# Symbolic computation for operators with matrix coefficients 

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## Introduction

Our goal is to provide an algebraic framework to compute with operators having matrix coefficients of generic size. For example, consider a system of homogeneous LODE

$$
Y^{\prime}-A Y=0 \quad \text { where } \quad A \in C^{\infty}(\mathbb{R})^{n \times n}
$$

Let $L=\partial-A$ where $\partial$ denotes the derivation operator and $A$ denotes the multiplication operator $F \mapsto A F$ induced by the matrix $A$. If $\Phi$ is a fundamental matrix of $L y=0$, then by the Leibniz rule, independently of the size $n$, the following identity holds

$$
L \circ \Phi=\Phi \circ \partial+\partial \Phi-A \Phi=\Phi \circ \partial .
$$

In what follows, let $(R, \partial)$ be a differential ring with ring of constants

$$
K=\{c \in R \mid \partial c=0\} .
$$

Considering the coefficient ring $R$ and its ring of constants $K$, we have two cases.
Scalars: $R$ and $K$ are commutative and hence $R$ and $R\langle\partial\rangle$ are $K$-algebras.
Matrices: $R$ and $K$ are noncommutative and $r c \neq c r$, so $R$ and $R\langle\partial\rangle$ are only $K$-rings.

## Tensor ring

We use the tensor ring for modelling additive operators and interpret $\otimes$ as composition of operators. The tensor product of $K$-bimodules $M \otimes_{K} N$ over an arbitrary ring $K$ satisfies $\left(m+m^{\prime}\right) \otimes n=m \otimes n+m^{\prime} \otimes n, \quad m \otimes\left(n+n^{\prime}\right)=m \otimes n+m \otimes n^{\prime}, \quad m k \otimes n=m \otimes k n$. Note that in general $k m \otimes n \neq m \otimes k n$. The $n$-fold tensor product of a $K$-bimodule $M$ is denoted by $M^{\otimes n}=M \otimes_{K} \cdots \otimes_{K} M$ such that $M^{\otimes 0}=K \varepsilon$, where $\varepsilon$ denotes the empty tensor. Then the $K$-tensor ring over the $K$-bimodule $M$ is defined as a $K$-bimodule by

$$
K\langle M\rangle=\bigoplus^{\infty} M^{\otimes n}
$$

Every $t \in K\langle M\rangle$ can be written in the form

$$
t=k \varepsilon+\sum_{i=1}^{p} m_{i, 1} \otimes \cdots \otimes m_{i, n_{i}}
$$

The $K$-tensor ring over the free $K$-bimodule on a set $X$ is generated as an additive group by $\left\{k_{1} x_{1} \otimes k_{2} x_{2} \otimes \cdots \otimes k_{n} x_{n} k_{n+1} \mid n \in \mathbb{N}_{0}, x_{i} \in X, k_{j} \in K\right\}$
and elements of the tensor ring do not have a unique representation. In contrast, if $K$ is commutative and $M$ is a free left $K$-module over some set $X$, elements of the $K$-tensor algebra $K\langle M\rangle$ have a unique representation as $K$-linear combinations of products $x_{1} \otimes x_{2} \otimes \cdots \otimes x_{n}$.

## Tensor reduction systems

Identities of operators which cannot be covered by the tensor ring are modelled by tensor reduction rules $r_{W}=(W, h)$ where $h: M_{W} \rightarrow K\langle M\rangle$ is a $K$-bimodule homomorphism.

## Example: Differential operators

Let $(R, \partial)$ be a differential ring with ring of constants $K$. We consider the $K$-bimodules

$$
M=R \oplus K \partial, \quad M_{\mathrm{K}}=K, \quad M_{\mathrm{R}}=R, \quad M_{\mathrm{D}}=K \partial
$$

and reduction rules acting on $M_{\mathrm{K}}, M_{\mathrm{R}} \otimes M_{\mathrm{R}}, M_{\mathrm{D}} \otimes M_{\mathrm{R}}$ defined by

$$
(\mathrm{K}, 1 \mapsto \varepsilon), \quad(\mathrm{RR}, f \otimes g \mapsto f g), \quad(\mathrm{DR}, \partial \otimes f \mapsto f \otimes \partial+\partial f)
$$

These rules induce the two-sided reduction ideal

$$
J=(1-\varepsilon, f \otimes g-f g, \partial \otimes f-f \otimes \partial-\partial f \mid f, g \in R),
$$

which is used to define the $K$-ring of differential operators as $R\langle\partial\rangle=K\langle M\rangle / J$.
More formally, the framework for tensor reduction systems is the following. Let $\left(M_{z}\right)_{z \in Z}$ be a family of a $K$-subbimodules of $M$ and let $X \subseteq Z$ be sets where $M=\bigoplus_{x \in X} M_{x}=\sum_{z \in Z} M_{z}$. To a word $W=w_{1} \ldots w_{n} \in\langle Z\rangle$ we associate the $K$-bimodule $M_{W}:=M_{w_{1}} \otimes \cdots \otimes M_{w_{n}}$. So

$$
K\langle M\rangle=\bigoplus_{n=0}^{\infty} M^{\otimes n}=\bigoplus_{W \in\langle X\rangle} M_{W}=\sum_{W \in\langle Z\rangle} M_{W}
$$

