# OreModules: A Symbolic Package for the Study of Multidimensional Linear Systems 

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## 1 Introduction

In the seventies, the study of transfer matrices of time-invariant linear systems of ordinary differential equations (ODEs) led to the development of the polynomial approach [20, 22, 44]. In particular, the univariate polynomial matrices play a central role in this approach (e.g., Hermite, Smith and Popov forms, invariant factors, primeness, Bézout/Diophantine equations).

In the middle of the seventies, while generalizing linear systems defined by ODEs to differential time-delay systems, ODEs with parameters, 2-D and 3-D filters..., one had to face the case of systems described by means of matrices with entries in multivariate commutative polynomial rings. All these new systems were called 2-D or 3-D linear systems and, more generally, $n$ - $D$ systems or multidimensional linear systems with constant coefficients [4, 16. It was quickly realized that no canonical forms such as Hermite, Smith and Popov forms existed for polynomial matrices with two and three variables (i.e., with entries in $k\left[x_{1}, x_{2}, x_{3}\right]$, where $k$ is a field such as $\left.\mathbb{Q}, \mathbb{R}, \mathbb{C}\right)$. Moreover, more than only one type of primeness was needed in order to classify $n$-D systems (e.g., factor/minor/zero primeness [48, 49]). Hence, it is not very surprising that, in the eighties, Gröbner bases were introduced in the study of multidimensional linear systems with constant coefficients [4, 16]. A Gröbner basis defines normal forms for polynomials with respect to a certain monomial ordering of the variables $x_{i}$ [2, 17, 23]. Given a Gröbner basis, there is a simple algorithm to effectively compute these normal forms. In many ways, the computation of these normal forms can be seen as an extension of the Gaussian elimination algorithm to commutative polynomial rings [2, 17].

In a pioneering work, R. E. Kalman developed a module-theoretic approach to time-invariant ordinary differential linear systems [21]. In his PhD thesis under the supervision of R. E. Kalman, Y. Rouchaleau considered Kalman-type

[^0]systems where the entries of $(A, B, C, D)$ belong to a commutative ring. In particular, he studied their structural properties using module theory. Such systems are nowadays called systems over rings and they have been considerably studied in the literature 43] since. An extension of the geometric approach [46] to linear systems over rings has also been recently developed [1, 12, 18, 19]. Using effective algebra methods (Gröbner bases, characteristic sets), the computational aspects of the systems over rings (e.g., differential time-delay systems) were firstly studied by L. Habets in 18, 19.

In the nineties, U . Oberst developed a general module-theoretic approach to multidimensional linear systems with constant coefficients [28]. Using B. Malgrange's approach [24], in which a finitely presented $D$-module $M$ is associated with a linear system of equations over a polynomial ring $D$, he showed how some structural properties of the system corresponded to algebraic properties of the $D$-module $M$. He then was able to develop a complete duality between his module-theoretic approach and the behavioural approach developed by J. C. Willems 30. Based on U. Oberst's ideas, the behavioural approach to multidimensional linear systems has been successfully developed in the recent years. See [30, 29, 36, 47, 49] and the references therein.

Within a similar module-theoretic approach, the concepts of flatness and $\pi$ freeness were introduced in [15, 26] for differential time-delay linear systems with constant coefficients. As it is shown in [26, 27] on different concrete examples, the detection of such structural properties is important for the study of the motion planning problem. In the behavioural approach, the concept of flatness corresponds to the existence of an observable image representation for the multidimensional system 32 .

In the same years as [28, J.-F. Pommaret studied underdetermined systems of partial differential equations (PDEs) coming from mathematical physics and differential geometry (e.g., elasticity, electromagnetism, hydrodynamics, general relativity). See also [3]. In particular, he showed how his mathematical approach was a generalization of U. Oberst's module-theoretic approach for multidimensional (linear) systems with varying coefficients. See 31 for more details and references. In particular, the problem of checking whether or not a multidimensional linear system described by PDEs with varying coefficients could be formally parametrized was solved within the theory of differential operators. Moreover, the work of M. Fliess on linear systems defined by ODEs with variable coefficients also illustrated the need to pass from the commutative polynomial viewpoint to the non-commutative one [14].

Based on B. Malgrange's approach [24, algebraic analysis has been developed in mathematics in order to study general linear systems of PDEs with variable coefficients using module theory, algebraic geometry, homological algebra and functional analysis. Algebraic analysis has recently been introduced in control theory in [38] in order to study multidimensional linear systems defined by PDEs with varying coefficients. In particular, using the formal theories of PDEs (Spencer's, Riquier-Janet's theories), it was shown in 31, 32, 33, 34, 38, how some structural properties of systems could be checked by means of constructive algorithms.

Finally, using the homological algebra approach developed in [38], we have recently shown in [9, 11] how the previous results could be generalized to some classes of multidimensional linear systems with varying coefficients encountered in the literature (e.g., ODEs, PDEs, differential time-delay systems, multidimensional discrete systems, partial differential delay systems). In order to do that, the concept of multidimensional linear systems over Ore algebras was introduced in [9, 11]. An Ore algebra is a ring of non-commutative polynomials in functional operators with polynomial or rational coefficients [5, 6, 7]. Characterizations of algebraic structural properties such as, for instance, controllability, parametrizability and flatness were obtained.

The recent progress of Gröbner bases over Ore algebras (i.e., over some classes of non-commutative polynomial rings) [5, 6, 7, 23] allows us to effectively test the algebraic properties of general multidimensional linear systems (e.g., controllability, observability, parametrizability, flatness, $\pi$-freeness) and compute different types of parametrizations and to propose feedback laws (motion planning, tracking, Bézout equations, optimal control).

In this paper, we shall develop the following methodology for the study of multidimensional linear systems over Ore algebras (see also [11]):

1. A linear system is defined by means of a $(q \times p)$-matrix $R$ with entries in an Ore algebra $D$, i.e., it corresponds to a system of linear equations $R z=0$, where $z$ is composed of the system variables (see Section (21).
2. We associate the finitely presented left $D$-module $M=D^{1 \times p} /\left(D^{1 \times q} R\right)$ with the system $R z=0$.
3. We develop a dictionary between the structural properties of the system and the properties of the left $D$-module $M$. Using module theory, we can then classify the properties of the left $D$-module $M$ (see Section 3).
4. Homological algebra permits to check these properties of the left $D$-module $M$ using extension and torsion functors (see Section (4).
5. Gröbner bases over Ore algebras allow to develop effective algorithms which check the properties of the left $D$-module $M$, and thus, of the system $R z=0$ (see Section 5).
6. Implementations of these algorithms in the package OreModules for the computer algebra system Maple (see Section 6).

The purpose of this paper is to give an introduction to the package OrEMODULES [8 for Maple which offers symbolic methods to investigate the structural properties of multidimensional linear systems over Ore algebras. The advantage of describing these properties in the language of homological algebra carries over to the implementation of OreModules: up to the choice of the domain of operators which occur in a given system, all algorithms are stated and implemented in sufficient generality such that ODEs, PDEs, differential time-delay systems, discrete systems with constant, polynomial or rational coefficients... are covered at the same time.

This paper is an extension of the congress paper [10].

## 2 Multidimensional Linear Systems over Ore Algebras

The mathematical framework of this paper is built on the concept of Ore algebras [5, 6, 7]. Ore algebras are non-commutative polynomial rings that represent linear functional operators in a natural way.

