_D(M) = 2.$$

- Checking $\text{ext}_D^1(D/(D^{1 \times 3} R^T), D) = 0$, we obtain that

$$Q = \begin{pmatrix} -4x_1 - 6x_2 + 4 & 6x_2^2 - 4x_2 & 0 \\ 4 & -4x_2 & 4x_1 x_2 \\ 2x_1 x_2 + 3x_2^2 - 2x_2 & -3x_2^3 + 2x_2^2 + 2 & -2x_1 - 3x_2 + 2 \end{pmatrix},$$

is a **parametrization of M** , i.e., $M \cong D^{1 \times 3} Q \subseteq D^{1 \times 3}$.

- Selecting the first two columns of Q , we obtain that

$$P = \begin{pmatrix} -4x_1 - 6x_2 + 4 & 6x_2^2 - 4x_2 \\ 4 & -4x_2 \\ 2x_1 x_2 + 3x_2^2 - 2x_2 & -3x_2^3 + 2x_2^2 + 2 \end{pmatrix},$$

is a **minimal parametrization of M** , i.e.:

$$M \cong D^{1 \times 3} Q \subseteq D^{1 \times 2} = D^{1 \times \text{rank}_D(M)}.$$

Example

- The parametrization P of M admits a **left-inverse** defined by:

$$T = \frac{1}{4} \begin{pmatrix} x_2^2 & 1 & 2x_2 \\ x_2 & 0 & 2 \end{pmatrix}.$$

- The set of generators $\{y_j = \pi(f_j)\}_{j=1,2,3}$ of M satisfies

$$(x_1 x_2^2 + 1) y_1 + (3 x_2/2 + x_1 - 1) y_2 + 2 x_1 x_2 y_3 = 0,$$

and a basis $\{z_1, z_2\}$ of M is defined by:

$$\begin{cases} z_1 = \frac{1}{4} (x_2^2 y_1 + y_2 + 2 x_2 y_3), \\ z_2 = \frac{1}{4} (x_2 y_1 + 2 y_3). \end{cases}$$

- The generators $\{y_1, y_2, y_3\}$ can be written in the basis $\{z_1, z_2\}$:

$$\begin{cases} y_1 = (-4 x_1 - 6 x_2 + 4) z_1 + (6 x_2^2 - 4 x_2) z_2, \\ y_2 = 4 z_1 - 4 x_2 z_2, \\ y_3 = (2 x_1 x_2 + 3 x_2^2 - 2 x_2) z_1 + (-3 x_2^3 + 2 x_2^2 + 2) z_2. \end{cases}$$

Computation of bases of general free modules

• Let $P \in D^{p \times m}$ and $D^{1 \times p} \xrightarrow{\cdot P} D^{1 \times m}$.

1. If $U = D^{1 \times p} P$ is free, then compute $R \in D^{q \times p}$ such that:

$$0 \longrightarrow D^{1 \times q} \xrightarrow{\cdot R} D^{1 \times p} \xrightarrow{\cdot P} D^{1 \times m} \quad \text{is exact.}$$

$$\begin{array}{ccccccc} 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\cdot P} & D^{1 \times p} P \longrightarrow 0 \\ \Rightarrow & & \parallel & & \parallel & & \uparrow \psi \\ 0 & \longrightarrow & D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M \longrightarrow 0, \end{array}$$

where $\psi(\pi(\lambda)) = \lambda P$, $\forall \lambda \in D^{1 \times p}$. We get $U = \psi(M)$ and:

$$\Rightarrow U = D^{1 \times (p-q)} (T P), \quad \text{where} \quad \begin{pmatrix} R \\ T \end{pmatrix} (S \quad Q) = I_p.$$

2. If $V = \ker_D(\cdot P)$ is free, then compute $R \in D^{q \times p}$ such that $\ker_D(\cdot P) = D^{1 \times q} R$ and go to 1 with $V = D^{1 \times p} R$.

3. If $W = D^{1 \times p} / \ker_D(\cdot P)$, then $W = D^{1 \times p} / (D^{1 \times q} R) = M$.

Quillen-Suslin theorem

- **Theorem:** Every **finitely generated projective module over the ring** $D = k[x_1, \dots, x_n]$, where k is a field, is **free**.
- **Corollary:** For every stably free D -module $M = D^{1 \times p} / (D^{1 \times q} R)$ defined by a **minimal presentation matrix** $R \in D^{q \times p}$, there exists $U \in GL_p(D)$, i.e., $\det U \in k \setminus \{0\}$, such that:

$$R U = (I_q \quad 0).$$

- **Corollary:** For every stably free D -module $M = D^{1 \times p} / (D^{1 \times q} R)$ defined by a **minimal presentation matrix** $R \in D^{q \times p}$, there exists $T \in D^{(p-q) \times p}$ such that:

$$\det \left(\begin{pmatrix} R \\ T \end{pmatrix} \right) \in k \setminus \{0\}.$$

- **Constructive proofs** of the Quillen-Suslin have been given in the literature (e.g., Logar-Sturmfels, Park, Lombardi-Yengui).