## Theorem (confluence criterion, [2, 4]):

Given a tensor reduction system, every tensor $t \in K\langle M\rangle$ has a unique normal form $t \downarrow$ iff all ambiguities are resolvable. In that case, the tensor ring factored by the reduction ideal is isomorphic to the ring of irreducible tensors with multiplication $s \cdot t:=(s \otimes t) \downarrow$.

## Integro-differential operators with linear substitutions

Recall the fundamental theorem of calculus and the definition of linear substitution operators: $\frac{d}{d x} \int_{a}^{x} f(s) d s=f(x), \quad f(a)=f(x)-\int_{a}^{x} f^{\prime}(s) d s, \quad \sigma_{a, b}(f(x))=f(a x-b)$.

Based on these identities we have the following general definition.

## Definition:

An integro-differential ring with linear substitutions $\left(R, \partial, \int, S\right)$ is a differential ring $(R, \partial)$ with the ring of constants $K$ where $\int$ is a $K$-bimodule homomorphism of $R$ such that $\partial \int f=f$, the evaluation $\mathrm{E} f:=f-\int \partial f$ is multiplicative, i.e. $\mathrm{E} f g=(\mathrm{E} f) \mathrm{E} g$, and $S$ is a group of multiplicative $K$-bimodule homomorphisms $\sigma_{a, b}: R \rightarrow R$ on $R$ fixing the constants $K$ and satisfying $\partial \sigma_{a, b} f=a \sigma_{a, b} \partial f$.

Let $\left(R, \partial, \int, S\right)$ be an integro-differential ring with linear substitutions and

$$
M=R \oplus K \partial \oplus K \int \oplus K \mathrm{E} \oplus K S
$$

Let $M_{\mathrm{K}}, M_{\mathrm{R}}, M_{\mathrm{D}}$ as before, $M_{\mathrm{I}}=K \int, M_{\mathrm{E}}=K \mathrm{E}, M_{\mathrm{G}}=K S$, and $M_{\mathrm{N}}=K \sigma_{1,0}$. Then we call $R\left\langle\partial, \int, \mathrm{E}, S\right\rangle:=K\langle M\rangle / J$ the ring of integro-differential operators with linear substitutions, where $J$ is the two-sided reduction ideal induced by the reduction system:
$\begin{array}{ccc}\text { K } & & 1 \\ \text { RR } & f \otimes g & \mapsto f g\end{array}$
$\mathrm{RR} \quad f \otimes g \mapsto f g$
$\mathrm{DR} \partial \otimes f \mapsto f \otimes \partial+\partial f$
ER $\mathrm{E} \otimes f \mapsto(\mathrm{E} f) \mathrm{E}$
$\mathrm{EE} \mathrm{E} \otimes \mathrm{E} \mapsto \mathrm{E}$
$\mathrm{EI} \mathrm{E} \otimes \int \mapsto 0$
IRE $\quad \int \otimes f \otimes \mathrm{E} \mapsto \int f \otimes \mathrm{E}$
DE $\partial \otimes E \mapsto 0$
DI $\partial \otimes \int \mapsto \varepsilon$
IE $\int \otimes \mathrm{E} \mapsto \int 1 \otimes \mathrm{E} \quad$ IG $\quad \int \otimes \sigma_{a, b} \mapsto a^{-1}(\varepsilon-\mathrm{E}) \otimes \sigma_{a, b} \otimes \int$
IRD $\int \otimes f \otimes \partial \mapsto f-\int \otimes \partial f-(\mathrm{E} f) \mathrm{E}$
IRI $\quad \int \otimes f \otimes \int \mapsto \int f \otimes \int-\int \otimes \int f$
$\begin{aligned} & \sigma_{1,0} \mapsto \varepsilon \\ & \sigma_{0,1} \\ & \sigma_{b,}\end{aligned}$
$\sigma_{a, b} \otimes \sigma_{c, d} \mapsto \sigma_{a c, b c+d}$
$\sigma_{a, b} \otimes f \mapsto \sigma_{a, b} f \otimes \sigma_{a, b}$
$\begin{array}{ll}\mathrm{GR} & \\ \text { GE,b} & \sigma_{a, b} \otimes \mathrm{E} \mapsto \mathrm{E}_{a, b} \\ \text { DG } & \partial \otimes \sigma_{a, b} \mapsto a \sigma_{a, b} \otimes \partial\end{array}$
ID $\int \otimes \partial \mapsto \varepsilon-\mathrm{E} \quad$ IRG $\int \otimes f \otimes \sigma_{a, b} \mapsto a^{-1}(\varepsilon-\mathrm{E}) \otimes \sigma_{a, b} \otimes \int \otimes \sigma_{a, b}^{-1} f$
II $\int \otimes \int \mapsto \int 1 \otimes \int-\int \otimes \int 1$