We recall that a ring with a unit 1 is a domain if the product of non-zero elements is non-zero. In what follows, we shall denote by $A$ domain which has a $k$-algebra structure, where $k$ is a field.
Definition 1. 1. 25] A skew polynomial ring $A[\partial ; \sigma, \delta]$ is a non-commutative ring consisting of all polynomials in $\partial$ with coefficients in $A$ obeying the commutation rule

$$
\begin{equation*}
\forall a \in A, \quad \partial a=\sigma(a) \partial+\delta(a), \tag{1}
\end{equation*}
$$

where $\sigma$ is a $k$-algebra endomorphism of $A$, namely, $\sigma: A \longrightarrow A$ satisfies

$$
\sigma(1)=1, \quad \forall a, b \in A, \quad \sigma(a+b)=\sigma(a)+\sigma(b), \quad \sigma(a b)=\sigma(a) \sigma(b),
$$

and $\delta$ is a $\sigma$-derivation of $A$, namely, $\delta: A \longrightarrow A$ satisfies:

$$
\forall a, b \in A, \quad \delta(a+b)=\delta(a)+\delta(b), \quad \delta(a b)=\sigma(a) \delta(b)+\delta(a) b
$$

2. [5, 7] Let $A=k\left[x_{1}, \ldots, x_{n}\right]$ be a commutative polynomial ring over a field $k$ (if $n=0$ then $A=k$ ). The skew polynomial ring

$$
D=A\left[\partial_{1} ; \sigma_{1}, \delta_{1}\right] \ldots\left[\partial_{m} ; \sigma_{m}, \delta_{m}\right]
$$

is called Ore algebra if the $\sigma_{i}$ 's and $\delta_{j}$ 's commute for $1 \leq i, j \leq m$ and satisfy:

$$
\sigma_{i}\left(\partial_{j}\right)=\partial_{j}, \quad \delta_{i}\left(\partial_{j}\right)=0, \quad j<i
$$

Example 1. In order to model an ordinary differential linear system with polynomial coefficients, we use the Weyl algebra $A_{1}(k)=k[t][\partial ; \sigma, \delta]$ which is a non-commutative $k$-algebra generated by $t$ and $\partial$. Elements of $A_{1}(k)$ are non-commutative polynomials in $t$ and $\partial$ with coefficients in the field $k$ (e.g., $k=\mathbb{Q}, \mathbb{R}, \mathbb{C}$ ) satisfying the following commutation rule:

$$
\forall a \in k[t], \quad \partial(a \cdot)=a \partial \cdot+\frac{d a}{d t} .
$$

Therefore, regarding Definition 1 we have $\sigma=\mathrm{id}_{k[t]}$ and $\delta=\frac{d}{d t}$.
More generally, for the study of partial differential linear systems, we shall use the Weyl algebra $A_{n}(k)=k\left[x_{1}, \ldots, x_{n}\right]\left[\partial_{1} ; \sigma_{1}, \delta_{1}\right] \ldots\left[\partial_{n} ; \sigma_{n}, \delta_{n}\right]$, where $\sigma_{i}$ and $\delta_{i}$ are the maps on $k\left[x_{1}, \ldots, x_{n}\right]$ defined by

$$
\sigma_{i}=\operatorname{id}_{k\left[x_{1}, \ldots, x_{n}\right]}, \quad \delta_{i}=\frac{\partial}{\partial x_{i}}, \quad i=1, \ldots, n
$$

and every other commutation rule is prescribed by Definition 1. We have:

$$
\partial_{i} x_{j}=x_{j} \partial_{i}+\delta_{i j}, \quad 1 \leq i, j \leq n, \quad \text { where } \quad \delta_{i j}=1 \text { if } i=j \text { and } 0 \text { else. }
$$

Example 2. The algebra of shift operators with polynomial coefficients is another special case of an Ore algebra. For $h$ in the field $k$ (e.g., $k=\mathbb{Q}, \mathbb{R}$ ), we define $S_{h}(k)=k[t]\left[\delta_{h} ; \sigma_{h}, \delta\right]$ by:

$$
\forall a \in k[t], \quad \sigma_{h}(a)(t)=a(t-h), \quad \delta(a)=0 .
$$

The commutation rule $\delta_{h} t=(t-h) \delta_{h}$ represents the action of the shift operator on polynomials. Forming equations over $S_{h}$, we model time-delay (resp., timeadvance) systems if $h>0$ (resp., $h<0$ ).

Example 3. For differential time-delay linear systems, we mix the constructions of the two preceding examples. For $h \in k$ (e.g., $k=\mathbb{Q}, \mathbb{R}$ ), we define the Ore algebra $D_{h}(k)=k[t]\left[\partial ; \sigma_{1}, \delta_{1}\right]\left[\delta_{h} ; \sigma_{2}, \delta_{2}\right]$ where:

$$
\sigma_{1}=\operatorname{id}_{k[t]}, \quad \delta_{1}=\frac{d}{d t}, \quad \forall a \in k[t], \quad \sigma_{2}(a)(t)=a(t-h), \quad \delta_{2}=0
$$

If the considered system also involves an advance operator, then we may work with the algebra defined by

$$
H_{(h, l)}(k)=k[t]\left[\partial ; \sigma_{1}, \delta_{1}\right]\left[\delta_{h} ; \sigma_{2}, \delta_{2}\right]\left[\tau_{l} ; \sigma_{3}, \delta_{3}\right]
$$

where $\sigma_{i}, \delta_{i}, i=1,2$, are as above and:

$$
\forall a \in k[t], \quad \sigma_{3}(a)(t)=a(t+l), \quad \delta_{3}=0, \quad l>0
$$

Example 4. In order to study multidimensional discrete linear systems, we can define the following Ore algebra $k\left[z_{1}, \ldots, z_{n}\right]\left[\partial_{1} ; \sigma_{1}, \delta_{1}\right] \ldots\left[\partial_{n} ; \sigma_{n}, \delta_{n}\right]$, where $\sigma_{i}$ and $\delta_{i}, i=1, \ldots, n$, are the maps on $k\left[z_{1}, \ldots, z_{n}\right]$ defined by $\delta_{i}=0$ and:

$$
\forall a \in k\left[z_{1}, \ldots, z_{n}\right], \quad \sigma_{i}(a)\left(z_{1}, \ldots, z_{n}\right)=a\left(z_{1}, \ldots, z_{i-1}, z_{i}+1, z_{i+1}, \ldots, z_{n}\right)
$$

We refer to 7 for more examples of Ore algebras using for instance the difference, the divided differences or the $q$-dilation functional operators.

We can "concatenate" different Ore algebras in order to combine different types of functional operators and, by this means, we get Ore algebras for most of the linear systems commonly considered in control theory. Moreover, we can also use different rings of coefficients such as the field of rational functions or the ring of analytic functions. However, as we shall develop computational aspects, we only consider here polynomial or rational coefficients over $\mathbb{Q}$. Finally, we can prove that the algebras defined in Examples 1, 2, 3 and 4 are left and right noetherian rings (namely, every left/right ideal is generated by means of a finite number of elements). Thus, they have the left and right Ore properties (namely, for any pair $\left(a_{1}, a_{2}\right)$ of elements, there exists a non-zero pair $\left(b_{1}, b_{2}\right)$ (resp., $\left.\left(c_{1}, c_{2}\right)\right)$ such that $b_{1} a_{1}=b_{2} a_{2}$ (resp., $a_{1} c_{1}=a_{2} c_{2}$ )) [5, 7, 11, 25.

Linear systems studied in control theory are generally defined by means of systems of ordinary or partial differential equations, time-delay equations, recurrence equations... These equations usually come from mathematical models. Hence, we can generally write a system as $R z=0$, where $R$ is a matrix with entries in a certain Ore algebra and $z$ contains the system variables including the inputs, the outputs, the states, the latent variables.

Example 5. - The linear system $P\left(\frac{d}{d t}\right) y=Q\left(\frac{d}{d t}\right) u$, where $P$ and $Q$ are two polynomial matrices in the differential operator $\frac{d}{d t}$ and with coefficients in $k[t]$, can be rewritten as $R z=0$, where the entries of the matrix

$$
R=\left(P\left(\frac{d}{d t}\right), \quad-Q\left(\frac{d}{d t}\right)\right)
$$

belong to the Weyl algebra $A_{1}(k)$ and $z=\left(y^{T}, u^{T}\right)^{T}$.

- The differential time-delay linear system $\dot{x}(t)=A(t) x(t)+B(t) u(t-h)$, where $A$ and $B$ are two matrices with entries in $k[t]$ and $h>0$, can be rewritten as $R z=0$, where the entries of the matrix

$$
R=\left(\begin{array}{cc}
\frac{d}{d t} I-A(t), & -B(t) \delta_{h}
\end{array}\right)
$$

belong to the Ore algebra $D_{h}(k)$ and $z=\left(x^{T}, u^{T}\right)^{T}$.