Particular case: principal ideal domain D

- Let D be a **principal ideal domain** D (e.g., $D = k[x]$, k a field).
- Computing a **Smith normal form of $R \in D^{q \times p}$** satisfying $RS = I_q$, we obtain $F \in GL_q(D)$ and $G \in GL_p(D)$ satisfying:

$$R = F \begin{pmatrix} I_q & 0 \end{pmatrix} G = F \begin{pmatrix} I_q & 0 \end{pmatrix} \begin{pmatrix} G_1 \\ G_2 \end{pmatrix} = F G_1 \Leftrightarrow G_1 = F^{-1} R.$$

$$\begin{pmatrix} F^{-1} R \\ G_2 \end{pmatrix} G^{-1} = I_p \Rightarrow \begin{pmatrix} F^{-1} & 0 \\ 0 & I_{p-q} \end{pmatrix} \begin{pmatrix} R \\ G_2 \end{pmatrix} G^{-1} = I_p,$$

$$\Rightarrow \begin{pmatrix} R \\ G_2 \end{pmatrix} G^{-1} \begin{pmatrix} F^{-1} & 0 \\ 0 & I_{p-q} \end{pmatrix} = I_p.$$

Then, the matrix $U = G^{-1} \begin{pmatrix} F^{-1} & 0 \\ 0 & I_{p-q} \end{pmatrix} \in GL_p(D)$ satisfies:

$$RU = \begin{pmatrix} I_q & 0 \end{pmatrix}.$$

Particular case: $R \in D^{(p-1) \times p}$

- Let D be a **commutative ring** and $R \in D^{(p-1) \times p}$ admitting a right-inverse $S \in D^{p \times (p-1)}$.
- Let us denote by m_i the $(p-1) \times (p-1)$ -minor of R obtained by removing the i^{th} column of R .
- The m_i 's satisfy a Bézout identity $\sum_{i=1}^p n_i m_i = 1$, with $n_i \in D$.
- Then, we can check that the matrix

$$V = \begin{pmatrix} & R & \\ (-1)^{p+1} n_1 & \dots & (-1)^{2p} n_p \end{pmatrix} \in D^{p \times p}$$

is such that $\det V = 1$ and its inverse $U = V^{-1} \in D^{p \times p}$ satisfies:

$$R U = (I_{p-1} \quad 0).$$

Reduction to the case of a single row

- Let $R \in D^{q \times p}$ a matrix admitting a **right-inverse** $S \in D^{p \times q}$.
- The computation of $U \in \text{GL}_p(D)$ satisfying $R U = (I_q \ 0)$ can be **reduced to the case of row vectors with entries in D** :

Let $U_1 \in \text{GL}_p(D)$ be such that $R_{1\bullet} U_1 = (1 \ 0 \ \dots \ 0)$

$$\Rightarrow R U_1 = \begin{pmatrix} 1 & 0 \\ C_1 & R_2 \end{pmatrix}.$$

$$R S = I_q \Leftrightarrow (R U_1)(U_1^{-1} S) = I_q$$

$$\Leftrightarrow \begin{pmatrix} 1 & 0 \\ C_1 & R_2 \end{pmatrix} \begin{pmatrix} W & X \\ Y & Z \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & I_{q-1} \end{pmatrix} \Leftrightarrow \begin{cases} W = 1, \\ X = 0, \\ C_1 + R_2 Y = 0, \\ R_2 Z = I_{q-1}, \end{cases}$$

$$\Rightarrow U_2 = \begin{pmatrix} 1 & 0 \\ Y & I_{p-1} \end{pmatrix} \in \text{GL}_p(D) : R(U_1 U_2) = \begin{pmatrix} 1 & 0 \\ 0 & R_2 \end{pmatrix} \cdots$$

Particular case: one invertible entry in R

- Let $R \in D^{1 \times p}$ a row vector admitting a **right-inverse** $S \in D^{p \times 1}$.
- If **one entry of R is invertible over D** , e.g., $R_1 \in U(D)$, then

$$(R_1 \ \dots \ R_p) \overbrace{\begin{pmatrix} R_1^{-1} & 0 \\ 0 & I_{p-1} \end{pmatrix}}^W = (1 \ R_2 \ \dots \ R_p),$$

and $\det W = R_1^{-1} \in D$. Denoting by $L = (R_2 \ \dots \ R_p)$, we get:

$$(1 \ L) \begin{pmatrix} 1 & -L \\ 0 & I_{p-1} \end{pmatrix} = (1 \ 0 \ \dots \ 0).$$

Then, the matrix $U = \begin{pmatrix} R_1^{-1} & 0 \\ 0 & I_{p-1} \end{pmatrix} \begin{pmatrix} 1 & -L \\ 0 & I_{p-1} \end{pmatrix} \in \mathrm{GL}_p(D)$ satisfies:

$$RU = (1 \ 0 \ \dots \ 0).$$

Particular case: 2 entries of R generate D

- Let D be a commutative ring.
- Let $R \in D^{1 \times p}$ a row vector admitting a right-inverse $S \in D^{p \times 1}$.
- We suppose that two entries of R , e.g., R_1 and R_2 generate D : there exist X_1 and $X_2 \in D$ such that $R_1 X_1 + R_2 X_2 = 1$.
- The matrix defined by

$$W = \begin{pmatrix} X_1 & -R_2 & 0 \\ X_2 & R_1 & 0 \\ 0 & 0 & I_{p-2} \end{pmatrix}$$

satisfies $\det W = 1$ and $RW = (1 \ 0 \ R_3 \ \dots \ R_p)$.