Checking confluence, reducing S-polynomials, and computing with operators are supported by our Mathematica package TenReS [3]. Using the above reduction system, every $t \in K\langle M\rangle$ has a unique normal form given by a sum of pure tensors

$$
f \otimes \mathrm{E} \otimes \sigma_{a, b} \otimes \partial^{\otimes j} \quad \text { and } \quad f \otimes \mathrm{E} \otimes \sigma_{a, b} \otimes \int \otimes g
$$

where $j \in \mathbb{N}_{0}$, each of $f, g \in \int R$ and $\sigma_{a, b} \in S \backslash\left\{\sigma_{1,0}\right\}$ may be absent.

## Example: Variation of constants

Identities in $R\left\langle\partial, \int, \mathrm{E}, S\right\rangle$ can be proven using the reduction system. Let $R=C^{\infty}(\mathbb{R})^{n \times n}$ $L=\partial-A \in R\left\langle\partial, \int, \mathrm{E}, S\right\rangle$, and let $\Phi \in R$ be an invertible solution of $L y=0$. Then $H:=\Phi \otimes \int \otimes \Phi^{-1}$ is a right inverse of $L$ since independent of the size $n$ we have
$L \otimes H=(\partial-A) \otimes \Phi \otimes \int \otimes \Phi^{-1} \rightarrow_{r_{\mathrm{DR}}} \Phi \otimes \partial \otimes \int \otimes \Phi^{-1} \rightarrow_{r_{\mathrm{DI}}} \Phi \otimes \Phi^{-1} \rightarrow_{r_{\mathrm{RR}}} \Phi \Phi^{-1} \rightarrow_{r_{\mathrm{K}}} \varepsilon$.

## Artstein's reduction of DTD control systems

Aim [5]: Algebraically find Artstein's reduction [1] of a differential time-delay control system

$$
x^{\prime}(t)=A_{1}(t) x(t)+A_{2}(t) u(t)+A_{3}(t) u(t-h)
$$

to a differential control system

$$
z^{\prime}(t)=B_{1}(t) z(t)+B_{2}(t) v(t) .
$$

We use our tensor setting to find an invertible transformation $\binom{z}{v}=\left(\begin{array}{cc}P_{11} & P_{12} \\ 0 & P_{22}\end{array}\right)\binom{x}{u}$, where $P_{11}$ and $P_{22}$ are arbitrary invertible multiplication operators, and make an ansatz for

$$
P_{12}=C_{0} \otimes \sigma_{1, h} \otimes \int \otimes C_{1}+C_{2} \otimes \int \otimes C_{3}+C_{4} \otimes \sigma_{1, h}+C_{5} \in R\left\langle\partial, \int, \mathrm{E}, S\right\rangle
$$

We consider the ring $R$ generated by the symbols $A_{i}, B_{i}, C_{i}, P_{i i}$ and we can check that only expressions occur in computations that are valid as operators. By plugging the ansatz into

$$
\left(\partial-B_{1}\right) \otimes P_{12}-B_{2} \otimes P_{22}=P_{11} \otimes\left(-A_{2}-A_{3} \otimes \sigma_{1, h}\right)
$$

and reducing both sides to normal form using TenReS, we obtain conditions for the multiplication operators $C_{i}$. Solving these conditions with partial support by our package, we obtain the following transformation

$$
v=P_{22} u, \quad z=P_{11} x+\left(\Phi \otimes\left(\varepsilon-\sigma_{1, h}\right) \otimes \int \otimes\left(\sigma_{1,-h} \Phi^{-1}\right) \sigma_{1,-h} A_{3}\right) u
$$

where $P_{11}^{-1} \Phi$ is a fundamental matrix of $y^{\prime}(t)=A_{1}(t) y(t)$.

## References

[1] Zvi Artstein. Linear systems with delayed controls: A reduction. IEEE Trans. Autom. Control 27, pp. 869-879, 1982 [2] George M. Bergman. The Diamond Lemma for Ring Theory. Adv. Math. 29, pp. 178-218, 1978
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[5] Alban Quadrat. A constructive algebraic analysis approach to Artstein's reduction of linear time-delay systems. Proceedings of 12th IFAC Workshop on Time Delay Systems, IFAC-PapersOnline 48 (12), pp. 209-214, 2015.