- The partial differential equation (heat equation)

$$
\frac{\partial y(t, x)}{\partial t}=\frac{\partial}{\partial x}\left(a(x) \frac{\partial y(t, x)}{\partial x}\right)+u(t, x)
$$

where the conductivity of the bar $a$ is assumed to be polynomial in $x$, can be rewritten as $R z=0$, where the entries of the matrix

$$
R=\left(\begin{array}{ll}
\frac{\partial}{\partial t}-\frac{\partial}{\partial x}\left(a(x) \frac{\partial}{\partial x}\right), & -1
\end{array}\right)
$$

belong to the Weyl algebra $A_{2}(k)$ with $x_{1}=t, x_{2}=x$ and $z=(y, u)^{T}$.
Real systems are generally nonlinear ones, meaning that the theory developed in this paper is not directly applicable to these systems. However, using a linearization around a (generic/given) trajectory of the system, then the linearized system has varying coefficients. Therefore, we can examine the structural properties of the linearized system by means of the approach described here and use them to study the ones of the nonlinear system.

## 3 A Module-Theoretical Approach to Linear Systems

In what follows, we denote by $D$ an Ore algebra. The main idea of algebraic analysis is to study a linear system of the form $R z=0$, where $R \in D^{q \times p}$, by means of the finitely presented (f.p.) left $D$-module $M=D^{1 \times p} /\left(D^{1 \times q} R\right)$ [24]. $M$ is associated with $R z=0$ in the sense that, if we denote by $z_{i}$ the residue class in $M$ of the row vector $e_{i} \in D^{1 \times p}$ defined by 1 in the $i^{\text {th }}$ position and 0 elsewhere and $z=\left(z_{1}, \ldots, z_{p}\right)^{T}$, then $M$ is defined by all left $D$-linear combinations of the system equations $R z=0$. See [11, 28, 31, 47] for more details.

The use of the residue class left $D$-module $M$ is natural as it is a generalization of the construction of the algebras commonly studied in algebraic or analytic
geometry and number theory (e.g., $\mathbb{C}=\mathbb{R}[x] /\left(x^{2}+1\right), \mathbb{Z}[i \sqrt{5}]=\mathbb{Z}[x] /\left(x^{2}+5\right)$, $\left.A=\mathbb{C}[x, y] /\left(x^{2}+y^{2}-1, x y-1\right)\right)$ [17]. For instance, $A=\mathbb{C}[x, y] / I$, where $I$ is the ideal $I=\left(x^{2}+y^{2}-1, x y-1\right)$ of $\mathbb{C}[x, y]$, can be defined by:

$$
A=\mathbb{C}[x, y] /\left(\mathbb{C}[x, y]^{1 \times 2} R\right), \quad R=\binom{x^{2}+y^{2}-1}{x y-1} \in \mathbb{C}[x, y]^{2 \times 1}
$$

The first main interest regarding the left $D$-module $M$ instead of the system $R z=0$ is that $M$ is intrinsically well-defined in the sense that it does not depend on the choice of the representation $R z=0$ of the system. Indeed, the same system can be represented in different equivalent forms having different numbers of unknowns and equations (e.g., state-space or input-output representations, Roesser or Fornasini-Marchesini models) [35, 38, 45].

The second main interest of using the finitely presented left $D$-module $M$ is that we can classify the structural properties of the system by means of the module properties of $M$. We introduce a few definitions [25, 45].

Definition 2. Let $M$ be a finitely generated left module over a left noetherian domain $D$. Then, we have the following definitions:

1. $M$ is free if there exists $r \in \mathbb{Z}_{+}$such that $M$ is isomorphic to $D^{1 \times r}$, a fact that we denote by $M \cong D^{1 \times r}$.
2. $M$ is stably free if there exist $r, s \in \mathbb{Z}_{+}$such that $M \oplus D^{1 \times s} \cong D^{1 \times r}$, where $\oplus$ denotes the direct sum.
3. $M$ is projective if there exist a left $D$-module $N$ and $r \in \mathbb{Z}_{+}$such that we have $M \oplus N \cong D^{1 \times r}$. Then, the left $D$-module $N$ is also projective.
4. $M$ is reflexive if the following canonical $D$-morphism (i.e., $D$-linear map)

$$
\varepsilon_{M}: M \longrightarrow \operatorname{hom}_{D}\left(\operatorname{hom}_{D}(M, D), D\right), \quad \varepsilon_{M}(m)(f)=f(m)
$$

- where $m \in M$ and $f$ belongs to the right $D$-module $\operatorname{hom}_{D}(M, D)$ formed by the left $D$-morphisms from $M$ to $D$ - is an isomorphism (i.e., $\varepsilon_{M}$ is both injective and surjective).

5. $M$ is torsion-free if the left $D$-submodule

$$
t(M)=\{m \in M \mid \exists 0 \neq P \in D, P m=0\}
$$

of $M$ is reduced to $0 . t(M)$ is called the torsion left $D$-submodule of $M$ and the elements of $t(M)$ are the torsion elements of $M$.
6. $M$ is torsion if $t(M)=M$.

We have the following important results [25, 45].
Theorem 1. 1. We have the following implications of module properties:

$$
\text { free } \Rightarrow \text { stably free } \Rightarrow \text { projective } \Rightarrow \text { reflexive } \Rightarrow \text { torsion-free. }
$$

2. Every torsion-free left module over $A_{1}(k)$ (resp., $k\left[\frac{d}{d t}\right], k(t)\left[\frac{d}{d t}\right]$ ) is stably free (resp., free).
3. Every projective module over the commutative polynomial ring $k\left[x_{1}, \ldots, x_{n}\right]$ over a field $k$ is free (Quillen-Suslin theorem).

In the recent years, a classification of properties of multidimensional linear systems has been established in terms of the properties of the corresponding left $D$-module $M$. Let us summarize some of them in Table 1. We refer the reader to $[15,28,26,29,31,32,33,34,37,38,47,49]$ for the precise definitions of the properties listed in the second and third column of Table 1.

Table 1. Classification of structural properties

| Module $M$ | Structural properties | Optimal control |
| :---: | :---: | :---: |
| Torsion | Poles/zeros classifications |  |
| With torsion | Existence of autonomous elements |  |
| Torsion-free | No autonomous elements, <br> Controllability, <br> Parametrizability, <br> $\pi$-flatness | Variational problem <br> without constraints <br> (Euler-Lagrange <br> equations) |
| Reflexive | Filter identification | Computation of the |
| Projective | Internal stabilizability, <br> Bézout identities, <br> Stabilizing controllers | wagrange parameters <br> without integration |
| Free | Flatness, Poles placement, <br> Doubly coprime factorization, <br> Youla-Kučera parametrization | Optimal controller |

## 4 Homological Algebra

The main issue of checking effectively the system properties via the properties of modules defined in Section 3 was still open until recently. Only the case of multidimensional systems defined by a full row rank matrix $R$ with entries in
the commutative polynomial ring $k\left[x_{1}, \ldots, x_{n}\right]$ was known using the different concepts of primeness [26, 34, 48, 49] developed in the middle of the seventies.

The concepts of syzygy modules, free resolutions, extension and torsion functors, projective and homotopic equivalences, projective dimensions... developed in homological algebra [45] form the basis of new algorithms checking the first column of Table 1, and thus, the system properties. These algorithms were obtained in [38] in the case of PDEs (see also [31, 32, 33, 34]).

We have recently shown in [9, 11] how these algorithms could be extended to some classes of Ore algebras including the interesting ones from the control theory point of view (e.g., ODEs, PDEs, recurrence operators, time-delay operators). The main steps of the algorithms developed in [9, 11] are:

1. Computation of free resolutions of f.p. left modules over an Ore algebra.
2. Dualization of the previous free resolutions using the $\operatorname{hom}_{D}(\cdot, D)$ functor.
3. Use of involutions in order to pass from right to left $D$-modules.
4. Computation of the quotient module of f.p. left $D$-modules.

Using the previous four points, we can then compute the extension modules $\operatorname{ext}_{D}^{i}(M, D), i \in \mathbb{Z}_{+}$, of any left $D$-module of the form $M=D^{1 \times p} /\left(D^{1 \times q} R\right)$. Let us explain the previous concepts. See [45] for more details.