- Denoting by $L = (R_3 \ \dots \ R_p)$, we finally obtain:

$$(1 \ 0 \ L) \begin{pmatrix} 1 & 0 & -L \\ 0 & 1 & 0 \\ 0 & 0 & I_{p-2} \end{pmatrix} = (1 \ 0 \ \dots \ 0).$$

H. A. Park's example

- Let us consider $D = \mathbb{Q}[x, y]$ and $R = (1 - xy \quad x^2 \quad y^2)$.
- R admits the right-inverse $S = (xy + 1 \quad y^2 \quad 0)$ over D .
- In particular, the first two entries $R_1 = 1 - xy$ and $R_2 = x^2$ of R generate D : $R_1 X_1 + R_2 X_2 = 1$, where $X_1 = xy + 1$ and $X_2 = y^2$.
- Then, the unimodular matrices defined by

$$W = \begin{pmatrix} xy + 1 & -x^2 & 0 \\ y^2 & 1 - xy & 0 \\ 0 & 0 & 1 \end{pmatrix}, \quad Z = \begin{pmatrix} 1 & 0 & -y^2 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix},$$

satisfy $\det W = 1$, $RW = (1 \quad 0 \quad y^2)$ and $R(WZ) = (1 \quad 0 \quad 0)$.

$$WZ = \begin{pmatrix} xy + 1 & -x^2 & -(xy + 1)y^2 \\ y^2 & 1 - xy & -y^4 \\ 0 & 0 & 1 \end{pmatrix} \in \mathrm{GL}_3(D).$$

Particular case: one entry of R is 0

- Let D be a commutative ring.
- Let $R \in D^{1 \times p}$ a row vector admitting a right-inverse $S \in D^{p \times 1}$.
- We suppose that one entry of R (e.g., R_1) is 0, $\sum_{i=2}^p S_i R_i = 1$.
- The matrix defined by

$$W = \begin{pmatrix} 1 & & & & \\ (1 - R_1) S_2 & 1 & & & \\ \vdots & & \ddots & & \\ (1 - R_1) S_p & & & & 1 \end{pmatrix}$$

satisfies $\det W = 1$ and:

$$R W = (R_1 + (1 - R_1) \sum_{i=2}^p S_i R_i = 1 \quad R_2 \quad \dots \quad R_p).$$

The row vector $R W = (1 \quad R_2 \quad \dots \quad R_p)$ can then be reduced to $(1 \quad 0 \quad \dots \quad 0)$ by means of elementary operations.

Particular case: first condition on the right-inverse

- Let D be a commutative ring.
- Let $R \in D^{1 \times p}$ a row vector admitting a right-inverse $S \in D^{p \times 1}$.
- Let us suppose that one entry of S , e.g., S_1 is invertible.
- The matrix defined by

$$W = \begin{pmatrix} S_1 & & & & \\ S_2 & 1 & & & \\ \vdots & & \ddots & & \\ S_p & & & & 1 \end{pmatrix}$$

satisfies $\det W = S_1 \in U(D)$ and $RW = (1 \ R_2 \ \dots \ R_p)$.

The row vector $RW = (1 \ R_2 \ \dots \ R_p)$ can then be reduced to $(1 \ 0 \ \dots \ 0)$ by means of elementary operations.

Particular case: second condition on the right-inverse

- Let D be a commutative ring.
- Let $R \in D^{1 \times p}$ a row vector admitting a right-inverse $S \in D^{p \times 1}$.
- Let us suppose that two entries of S , e.g., S_1 and S_2 generate D : there exist X_1 and $X_2 \in D$ such that $X_1 S_1 + X_2 S_2 = 1$.
- The matrix defined by

$$W = \begin{pmatrix} S_1 & -X_2 & & & \\ S_2 & X_1 & & & \\ S_3 & & 1 & & \\ \vdots & & & \ddots & \\ S_p & & & & 1 \end{pmatrix}$$

satisfies $\det W = 1$ and $RW = (1 \quad * \quad R_3 \quad \dots \quad R_p)$, which can be reduced to $(1 \quad 0 \quad \dots \quad 0)$ by means of elementary operations.

Example: locally free modules

- Let us consider the $D = \mathbb{Q}[x_1, x_2]$ -module $M = D^{1 \times 3} / (D R)$:

$$R = (x_1^2 - x_2^2 - 1 \quad x_1^2 + x_2^2 - 1 \quad x_1 - x_2).$$

- The matrix $S = (-1 \quad 0 \quad x_1 + x_2)$ is a **right-inverse of R** , a fact proving that M is a **projective**, i.e., **free** D -module of rank 2.
- Checking that $\text{ext}_D^1(D/(D^{1 \times 3} R^T), D) = 0$, we obtain that

$$Q = \begin{pmatrix} x_1 - x_2 & -x_1 + x_2 & x_1^2 + x_2^2 - 1 \\ -x_1 + x_2 & -x_1 + x_2 & -x_1^2 + x_2^2 \\ 2x_2^2 & 2x_1^2 - 2 & 0 \end{pmatrix}$$

defines a **parametrization** of M , i.e., $M \cong D^{1 \times 3} Q \subseteq D^{1 \times 3}$.

- The parametrization Q is not injective because $\text{rank}_D(M) = 2$.

Example: locally free modules

- We have the following **3 minimal parametrizations of M** :

$$Q_1 = \begin{pmatrix} -x_1 + x_2 & x_1^2 + x_2^2 - 1 \\ -x_1 + x_2 & -x_1^2 + x_2^2 \\ 2x_1^2 - 2 & 0 \end{pmatrix}, \quad Q_2 = \begin{pmatrix} x_1 - x_2 & x_1^2 + x_2^2 - 1 \\ -x_1 + x_2 & -x_1^2 + x_2^2 \\ 2x_2^2 & 0 \end{pmatrix},$$
$$Q_3 = \begin{pmatrix} x_1 - x_2 & -x_1 + x_2 \\ -x_1 + x_2 & -x_1 + x_2 \\ 2x_2^2 & 2x_1^2 - 2 \end{pmatrix}.$$

None of them admits a left-inverse over D .

- The annihilators of the torsion D -modules $L_i = D^{1 \times 2} / (D^{1 \times 3} Q_i)$

$$\begin{cases} \text{ann}_D(L_1) = (x_1^2 - 1), \\ \text{ann}_D(L_2) = (x_2^2), \\ \text{ann}_D(L_3) = (x_1 - x_2). \end{cases}$$

satisfy the **Bézout identity** $-p_1 + p_2 + (x_1 + x_2)p_3 = 1$, where:

$$p_1 = x_1^2 - 1, \quad p_2 = x_2^2, \quad p_3 = x_1 - x_2.$$

Example: locally free modules

- Over the **localizations** $D_{p_i} = \{a/p_i^r \mid a \in D, r \in \mathbb{N}\}$ of D , the minimal parametrizations Q_i 's admit the following **left-inverses**:

$$T_1 = \frac{1}{2(x_1^2 - 1)} \begin{pmatrix} 0 & 0 & 1 \\ 1 & -1 & 0 \end{pmatrix}, \quad T_2 = \frac{1}{2x_2^2} \begin{pmatrix} 0 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix},$$

$$T_3 = -\frac{1}{2(x_1 - x_2)} \begin{pmatrix} -1 & 1 & 0 \\ 1 & 1 & 0 \end{pmatrix},$$

i.e., satisfy $T_i Q_i = I_2$, $i = 1, 2, 3$.

- Projective D -modules** are **locally free**.
- Computation of minimal parametrizations gives us **local bases**.

A constructive proof of the Quillen-Suslin theorem

- We shortly explain the idea of a **constructive proof** of the Quillen-Suslin theorem (Logar and Sturmfels).
- **Normalization step**: Let us consider $a \in k[y_1, \dots, y_n]$ and let us denote by $m = \deg a + 1$, where $\deg a$ is the total degree of a . Using the following **reversible transformation**

$$\begin{cases} x_n = y_n, \\ x_i = y_i - y_n^{m^{n-i}}, \end{cases} \Leftrightarrow \begin{cases} y_n = x_n, \\ y_i = x_i + x_n^{m^{n-i}}, \quad i = 1, \dots, n-1, \end{cases}$$

we obtain $a(y_1, \dots, y_n) = c b(x_1, \dots, x_n)$, where $0 \neq c \in k$ and b is a **monic polynomial** in x_n , i.e., the leading coefficient of $b \in E[x_n]$ is 1, where $E = k[x_1, \dots, x_{n-1}]$.

- If k is a **infinite field**, then we can obtain this result by means of a simpler transformation.

A constructive proof of the Quillen-Suslin theorem

- A ring A is called **local** if it contains **only one maximal ideal** \mathfrak{m} , namely, a proper ideal \mathfrak{m} of A which is not properly contained in any ideal of A other than A itself.
- **Computation of local bases (Horrocks's theorem)**: Let A be a commutative **local ring** and R a row vector admitting a right-inverse over $A[x]$. If **one of the components** R_i of R is **monic**, then there exists $U \in \text{GL}_p(A[x])$, satisfying:

$$R U = (1 \ 0 \ \dots \ 0).$$

- **Constructive proof of Horrocks's theorem can easily be obtained and implemented (QUILLEN SUSLIN).**

A constructive proof of the Quillen-Suslin theorem

- Main algorithm:
 - **Input:** $R \in D^{1 \times p}$ a row vector which admits a right-inverse over D and a monic component in the last variable x_n .
 - **Output:** A finite number of maximal ideals $\{\mathfrak{m}_i\}_{i \in I}$ of the ring $E = k[x_1, \dots, x_{n-1}]$ and unimodular matrices $\{H_i\}_{i \in I}$ over the ring $E_{\mathfrak{m}_i}[x_n]$, i.e., $H_i \in \text{GL}_p(E_{\mathfrak{m}_i}[x_n])$, satisfying

$$R H_i = (1, 0, \dots, 0),$$

and such that the ideal defined by the denominators of the matrices H_i 's, $i \in I$, generates E .