Definition 3. We have the following definitions:

- A complex of left $D$-modules is a sequence formed by left $D$-modules $P_{i}$ and left $D$-morphisms $d_{i}: P_{i} \longrightarrow P_{i-1}$ which satisfy $\operatorname{im} d_{i+1} \subseteq \operatorname{ker} d_{i}$ for all $i \in \mathbb{Z}_{+}$. Such a complex is denoted by:

$$
\begin{equation*}
\ldots \xrightarrow{d_{i+2}} P_{i+1} \xrightarrow{d_{i+1}} P_{i} \xrightarrow{d_{i}} P_{i-1} \xrightarrow{d_{i-1}} P_{i-2} \xrightarrow{d_{i-2}} \ldots \tag{2}
\end{equation*}
$$

- The left $D$-module $H\left(P_{i}\right)=\operatorname{ker} d_{i} / \operatorname{im} d_{i+1}$ is called the defect of exactness of (21) at $P_{i}$. The complex (2) is said to be exact at $P_{i}$ if $H\left(P_{i}\right)=0$, i.e., $\operatorname{ker} d_{i}=\operatorname{im} d_{i+1}$, and exact if $H\left(P_{i}\right)=0$ for all $i \in \mathbb{Z}_{+}$.
- Let $M=D^{1 \times p} /\left(D^{1 \times q} R\right)$ be a finitely presented left $D$-module. A free resolution of $M$ is an exact sequence of the form

$$
\begin{equation*}
\ldots \xrightarrow{. R_{3}} D^{1 \times p_{2}} \xrightarrow{. R_{2}} D^{1 \times p_{1}} \xrightarrow{. R_{1}} D^{1 \times p_{0}} \xrightarrow{\pi} M \longrightarrow 0 \tag{3}
\end{equation*}
$$

where $p_{0}=p, p_{1}=q, R_{1}=R, R_{i} \in D^{p_{i} \times p_{i-1}}$ and $. R_{i}: D^{1 \times p_{i}} \longrightarrow D^{1 \times p_{i-1}}$ is defined by $\left(. R_{i}\right)(\lambda)=\lambda R_{i}$ for all $\lambda \in D^{1 \times p_{i}}$.

- Let us consider the free resolution (3) of $M$ and the following complex

$$
\begin{equation*}
\ldots \stackrel{R_{3}}{\longleftarrow} D^{p_{2}} \stackrel{R_{2}}{\longleftarrow} D^{p_{1}} \stackrel{R_{1} \cdot}{\longleftarrow} D^{p_{0}} \longleftarrow 0 \tag{4}
\end{equation*}
$$

where $R_{i}$. : $D^{p_{i-1}} \longrightarrow D^{p_{i}}$ is defined by $\left(R_{i}.\right)(\lambda)=R_{i} \lambda$ for all $\lambda \in D^{p_{i-1}}$. Then, the defects of exactness of the complex (4) are denoted by:

$$
\left\{\begin{array}{l}
\operatorname{ext}_{D}^{0}(M, D)=\operatorname{ker}\left(R_{1} .\right) \\
\operatorname{ext}_{D}^{i}(M, D)=\operatorname{ker}\left(R_{i+1} \cdot\right) /\left(R_{i} D^{1 \times p_{i-1}}\right), \quad i \geq 1
\end{array}\right.
$$

$\operatorname{ext}_{D}^{i}(M, D)$ inherits a right module structure by the right action of $D$.

Proposition 1. 45] The right $D$-module $\operatorname{ext}_{D}^{i}(M, D)$ only depends on $M$, i.e., we can choose any free resolution of $M$ to compute $\operatorname{ext}_{D}^{i}(M, D), i \in \mathbb{Z}_{+}$. Moreover, we have $\operatorname{ext}_{D}^{0}(M, D)=\operatorname{hom}_{D}(M, D)$.

Coming back to the four main algorithmic steps outlined above, we recall that an involution $\theta$ of $D$ is a $k$-linear map $\theta: D \longrightarrow D$ satisfying:

$$
\begin{equation*}
\forall a_{1}, a_{2} \in D, \quad \theta\left(a_{1} \cdot a_{2}\right)=\theta\left(a_{2}\right) \cdot \theta\left(a_{1}\right), \quad \theta \circ \theta=\operatorname{id}_{D} \tag{5}
\end{equation*}
$$

Example 6. We have the following examples of involutions:

1. If $D=k\left[x_{1}, \ldots, x_{n}\right]$ is a commutative polynomial algebra, then $\theta=i d$ is a trivial involution.
2. If $D=A_{n}$ is the Weyl algebra and $P \in D$, then we let $\theta(P)$ be the classical formal adjoint of $P$ obtained by multiplying a test function on the left of $P z$ and by integrating by parts [31, 33, 34]. Equivalently, $\theta$ is defined by $\theta\left(x_{i}\right)=x_{i}$ and $\theta\left(\partial_{i}\right)=-\partial_{i}, i=1, \ldots, n$.
3. Let $S_{h}(k)$ be the Ore algebra of shift operators defined in Example 2, Then, an involution $\theta$ of $S_{h}(k)$ is defined by $\theta(t)=-t$ and $\theta\left(\delta_{h}\right)=\delta_{h}$.
4. If $D_{h}(k)$ is the Ore algebra of differential time-delay operators defined in Example 3, then an involution $\theta$ of $D_{h}(k)$ can be defined by $\theta(t)=-t$, $\theta\left(\delta_{h}\right)=\delta_{h}$ and $\theta(\partial)=\partial$. This last result shows that a simple involution of $D_{h}(k)$ exists contrary to what was written in [11] (we thank V. Levandovskyy for pointing out to us this trivial mistake).

Now, if $R$ is a matrix with entries in an Ore algebra having an involution $\theta$ (e.g., $\left.A_{n}(k), S_{h}(k), D_{h}(k)\right)$, then we can define $\theta(R)=\left(\theta\left(R_{i j}\right)\right)^{T}$ and the left $D$ module $\widetilde{N}=D^{1 \times q} /\left(D^{1 \times p} \theta(R)\right)$. The main idea developed in 9, 11, 31, 34, 38, is that the module properties in the first column of Table 1 are characterized by the vanishing of $\operatorname{certain}_{\operatorname{ext}}^{D}{ }_{D}(\widetilde{N}, D)$ as it is shown in Table 2. We refer the reader to [42] for a constructive algorithm which checks freeness and computes bases of free modules.

The last column of Table 2 explains the correspondence between module properties and primeness for a multidimensional system defined by a full row rank matrix $R$ with entries in the commutative polynomial ring $D=k\left[x_{1}, \ldots, x_{n}\right]$ [34]. The third column generalizes the last column to multidimensional systems defined by a full row rank matrix $R$ with entries in the ring of differential operators with rational coefficients and $d(\widetilde{N})$ denotes the Krull dimension of the characteristic variety of $\tilde{N}$ (see [34, 39).

## 5 Computation of $\operatorname{ext}_{D}^{i}(\widetilde{N}, D)$

The main difficulty in the computation of $\operatorname{ext}_{D}^{i}(\widetilde{N}, D)$ is to be able to construct a free resolution for the left $D$-module $\widetilde{N}=D^{1 \times q} /\left(D^{1 \times p} \theta(R)\right)$ (see point 1 in the previous section), i.e., an exact sequence of the form

$$
\ldots \xrightarrow{\widetilde{R}_{4}} D^{1 \times q_{3}} \xrightarrow{\widetilde{R}_{3}} D^{1 \times q_{2}} \xrightarrow{\widetilde{R}_{2}} D^{1 \times q_{1}} \xrightarrow{\widetilde{R}_{1}} D^{1 \times q_{0}} \longrightarrow \widetilde{N} \longrightarrow 0,
$$

where $\widetilde{R}_{1}=\theta(R), q_{0}=q, q_{1}=p$ and $. \widetilde{R}_{i}: D^{1 \times q_{i}} \rightarrow D^{1 \times q_{i-1}}$ is defined by $\left(. \widetilde{R}_{i}\right)(\lambda)=\lambda \widetilde{R}_{i}$. The left $D$-module

$$
S_{i}(\widetilde{N})=\operatorname{ker}\left(. \widetilde{R}_{i}\right)=\left\{\lambda \in D^{1 \times q_{i}} \mid \lambda \widetilde{R}_{i}=0\right\}
$$

is called the $i^{\text {th }}$ syzygy left $D$-module of $\widetilde{N}$. If $D$ is a noetherian ring, which is the case for a large class of algebras (e.g., $A_{n}, S_{h}, D_{h}$ and $H_{(h, l)}$ ) [9, 11, then free resolutions always exist for finitely generated left $D$-modules [25, 45].

The computation of the matrix $\widetilde{R}_{i+1}$ is an elimination problem [2, 17]. Indeed, multiplying $\lambda \in S_{i}(\widetilde{N})$ on the left of the inhomogeneous system $\widetilde{R}_{i} y=u$, we then obtain $\lambda u=0$. Hence, finding a family of generators for $S_{i}(\widetilde{N})$, i.e., $\left\{\lambda_{j}\right\}_{1 \leq j \leq q_{i+1}}$, $\lambda_{j} \in D^{1 \times q_{i}}$ satisfying $S_{i}(\widetilde{N})=D \lambda_{1}+\ldots+D \lambda_{q_{i+1}}$ is equivalent to finding a family of generators for the compatibility conditions of the inhomogeneous system $\widetilde{R}_{i} y=u$. Then, if we denote by $\widetilde{R}_{i+1}=\left(\lambda_{1}^{T}, \ldots, \lambda_{q_{i+1}}^{T}\right)^{T}$, we finally obtain $S_{i}(\widetilde{N})=D^{1 \times q_{i+1}} \widetilde{R}_{i+1}$.

Such a difficult problem has largely been studied for linear systems of PDEs since the $19^{\text {th }}$ century 31, 38, 39. But, only recently some computational answers were found based on the concepts of Janet and Gröbner bases for noncommutative polynomial rings. We recall the definition of Gröbner bases for polynomial ideals. This definition can easily be extended to modules [2, 17. The algorithmic methods used in the theory of Gröbner bases require that a monomial order is chosen to compare polynomials.

Table 2. Characterization of module properties

| Module $M$ | $\operatorname{ext}^{i}(\widetilde{N}, D)$ | $d(\widetilde{N})$ | Primeness |
| :---: | :---: | :---: | :---: |
| With torsion | $\operatorname{ext}_{D}^{1}(\widetilde{N}, D) \cong t(M)$ | $n-1$ | $\emptyset$ |
| Torsion-free | $\operatorname{ext}_{D}^{1}(\widetilde{N}, D)=0$ | $n-2$ | Minor left-prime |
| Reflexive | $\begin{gathered} \operatorname{ext}_{D}^{i}(\tilde{N}, D)=0, \\ i=1,2 \end{gathered}$ | $n-3$ |  |
| Projective | $\begin{gathered} \operatorname{ext}_{D}^{i}(\widetilde{N}, D)=0, \\ 1 \leq i \leq n \end{gathered}$ | -1 | Zero left-prime |

Definition 4. 1. Let $D$ be an Ore algebra. A monomial order $<$ on $D$ is defined as a total order on the set of monomials $\operatorname{Mon}(D)$ satisfying the following two conditions:
a) For all monomials $m \in \operatorname{Mon}(D) \backslash\{1\}$, we have $1<m$.
b) If $m_{1}<m_{2}$ holds for two monomials $m_{1}, m_{2} \in \operatorname{Mon}(D)$, then, for all $n \in \operatorname{Mon}(D)$, we have $n \cdot m_{1}<n \cdot m_{2}$.
2. Given a polynomial $P \in D \backslash\{0\}$ and a monomial order $<$ on $D$, we can compare the monomials with a non-zero coefficient in $P$ w.r.t. $<$. The greatest of these monomials is the leading monomial $\operatorname{lm}(P)$ of $P$.

Definition 5. 2, 17] Let $D$ be a polynomial ring and $I$ a (left) ideal of $D$. A set of non-zero polynomials $G=\left\{g_{1}, \ldots, g_{t}\right\} \subset I$ is called a Gröbner basis for $I$ if for all $0 \neq f \in I$, there exists $1 \leq i \leq t$ such that $\operatorname{lm}\left(g_{i}\right)$ divides $\operatorname{lm}(f)$.

One consequence of the condition that defines Gröbner bases is that every polynomial $f$ in $I$ is reduced to 0 modulo $G$, i.e., by subtraction of suitable left multiples of the $g_{i} \in G$ from $f$, we then obtain the zero polynomial.

For the case of commutative polynomial rings, Buchberger's algorithm [2, 17) computes Gröbner bases of polynomial ideals. Recently, Buchberger's algorithm was extended to some non-commutative polynomial rings and, in particular, to some classes of Ore algebras [5, 7] that are important for the study of multidimensional linear systems. Hence, manipulations of (one-sided) ideals and modules over many classes of Ore algebras have been turned effective. Moreover, the Maple library Mgfun [6] has been developed for the symbolic manipulation of a large class of special functions and combinatorial sequences. It offers implementations of Gröbner bases for some classes of Ore algebras.

## 6 The Package OreModules

Using the Maple library Mgfun, the authors of this paper have recently been developing the package OreModules [8, 10]. OreModules as well as a library of examples are freely available at:

```
http://wwwb.math.rwth-aachen.de/OreModules.
```

This second release of OreModules focuses on the following problems:

- Compute free resolutions, formal adjoints, extension functors, duals and biduals of f.p. left $D$-modules over some classes of Ore algebras $D$.
- Recognize the properties of a finitely presented left $D$-module $M$ (torsionfree, reflexive, projective, stably free, free).
- Decide the existence of torsion elements in the corresponding system and, if so, compute a family of generators for them.
- Compute left/right/generalized inverses of matrices with entries in $D$.
- Check whether or not a multidimensional linear system is controllable in the sense of [14, 15, 26, 30, 29, 31, 32, 47, 49] or compute the autonomous elements of the system [30, 31, 32, 47, 49.
- Check whether or not a multidimensional linear system is parametrizable in the sense of 11, 31, 32, 33].
- Check whether or not a multidimensional linear system is flat and, if so, compute an injective parametrization and a flat output [15, 26, 31, 33, 42.
- Check whether or not a multidimensional linear system with constant coefficients is $\pi$-free and, if so, compute the ideal of all the $\pi$-polynomials [11, 15, 26].
A list of the most important functions of OreModules is given in Table 3 (the suffix "Rat" distinguishes the procedures which deal with polynomial/rational coefficients). Detailed documentation of OreModules is available in form of Maple help pages.


## 7 Worked Examples Using OreModules

OreModules comes with a library of examples which demonstrates the above features by means of systems like two pendula mounted on a cart, stirred tank models, electric transmission line, wind tunnel model, Maxwell equations, Einstein equations, equations of linear elasticity, Lie-Poisson structures... We only give here four simple examples but we refer the reader to [8, 9, 11] for more sophisticated examples. All examples were run on a Pentium III, 1 GHz with 1 GB RAM using Maple 8 (OreModules is available for Maple V release 5, Maple 6, Maple 8, Maple 9 and Maple 10).

Example 7. We study a linearized bipendulum [31, i.e., a system composed of a bar, where two pendula are fixed, one of length $l 1$ and one of length $l 2$. The appropriate Ore algebra for this example is the Weyl algebra $\operatorname{Alg}=A_{1}$, i.e., $A_{1}=\mathbb{Q}(l 1, l 2, g)[t][\mathrm{D}]$, where $\mathrm{D}=\frac{d}{d t}$ is the differential operator w.r.t. time $t$ :

```
> with(OreModules):
> Alg:=DefineOreAlgebra(diff=[D,t],polynom=[t],comm=[g,l1,l2]):
```

Note that we have to declare all constants appearing in the system equations (the gravitational constant $g$ and the lengths $l 1$ and $l 2$ ) as variables that commute with D and $t$. Next, we enter the system matrix:

$$
\begin{gathered}
>R:=\operatorname{evalm}\left(\left[\left[\mathrm{D}^{\wedge} 2+\mathrm{g} / l 1,0,-\mathrm{g} / \mathrm{l1}\right],[0, \quad \mathrm{D} \wedge 2+\mathrm{g} / 12,-\mathrm{g} / 12]\right]\right) ; \\
R:=\left[\begin{array}{ccc}
\mathrm{D}^{2}+\frac{g}{l 1} & 0 & -\frac{g}{l 1} \\
0 & \mathrm{D}^{2}+\frac{g}{l 2} & -\frac{g}{l 2}
\end{array}\right]
\end{gathered}
$$