A constructive proof of the Quillen-Suslin theorem

- 1 Let \mathfrak{m}_1 be an arbitrary maximal ideal of the ring E . Using Horrocks' theorem, compute a unimodular matrix H_1 over $E_{\mathfrak{m}_1}[x_n]$ which satisfies that $R H_1 = (1 \ 0 \ \dots \ 0)$.
- 2 Let $d_1 \in E$ be the common denominator of all the entries of H_1 and J the ideal of E generated by d_1 . Set $i = 1$.
- 3 While $J \neq E$, do:
 - 1 For $i \leftarrow i + 1$, compute a maximal ideal \mathfrak{m}_i of E such that:

$$J \subset \mathfrak{m}_i.$$

- 2 Using Horrocks' theorem, compute a matrix H_i over the ring $E_{\mathfrak{m}_i}[x_n]$ such that $\det H_i$ is invertible in $E_{\mathfrak{m}_i}[x_n]$ and:
- $$R H_i = (1 \ 0 \ \dots \ 0).$$
- 3 Let d_i be the denominator of the matrix H_i and consider the ideal $J = (d_1, \dots, d_i)$.
 - 4 Return $\{\mathfrak{m}_i\}_{i \in I}$, $\{H_i\}_{i \in I}$ and $\{d_i\}_{i \in I}$.

A constructive proof of the Quillen-Suslin theorem

- Patching the local bases: Let $R \in D^{1 \times p}$ be a vector admitting a right-inverse over $D = k[x_1, \dots, x_n]$ and $U \in GL_p(E_{\mathfrak{m}}[x_n])$, where \mathfrak{m} is a maximal ideal of $E = k[x_1, \dots, x_{n-1}]$, which satisfies:

$$R U = (1 \ 0 \ \dots \ 0).$$

Let $d \in E \setminus \mathfrak{m}$ be a common denominator of the entries of U .

Then, the matrix defined by

$$\Delta(\bullet, x_n, z) = U(\bullet, x_n) U^{-1}(\bullet, x_n + z) \in GL_p(E_{\mathfrak{m}}[x_n, z])$$

is such that

$$\forall z \in D, \quad R(\bullet, x_n) \Delta(\bullet, x_n, z) = R(\bullet, x_n + z),$$

d^p is a common denominator of the entries of $\Delta(\bullet, x_n, z)$ and:

$$\Delta(\bullet, x_n, d^p z) \in GL_p(E[x_n, z]).$$

A constructive proof of the Quillen-Suslin theorem

- Let $\{m_i\}_{i \in I}$, $\{H_i\}_{i \in I}$ and $\{d_i\}_{i \in I}$ be the output of the main algorithm, where $I = \{1, \dots, m\}$. Let us define the matrices:

$$\Delta_i(\bullet, x_n, z) = H_i(\bullet, x_n) H_i^{-1}(\bullet, x_n + z), \quad i = 1, \dots, m.$$

Let $a_n \in k$. We have $(d_1, \dots, d_m) = E = k[x_1, \dots, x_{n-1}]$

$$\Rightarrow \exists c_i \in E, i = 1, \dots, m, \quad \sum_{i=1}^m c_i d_i^p = 1.$$

$$\begin{aligned} R(\bullet, x_n) \Delta_1(\bullet, x_n, (a_n - x_n) c_1 d_1^p) &= R(\bullet, x_n + (a_n - x_n) c_1 d_1^p), \\ R(\bullet, x_n + (a_n - x_n) c_1 d_1^p) \Delta_2(\bullet, x_n + (a_n - x_n) c_1 d_1^p, (a_n - x_n) c_2 d_2^p) \\ &= R\left(\bullet, x_n + (a_n - x_n) \left(\sum_{i=1}^2 c_i d_i^p\right)\right), \\ &\quad \dots \\ &\quad R\left(\bullet, x_n + (a_n - x_n) \left(\sum_{i=1}^{m-1} c_i d_i^p\right)\right) \\ \Delta_l\left(\bullet, x_n + (a_n - x_n) \left(\sum_{i=1}^{m-1} c_i d_i^p\right), (a_n - x_n) c_l d_l^p\right) &= R(\bullet, a_n). \end{aligned}$$

A constructive proof of the Quillen-Suslin theorem

- We finally obtain that the matrix

$$U(\bullet, x_n) = \Delta_1(\bullet, x_n, (a_n - x_n) c_1 d_1^p) \Delta_2(\bullet, x_n + (a_n - x_n) c_1 d_1^p, (a_n - x_n) c_2 d_2^p) \dots \Delta_l \left(\bullet, x_n + (a_n - x_n) \left(\sum_{i=1}^{l-1} c_i d_i^p \right), (a_n - x_n) c_l d_l^p \right) \in \text{GL}_p(D)$$

satisfies $R(\bullet, x_n) U(\bullet, x_n) = R(\bullet, a_n)$.