In terms of equations, the linearized bipendulum is described by:

$$
\begin{aligned}
& >\text { ApplyMatrix }(\mathrm{R},[\mathrm{x} 1(\mathrm{t}), \mathrm{x} 2(\mathrm{t}), \mathrm{u}(\mathrm{t})], \mathrm{Alg})=\text { evalm }([[0],[0]]) ; \\
& \qquad\left[\begin{array}{l}
\frac{g \mathrm{x} 1(t)}{l 1}+\left(\frac{d^{2}}{d t^{2}} \mathrm{x} 1(t)\right)-\frac{g \mathrm{u}(t)}{l 1} \\
\frac{g \mathrm{x} 2(t)}{l 2}+\left(\frac{d^{2}}{d t^{2}} \mathrm{x} 2(t)\right)-\frac{g \mathrm{u}(t)}{l 2}
\end{array}\right]=\left[\begin{array}{l}
0 \\
0
\end{array}\right]
\end{aligned}
$$

Table 3. List of the most important functions of OreModules


We compute the formal adjoint of $R$ :

```
> R_adj:=Involution(R, Alg);
```

$$
R_{\_} a d j:=\left[\begin{array}{cc}
\mathrm{D}^{2}+\frac{g}{l 1} & 0 \\
0 & \mathrm{D}^{2}+\frac{g}{l 2} \\
-\frac{g}{l 1} & -\frac{g}{l 2}
\end{array}\right]
$$

By computing $\operatorname{ext}_{A_{1}}^{1}\left(A_{1}^{1 \times 2} /\left(A^{1 \times 3} R_{\_} a d j\right), A_{1}\right)$, we check whether or not the left $A_{1}$-module $M=A_{1}^{1 \times 3} /\left(A_{1}^{1 \times 2} R\right)$ is torsion-free, i.e., whether or not the bipendulum is controllable and parametrizable:
> Ext:=Exti(R_adj, Alg, 1);

$$
\left.\left.\left.\begin{array}{c}
\text { Ext }:=\left[\left[\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right],\left[\begin{array}{c}
\mathrm{D}^{2} l 1+g \\
0
\end{array} \mathrm{D}^{2} l 2+g-g\right.\right.
\end{array}\right], ~ \begin{array}{c}
l 2 \mathrm{D}^{2} g+g^{2} \\
g^{2}+\mathrm{D}^{2} l 1 g \\
l 2 \mathrm{D}^{2} g+l 2 l 1 \mathrm{D}^{4}+\mathrm{D}^{2} l 1 g+g^{2}
\end{array}\right]\right] .\left[\begin{array}{c} 
\\
l
\end{array}\right.
$$

From the output, we can see that the system is generically controllable because $\operatorname{Ext}[1]$ is the identity matrix which means that there are no torsion elements in the left $A_{1}$-module $M$ associated with the system. The interpretation of this structural fact is that the system has no autonomous elements in the generic case (see Section (3). There may be a few configurations of the constants $g, l 1$, $l 2$, in which the bipendulum is not controllable. We will actually find the only configuration where it is not controllable below. Let us write down the generic parametrization $\operatorname{Ext}[3]$ in a more familiar way with a free function $\xi_{1}$.

$$
\begin{aligned}
& >P:=P a r a m e t r i z a t i o n(R, A l g) ; \\
& \\
& P:=\left[\begin{array}{c}
g\left(g \xi_{1}(t)+l 2 \frac{d^{2}}{d t^{2}} \xi_{1}(t)\right. \\
g\left(g \xi_{1}(t)+l 1 \frac{d^{2}}{d t^{2}} \xi_{1}(t)\right. \\
g^{2} \xi_{1}(t)+g l 2 \frac{d^{2}}{d t^{2}} \xi_{1}(t)+g l 1 \frac{d^{2}}{d t^{2}} \xi_{1}(t)+l 1 l 2\left(\frac{d^{4}}{d t^{4}} \xi_{1}(t)\right)
\end{array}\right]
\end{aligned}
$$

Therefore, all smooth solutions of the system are parametrized by $P$, i.e.,

$$
R\left(x_{1}, x_{2}, u\right)^{T}=0 \Leftrightarrow\left(x_{1}, x_{2}, u\right)^{T}=\operatorname{Ext}[3] \xi_{1}=P \xi_{1} .
$$

Since the bipendulum is generically a time-invariant controllable system, it is also generically a flat system. A flat output of the system can be computed as a left-inverse of the parametrization $\operatorname{Ext}[3]$ :
> S:=LeftInverse(Ext[3], Alg);

$$
S:=\left[\frac{l 1}{g^{2}(l 1-l 2)}-\frac{l 2}{g^{2}(l 1-l 2)} 0\right]
$$

i.e., a flat output is defined by $\xi_{1}=S\left(x_{1}, \quad x_{2}, \quad u\right)^{T}$, namely:

$$
\begin{gathered}
>\mathrm{xi}[1](\mathrm{t})=A p p l y \operatorname{Matrix}(\mathrm{~S},[\mathrm{x} 1(\mathrm{t}), \mathrm{x} 2(\mathrm{t}), \mathrm{u}(\mathrm{t})], \mathrm{Alg})[1,1] ; \\
\xi_{1}(\mathrm{t})=\frac{l 1 \mathrm{x} 1(t)}{g^{2}(l 1-l 2)}-\frac{l 2 \mathrm{x} 2(t)}{g^{2}(l 1-l 2)}
\end{gathered}
$$

We remark that this flat output is defined only if $l 1-l 2 \neq 0$. In this case, let us compute the Brunovsky canonical form of the system.

$$
\begin{aligned}
& >B:=\operatorname{Brunovsky}(\mathrm{R}, \mathrm{Alg}) ; \\
& \qquad B:=\left[\begin{array}{ccc}
\frac{l 1}{g^{2}(l 1-l 2)} & -\frac{l 2}{g^{2}(l 1-l 2)} & 0 \\
\frac{\mathrm{D} l 1}{g^{2}(l 1-l 2)} & -\frac{\mathrm{D} l 2}{g^{2}(l 1-l 2)} & 0 \\
-\frac{1}{g(l 1-l 2)} & \frac{1}{g(l 1-l 2)} & 0 \\
-\frac{\mathrm{D}}{g(l 1-l 2)} & \frac{\mathrm{D}}{g(l 1-l 2)} & 0 \\
\frac{1}{(l 1-l 2) l 1} & -\frac{1}{(l 1-l 2) l 2} & \frac{1}{l 1 l 2}
\end{array}\right]
\end{aligned}
$$

In other words, we have the following transformation between the system variables $x_{1}, x_{2}$ and $u$ and the Brunovský variables $z_{i}, i=1, \ldots, 4$, and $v$ :

```
> evalm([seq([z[i](t)],i=1..4),[v(t)]])=ApplyMatrix(B,
> [x1(t),x2(t),u(t)],Alg);
```

$$
\left[\begin{array}{c}
z_{1}(t) \\
z_{2}(t) \\
z_{3}(t) \\
z_{4}(t) \\
\mathrm{v}(t)
\end{array}\right]=\left[\begin{array}{c}
\frac{l 1 \mathrm{x} 1(t)}{g^{2}(l 1-l 2)}-\frac{l 2 \times 2(t)}{g^{2}(l 1-l 2)} \\
\frac{l 1\left(\frac{d}{d t} \times 1(t)\right)}{g^{2}(l 1-l 2)}-\frac{l 2\left(\frac{d}{d t} \times 2(t)\right)}{g^{2}(l 1-l 2)} \\
-\frac{\mathrm{x} 1(t)}{g(l 1-l 2)}+\frac{\mathrm{x} 2(t)}{g(l 1-l 2)} \\
-\frac{\frac{d}{d t} \times 1(t)}{g(l 1-l 2)}+\frac{\frac{d}{d t} \times 2(t)}{g(l 1-l 2)} \\
\frac{\mathrm{x} 1(t)}{(l 1-l 2) l 1}-\frac{\mathrm{x} 2(t)}{(l 1-l 2) l 2}+\frac{\mathrm{u}(t)}{l 1 l 2}
\end{array}\right]
$$

Let us check that the new variables $z_{i}, i=1, \ldots, 4$, and $v$ satisfy a Brunovský canonical form:

```
> F:=Elimination(linalg[stackmatrix](B, R), [x1,x2,u],
> [seq(z[i],i=1..4),v,0,0], Alg):
> ApplyMatrix(F[1], [x1(t),x2(t),u(t)], Alg)=ApplyMatrix(F[2],
> [seq(z[i](t),i=1..4),v(t)],Alg);
```

$$
\left[\begin{array}{c}
0 \\
0 \\
0 \\
0 \\
\mathrm{u}(t) \\
\mathrm{x} 2(t) \\
\mathrm{x} 1(t)
\end{array}\right]=\left[\begin{array}{c}
-\left(\frac{d}{d t} z_{4}(t)\right)+\mathrm{v}(t) \\
-\left(\frac{d}{d d} z_{3}(t)\right)+z_{4}(t) \\
-\left(\frac{d}{d d} z_{2}(t)\right)+z_{3}(t) \\
-\left(\frac{d}{d t} z_{1}(t)\right)+z_{2}(t) \\
g^{2} z_{1}(t)+(g l 2+g l 1) z_{3}(t)+l 1 l 2 \mathrm{v}(t) \\
g^{2} z_{1}(t)+g l 1 z_{3}(t) \\
g^{2} z_{1}(t)+g l 2 z_{3}(t)
\end{array}\right]
$$

The last three equations give $u, x_{1}$ and $x_{2}$ in terms of the $z_{i}$ and $v$.
$l 1=l 2$ describes the only case in which the bipendulum may be uncontrollable. We now turn to the case where the lengths of the pendula are equal:

$$
\begin{aligned}
& >\text { R_mod:=subs }(12=11, \text { evalm (R)); } \\
& \qquad R \_m o d:=\left[\begin{array}{ccc}
\mathrm{D}^{2}+\frac{g}{l 1} & 0 & -\frac{g}{l 1} \\
0 & \mathrm{D}^{2}+\frac{g}{l 1} & -\frac{g}{l 1}
\end{array}\right] \\
& >\text { Ext_mod }:=\text { Exti (Involution(R_mod, Alg), Alg, 1); } \\
& \quad \text { Ext_mod }:=\left[\left[\begin{array}{rr}
\mathrm{D}^{2} l 1+g & 0 \\
0 & 1
\end{array}\right],\left[\begin{array}{ccc}
1 & -1 & 0 \\
0 & \mathrm{D}^{2} l 1+g & -g
\end{array}\right],\left[\begin{array}{c}
g \\
g \\
\mathrm{D}^{2} l 1+g
\end{array}\right]\right]
\end{aligned}
$$

The computation of $\operatorname{ext}_{A_{1}}^{1}\left(A_{1}^{1 \times 2} /\left(A_{1}^{1 \times 3} \theta\left(R_{m o d}\right)\right), A_{1}\right)$ gives the torsion submodule $t(M)$ of $M$ : it is generated by the residue class of the row $r$ of Ext_mod[2] which corresponds to the row with entry $l 1 \mathrm{D}^{2}+g$ in Ext_mod[1]. This means that $\left(l 1 \mathrm{D}^{2}+g\right) r=0$ in $M$, and the difference of the positions of the pendula (relative to the bar) is an autonomous element of the system. We can conclude that the bipendulum is controllable if and only if $l 1 \neq l 2$.

We can directly obtain the torsion elements of $M$ as follows:

```
> TorsionElements(R_mod, [x1(t),x2(t),u(t)],Alg);
```

$$
\left[\left[g \theta_{1}(t)+l 1\left(\frac{d^{2}}{d t^{2}} \theta_{1}(t)\right)=0\right],\left[\theta_{1}(t)=\mathrm{x} 1(t)-\mathrm{x} 2(t)\right]\right]
$$

We can also explicitly integrate this torsion element of $M$ :
$>$ AutonomousElements(R_mod, [x1(t), x2(t),u(t)], Alg)[2];

$$
\left[\theta_{1}={ }_{-} C 1 \sin \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)+{ }_{-} C 2 \cos \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)\right]
$$

The fact that there exists an autonomous element in the system is equivalent to the existence of a first integral of motion in the system. Indeed, we recall that we have a one-to-one correspondence between the torsion elements and the first integrals of motion. For more details, see [33]. We can compute this first integral of motion by using the command FirstIntegral:

$$
\begin{aligned}
> & \mathrm{V}:=\text { FirstIntegral (R_mod, }[\mathrm{x} 1(\mathrm{t}), \mathrm{x} 2(\mathrm{t}), \mathrm{u}(\mathrm{t})], \mathrm{Alg}) ; \\
V & :=-\left(-\left(\frac{d}{d t} \times 1(t)\right) \_C 1 \sin \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right) \sqrt{l 1}-\left(\frac{d}{d t} \times 1(t)\right) \_C 2 \cos \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right) \sqrt{l 1}\right. \\
& +\sqrt{g} \times 1(t) \_C 1 \cos \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)-\sqrt{g} \times 1(t) \_C 2 \sin \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right) \\
& +\left(\frac{d}{d t} \times 2(t)\right) \_C 1 \sin \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right) \sqrt{l 1}+\left(\frac{d}{d t} \times 2(t)\right) \_C 2 \cos \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right) \sqrt{l 1} \\
& \left.-\sqrt{g} \times 2(t) \_C 1 \cos \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)+\sqrt{g} \times 2(t) \_C 2 \sin \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)\right) / \sqrt{l 1}
\end{aligned}
$$

We let the reader check by himself that we have $\dot{V}(t)=0$. For the explicit computations, see [8].

Finally, even if we have some autonomous elements in the system, we can parametrize all solutions of the system in terms of one arbitrary function $\xi_{1}$ and two arbitrary constants $C_{1}$ and $C_{2}$ (these constants can easily be computed in terms of the initial conditions of the system):
> P2:=Parametrization(R_mod, Alg);

$$
P 2:=\left[\begin{array}{c}
g \xi_{1}(t) \\
--C 1 \sin \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)-\_C 2 \cos \left(\frac{\sqrt{g} t}{\sqrt{l 1}}\right)+g \xi_{1}(t) \\
l 1\left(\frac{d^{2}}{d t^{2}} \xi_{1}(t)\right)+g \xi_{1}(t)
\end{array}\right]
$$

i.e., we have $R\left(x_{1}, x_{2}, u\right)^{T}=0 \Leftrightarrow\left(x_{1}, x_{2}, u\right)^{T}=P_{2}\left(\xi_{1},{ }_{\_} C 1,{ }_{\_} C 2\right)$. We can easily check that $P_{2}$ gives a parametrization of some solutions of the system as we have:

```
> simplify(ApplyMatrix(R_mod, P2, Alg));
```

$$
\left[\begin{array}{l}
0 \\
0
\end{array}\right]
$$

We can prove that $P_{2}$ parametrizes all smooth solutions of the system 40.
Example 8. This example demonstrates the study of structural properties of a simple linear time-varying ordinary differential system [41, 43]. See [8] for more sophisticated examples.
$>$ Alg:=DefineOreAlgebra(diff=[D,t], polynom=[t]):
Let us consider the following matrix of ordinary differential operators:

```
> R:=evalm([[D, -t]]);
```

$$
R:=[\mathrm{D}-t]
$$

The matrix $R$ corresponds to the following time-varying linear system:
$>$ ApplyMatrix(R, $[x(t), u(t)], A l g)[1,1]=0$;

$$
\left(\frac{d}{d t} \mathrm{x}(t)\right)-t \mathrm{u}(t)=0
$$

Let us check whether or not this system is controllable and flat. In order to do that, let us define the formal adjoint $R \_a d j$ of $R$.
> R_adj:=Involution(R, Alg);

$$
R \_a d j:=\left[\begin{array}{c}
-\mathrm{D} \\
-t
\end{array}\right]
$$

We compute the first extension module $\operatorname{ext}_{A_{1}(\mathbb{Q})}^{1}\left(A_{1}(\mathbb{Q}) /\left(A_{1}(\mathbb{Q})^{1 \times 2} R_{\_} a d j\right)\right.$, $\left.A_{1}(\mathbb{Q})\right)$ of the left $A l g$-module associated with $R_{-} a d j$ :

$$
\begin{aligned}
& >\operatorname{Ext}:=\operatorname{Exti}\left(\mathrm{R}_{-} \operatorname{adj}, \mathrm{Alg}, 1\right) ; \\
& \qquad E x t:=\left[[1],[\mathrm{D}-t],\left[\begin{array}{cc}
-t^{2} & -1+t \mathrm{D} \\
-2-t \mathrm{D} & \mathrm{D}^{2}
\end{array}\right]\right.
\end{aligned}
$$