- **Theorem:** Let $D = k[x_1, \dots, x_n]$ be a commutative polynomial ring over a field k and $R \in D^{1 \times p}$ a row vector admitting a right-inverse over D . Then, for all $a_n \in k$, there exists $U \in \text{GL}_p(D)$ s.t.:

$$R(\bullet, x_n) U(\bullet, x_n) = R(\bullet, a_n).$$

- **Implementation** of the previous theorem was done in the package **QUILLEN****SUSLIN** (Fabiańska, Aachen University):

<http://wwwb.math.rwth-aachen.de/QuillenSuslin/>

Example

- We consider the $D = \mathbb{Q}[x_1, x_2]$ -module $M = D^{1 \times 3} / (D R)$, where:

$$R = (x_1 x_2^2 + 1 \quad 3x_2/2 + x_1 - 1 \quad 2x_1 x_2).$$

- Normalized entry** $3x_2/2 + x_1 - 1$ over $D = E[x_2]$ ($E = \mathbb{Q}[x_1]$).
- We consider the maximal ideal $\mathfrak{m}_1 = (x_1)$ of E . Using an effective version of **Horrocks' theorem**, we get that the matrix

$$\frac{1}{d_1} \begin{pmatrix} 4 & -2(3x_1 + 2x_2 - 2) & 4x_1(3x_1 - 2) \\ 2x_1(3x_1 - 2x_2 - 2) & 4(x_1 x_2^2 + 1) & -4x_1(3x_1^2 x_2 - 2x_1 x_2 + 2) \\ 0 & 0 & 9x_1^3 - 12x_1^2 + 4x_1 + 4 \end{pmatrix},$$

where $d_1 = 9x_1^3 - 12x_1^2 + 4x_1 + 4 \notin \mathfrak{m}_1$, is such that:

$$\begin{cases} \det H_1 = 4/d_1 \Rightarrow H_1 \in \text{GL}_3(E_{\mathfrak{m}_1}[x_2]), \\ R H_1 = (1 \quad 0 \quad 0). \end{cases}$$

Example

- We have $J = (d_1) \subsetneq E$. Then, we consider another maximal ideal \mathfrak{m}_2 such that $J \subseteq \mathfrak{m}_2$, e.g., $\mathfrak{m}_2 = (9x_1^3 - 12x_1^2 + 4x_1 + 4)$.
- Using an effective version of **Horrocks' theorem**, we obtain that

$$H_2 = \frac{1}{d_2} \begin{pmatrix} 0 & 0 & 4x_1(3x_1 - 2) \\ 8x_1 & -8x_1x_2 & -4x_1(3x_1^2x_2 - 2x_1x_2 + 2) \\ -4 & 2(3x_1 + 2x_2 - 2) & 9x_1^3 - 12x_1^2 + 4x_1 + 4 \end{pmatrix},$$

where $d_2 = 4x_1(3x_1 - 2) \notin \mathfrak{m}_2$, is such that:

$$\begin{cases} \det H_2 = -1/(x_1(3x_1 - 2)) \Rightarrow H_2 \in \text{GL}_3(E_{\mathfrak{m}_2}[x_2]), \\ RH_2 = (1 \ 0 \ 0). \end{cases}$$

- We have the Bézout identity

$$c_1 d_1 + c_2 d_2 = 1, \quad c_1 = 1/4, \quad c_2 = -(3x_1 - 2)/16,$$

i.e., $(d_1, d_2) = E$ and the **main algorithm stops**.

Example

- The matrix defined by

$$\Delta_1(x_1, x_2, -c_1 d_1 x_2) = H_1(x_1, x_2) H_1^{-1}(x_1, x_2 - c_1 d_1 x_2),$$

$$\begin{pmatrix} (9x_1^4/4 - 3x_1^3 + x_1^2)x_2^2 + (3x_1^2/2 - x_1)x_2 + 1 \\ -(18x_1^4 - 24x_1^3 + 8x_1^2)x_1x_2^3/8 + (27x_1^5 - 54x_1^4 + 36x_1^3 - 20x_1^2 + 8x_1)x_1x_2^2/8 - x_1x_2 \\ 0 \\ -x_2 & -2x_1x_2 \\ x_1x_2^2 + (-3x_1^2/2 + x_1)x_2 + 1 & 2x_1^2x_2^2 - x_1^2(3x_1 - 2)x_2 \\ 0 & 1 \end{pmatrix},$$

satisfies:

$$\begin{cases} \Delta_1(x_1, x_2, -c_1 d_1 x_2) \in GL_3(D), \\ R(x_1, x_2) \Delta_1(x_1, x_2, -c_1 d_1 x_2) = R(x_1, x_2 - c_1 d_1 x_2). \end{cases}$$

Example

- The matrix defined by

$$\Delta_2(x_1, x_2 - c_1 d_1 x_2, -c_2 d_2 x_2) = H_2(x_1, x_2 - c_1 d_1 x_2) H_2(x_2, 0)^{-1},$$

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & (3x_1^2/2 - x_1)x_2 + 1 & x_1^2(3x_1 - 2)x_2 \\ (9x_1^2 - 12x_1 + 4)x_1 x_2/8 & (-3x_1 + 2)x_2/4 & (-3x_1^2/2 + x_1)x_2 + 1 \end{pmatrix},$$

satisfies:

$$\begin{cases} \Delta_2(x_1, x_2 - c_1 d_1 x_2, -c_2 d_2 x_2) \in GL_3(D), \\ R(x_1, x_2 - c_1 d_1 x_2) \Delta_2(x_1, x_2 - c_1 d_1 x_2, -c_2 d_2 x_2) = R(x_1, 0). \end{cases}$$

$$U_1 = \Delta_1(x_1, x_2, -c_1 d_1 x_2) \Delta_2(x_1, x_2 - c_1 d_1 x_2, -c_2 d_2 x_2) \in GL_3(D),$$

$$R(x_1, x_2) U_1 = R(x_1, 0) = \begin{pmatrix} 1 & 3x_1/2 - 1 & 0 \end{pmatrix}.$$

Example

- We easily check that the matrix

$$U_2 = \begin{pmatrix} 1 & -3x_1/2 + 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \in \text{GL}_3(D)$$

satisfies $R(x_1, 0) U_2 = (1 \ 0 \ 0)$.

- Finally, if we define the matrix $U = U_1 U_2 \in \text{GL}_3(D)$, namely,

$$U = \begin{pmatrix} (3x_1^2/2 - x_1)x_2 + 1 & (-9x_1^3/4 + 3x_1^2 - x_1 - 1)x_2 - 3x_1/2 + 1 & \\ (-3x_1^3/2 + x_1^2)x_2^2 - x_1x_2 & (9x_1^4/4 - 3x_1^3 + x_1^2 + x_1)x_2^2 + (3x_1^2/2 - x_1)x_2 + 1 & \\ (9x_1^2 - 12x_1 + 4)x_1x_2/8 & (-27x_1^4/16 + 27x_1^3/8 - 9x_1^2/4 - x_1/4 + 1/2)x_2 & \\ & & -2x_1x_2 \\ & & 2x_1^2x_2^2 \\ & & (-3x_1^2/2 + x_1)x_2 + 1 \end{pmatrix},$$

we obtain $RU = (1 \ 0 \ 0)$!

Application: flat linear OD time-delay control system

- Let us consider the following OD time-delay linear system:

$$\begin{cases} \dot{y}_1(t) - y_1(t-h) + 2y_1(t) + 2y_2(t) - 2u(t-h) = 0, \\ \dot{y}_1(t) + \dot{y}_2(t) - \dot{u}(t-h) - u(t) = 0. \end{cases} \quad (\star)$$

- We consider $D = \mathbb{Q}(a) \left[\partial; \text{id}, \frac{d}{dt} \right] [\delta; \sigma, 0]$ and the two matrices:

$$R = \begin{pmatrix} \partial - \delta + 2 & 2 & -2\delta \\ \partial & \partial & -\partial\delta - 1 \end{pmatrix}, \quad S = \begin{pmatrix} 0 & 0 \\ \frac{1}{2}(\partial\delta + 1) & -\delta \\ \frac{1}{2}\partial & -1 \end{pmatrix}.$$

- We can easily check that $RS = I_2$, which proves that the D -module $M = D^{1 \times 3} / (D^{1 \times 2} R)$ is **free** (Quillen-Suslin theorem), and thus, (\star) admits an **injective parametrisation**.

Application: flat linear OD time-delay control system

- We have the following **system equivalence**

$$\begin{cases} \dot{y}_1(t) - y_1(t-h) + 2y_1(t) + 2y_2(t) - 2u(t-h) = 0, \\ \dot{y}_1(t) + \dot{y}_2(t) - \dot{u}(t-h) - u(t) = 0. \end{cases}$$

$$\Leftrightarrow \begin{cases} \dot{z}_1(t) + 2z_1(t) + 2z_2(t) = 0, \\ \dot{z}_1(t) + \dot{z}_2(t) - v(t) = 0, \end{cases}$$

defined by the following **reversible transformations**:

$$\begin{cases} y_1(t) = z_1(t), \\ y_2(t) = \frac{1}{2}(\dot{z}_1(t-2h) + z_1(t-h)) + z_2(t) + v(t-h), \\ u(t) = \frac{1}{2}\dot{z}_1(t-h) + v(t). \end{cases}$$

$$\Leftrightarrow \begin{cases} z_1(t) = y_1(t), \\ z_2(t) = -\frac{1}{2}y_1(t-h) + y_2(t) - u(t-h), \\ v(t) = -\frac{1}{2}\dot{y}_1(t-h) + u(t), \end{cases}$$

Application: flat linear OD time-delay control system

- Moreover, we have the following **system equivalence**

$$\begin{cases} \dot{z}_1(t) + 2z_1(t) + 2z_2(t) = 0, \\ \dot{z}_1(t) + \dot{z}_2(t) - v(t) = 0, \end{cases} \Leftrightarrow \begin{cases} 2x_1(t) + 2x_2(t) = 0, \\ -w(t) = 0, \end{cases}$$

defined by the following **reversible transformations**:

$$\begin{cases} z_1(t) = x_1(t), \\ z_2(t) = x_2(t) - \frac{1}{2}\dot{x}_1(t), \\ v(t) = w(t) - \frac{1}{2}\ddot{x}_1(t) + \dot{x}_1(t) + \dot{x}_2(t), \end{cases} \Leftrightarrow \begin{cases} x_1(t) = z_1(t), \\ x_2(t) = z_2(t) + \frac{1}{2}\dot{z}_1(t), \\ w(t) = v(t) + \dot{z}_1(t) + \dot{z}_2(t). \end{cases}$$

- We finally obtain the following **injective parametrisation**:

$$\begin{cases} y_1(t) = x_1(t), \\ y_2(t) = \frac{1}{2}(-\ddot{x}_1(t-h) + \dot{x}_1(t-2h) - \dot{x}_1(t) + x_1(t-h) - 2x_1(t)), \\ u(t) = \frac{1}{2}(\dot{x}_1(t-h) - \ddot{x}_1(t)). \end{cases}$$

Application: δ -flat linear OD time-delay systems

- Flexible rod with a mass:

$$\left\{ \begin{array}{l} \sigma^2 \frac{\partial^2 q(\tau, x)}{\partial \tau^2} - \frac{\partial^2 q(\tau, x)}{\partial x^2} = 0, \\ \frac{\partial q}{\partial x}(\tau, 0) = -u(\tau), \\ \frac{\partial q}{\partial x}(\tau, L) = -J \frac{\partial^2 q}{\partial \tau^2}(\tau, L), \\ y(\tau) = q(\tau, L). \end{array} \right. \quad (*)$$

- $q(\tau, x) = \phi(\tau + \sigma x) + \psi(\tau - \sigma x)$, $t = (\sigma/J)\tau$, $v = (2J/\sigma^2)u$,
 $(*) \Rightarrow \ddot{y}(t+1) + \ddot{y}(t-1) + \dot{y}(t+1) - \dot{y}(t-1) = v(t)$

$$\Leftrightarrow \left\{ \begin{array}{l} y(t) = \xi(t-1), \\ v(t) = \ddot{\xi}(t) + \ddot{\xi}(t-2) + \dot{\xi}(t) - \dot{\xi}(t-2). \end{array} \right.$$

- If y_r is a **desired trajectoire**, then $\xi_r(t) = y_r(t+1)$ and:

$$v_r(t) = \ddot{y}_r(t+1) + \ddot{y}_r(t-1) + \dot{y}_r(t+1) - \dot{y}_r(t-1).$$

Application: π -flat linear OD time-delay systems

- Wind tunnel model (Manitius, IEEE TAC 84):

$$\begin{cases} \dot{x}_1(t) + a x_1(t) - k a x_2(t - h) = 0, \\ \dot{x}_2(t) - x_3(t) = 0, \\ \dot{x}_3(t) + \omega^2 x_2(t) + 2 \zeta \omega x_3(t) - \omega^2 u(t) = 0, \end{cases}$$

$$\Leftrightarrow \begin{cases} x_1(t) = \omega^2 a k \xi(t - h), \\ x_2(t) = \omega^2 \dot{\xi}(t) + \omega^2 a \xi(t), & \xi(t) = x_1(t + h) / (\omega^2 a), \\ x_3(t) = \omega^2 \dot{\xi}(t) + \omega^2 a \dot{\xi}(t), \\ u(t) = \xi^{(3)}(t) + (2 \zeta \omega + a) \ddot{\xi}(t) + (\omega^2 + 2 a \zeta \omega) \dot{\xi}(t) + a \omega^2 \xi(t). \end{cases}$$

- Simple network model (Fliess-Mounier, IFAC TDS98):

$$\begin{cases} \dot{x}_1(t) + u_1(t) - u_2(t - h_1) = 0, \\ \dot{x}_2(t) - u_1(t - h_2) = 0, \end{cases} \Leftrightarrow \begin{cases} x_1(t) = \xi_1(t - h_1) - \xi_2(t), \\ x_2(t) = \xi_2(t - h_2), \\ u_1(t) = \dot{\xi}_2(t), \\ u_2(t) = \dot{\xi}_1(t). \end{cases}$$

$$\xi_1(t) = x_1(t + h_1) + x_2(t + h_1 + h_2), \quad \xi_2(t) = x_2(t + h_2).$$

Conclusion

- We have studied **stably free** and **free modules**.
- We have briefly explained the **Quillen-Suslin theorem**.
- ① **Constructive computation of bases** of free D -modules can be obtained by means of the package `QUILLENUSLIN`:

`http://wwwb.math.rwth-aachen.de/QuillenSuslin/`

- ② **More applications in mathematical systems theory**:
constructive solutions of the Lin-Bose's conjectures, effective computation of (weakly) coprime factorizations of rational transfer matrices, reduction and decomposition problems. . .
A. Fabiańska, A. Quadrat, "Applications of the Quillen-Suslin theorem to multidimensional systems theory", in *Gröbner Bases in Control Theory and Signal Processing*, H. Park and G. Regensburger, Radon Series on Computation and Applied Mathematics 3, de Gruyter publisher, 2007, 23-106.