Therefore, we obtain that the left $A_{1}$-module $M=A_{1}(\mathbb{Q})^{1 \times 2} /\left(A_{1}(\mathbb{Q}) R\right)$ associated with $R$ is torsion-free, and thus, stably free as $A_{1}(\mathbb{Q})$ is a hereditary ring. A parametrization of the system is given by Ext[3]. This result can directly be obtained by using the following command:
> Parametrization(R, Alg);

$$
\left[\begin{array}{c}
-t^{2} \xi_{1}(t)-\xi_{2}(t)+t\left(\frac{d}{d t} \xi_{2}(t)\right) \\
-2 \xi_{1}(t)-t\left(\frac{d}{d t} \xi_{1}(t)\right)+\left(\frac{d^{2}}{d t^{2}} \xi_{2}(t)\right)
\end{array}\right]
$$

Let us notice that the previous parametrization depends on two arbitrary functions $\xi_{1}$ and $\xi_{2}$. However, the system has only 1 input, and thus, the rank of the left $A_{1}(\mathbb{Q})$-module $M$ is 1 . Let us check this result:

```
> OreRank(R, Alg);
```

Hence, we deduce that there exist some minimal parametrizations of the system which depend on 1 arbitrary function. Let us compute some of them.

```
> P:=MinimalParametrizations(R, Alg);
```

$$
P:=\left[\left[\begin{array}{c}
-t^{2} \\
-2-t \mathrm{D}
\end{array}\right],\left[\begin{array}{c}
-1+t \mathrm{D} \\
\mathrm{D}^{2}
\end{array}\right]\right]
$$

Let us check whether or not the first minimal parametrization $P[1]$ is injective.

```
> LeftInverse(P[1], Alg);
```

We obtain that $P[1]$ is not an injective parametrization of the system. Let us examine the second minimal parametrization $P[2]$ in a similar way:

```
> LeftInverse(P[2], Alg);
```

We find that $P[2]$ is not an injective parametrization of the system either. Therefore, we cannot conclude that the left $A_{1}(\mathbb{Q})$-module $M$ associated with the system is free. In fact, we can prove that $M$ is not a free left $A_{1}(\mathbb{Q})$-module, and thus, the corresponding time-varying system is not flat. See [39, 42] for more details. However, we already know that $M$ is a stably free left $A_{1}(\mathbb{Q})$ module as the matrix $R$ has full row rank and $R$ admits a right-inverse defined by:

```
> RightInverse(R, Alg);
```

$$
\left[\begin{array}{c}
t \\
\mathrm{D}
\end{array}\right]
$$

See 11 for more details. One of the main interests of the non-minimal parametrization $E x t[3]$ is that it admits a generalized inverse, namely, there exists a matrix $G$ with entries in $A_{1}(\mathbb{Q})$ satisfying $\operatorname{Ext}[3] G \operatorname{Ext}[3]=\operatorname{Ext}[3]$ (contrary to $P[1]$ and $P[2])$. This last result implies that the non-minimal parametrization parametrizes all the solutions of the corresponding time-varying system which belong to any left $A_{1}(\mathbb{Q})$-module $\mathcal{F}$ (e.g., $\left.\mathcal{F}=C^{\infty}(\mathbb{R}), \mathbb{R}(t), \mathbb{R}[t]\right)$. Let us compute one generalized inverse of $\operatorname{Ext}[3]$ :

$$
\begin{aligned}
& >\mathrm{G}:=\text { GeneralizedInverse(Ext[3], Alg); } \\
& \qquad G:=\left[\begin{array}{rr}
0 & -1 \\
-1 & 0
\end{array}\right] \\
& >\operatorname{Mult}(\operatorname{Ext}[3], \mathrm{G}, \operatorname{Ext}[3], \mathrm{Alg}) ; \\
& {\left[\begin{array}{cc}
-t^{2} & -1+t \mathrm{D} \\
-2-t \mathrm{D} & \mathrm{D}^{2}
\end{array}\right]}
\end{aligned}
$$

Let us determine the obstruction of flatness. In order to do that, we study the system over the ring $\mathbb{Q}(t)\left[\frac{d}{d t}\right]$ of ordinary differential operators with rational coefficients in $t$. Let us compute a parametrization of the system by allowing to invert non-zero polynomials in $t$ :

```
> Extrat:=ExtiRat(R_adj, Alg, 1);
```

$$
\text { Extrat }:=\left[[1],[\mathrm{D}-t],\left[\begin{array}{c}
-t^{2} \\
-2-t \mathrm{D}
\end{array}\right]\right]
$$

We obtain that the left $\mathbb{Q}(t)\left[\frac{d}{d t}\right]$-module $M^{\prime}=\mathbb{Q}(t)\left[\frac{d}{d t}\right]^{1 \times 2} /\left(\mathbb{Q}(t)\left[\frac{d}{d t}\right] R\right)$ is torsion-free, and thus, free because $\mathbb{Q}(t)\left[\frac{d}{d t}\right]$ is a left principal ideal domain.

Moreover, a (minimal) parametrization of the system is defined by Extrat[3]. This result can directly be obtained by using ParametrizationRat:
> ParametrizationRat(R, Alg);

$$
\left[\begin{array}{c}
-t^{2} \xi_{1}(t) \\
-2 \xi_{1}(t)-t\left(\frac{d}{d t} \xi_{1}(t)\right)
\end{array}\right]
$$

The fact that the left $\mathbb{Q}(t)[D]$-module $M$ associated with $R$ is free implies that, away from some singularities that we are going to determine, the system is flat. Let us compute a basis for this module which gives a flat output of the system.

$$
\begin{aligned}
& >\mathrm{S}:=\text { LeftInverseRat (Extrat [3], Alg); } \\
& \qquad S:=\left[-\frac{1}{t^{2}} 0\right]
\end{aligned}
$$

Therefore, we obtain that a basis of the left $\mathbb{Q}(t)[D]$-module $M$ is defined by $\xi_{1}$ $=S(x, u)^{T}$ and satisfies:

$$
(x, u)^{T}=\operatorname{Extrat}[3] \xi_{1} .
$$

In particular, we see that this parametrization is not defined for $t=0$ as we have a singularity. Therefore, the system is flat except for $t=0$. Finally, we note that, away from $t=0$, we have another right-inverse of $R$ defined by:

```
> RightInverseRat(R, Alg);
```

$$
\left[\begin{array}{c}
0 \\
-\frac{1}{t}
\end{array}\right]
$$

Let us compute the Brunovský canonical form:

```
> B:=BrunovskyRat(R,Alg);
```

$$
B:=\left[\begin{array}{cc}
-\frac{1}{t^{2}} & 0 \\
\frac{2}{t^{3}} & -\frac{1}{t}
\end{array}\right]
$$

Let us check that the variables $z$ and $v$ defined by $(z, v)^{T}=B(x, u)^{T}$ satisfy a Brunovský canonical form:

```
> E:=EliminationRat(linalg[stackmatrix](B, R), [x,u], [z,v,0],
> Alg):
> ApplyMatrix(E[1], [x(t),u(t)], Alg)=ApplyMatrix(E[2],
> [z(t),v(t)],Alg);
\[
\left[\begin{array}{c}
0 \\
\mathrm{u}(t) \\
\mathrm{x}(t)
\end{array}\right]=\left[\begin{array}{c}
-\left(\frac{d}{d t} \mathrm{z}(t)\right)+\mathrm{v}(t) \\
-2 \mathrm{z}(t)-t \mathrm{v}(t) \\
-t^{2} \mathrm{z}(t)
\end{array}\right]
\]
```

The first equation shows that $z$ and $v$ satisfy a Brunovský canonical form. The last two equations give $x$ and $u$ in terms of $z$ and $v$. We refer to 8 ] for more difficult examples of time-varying ordinary differential linear systems.

Example 9. Let us consider the example of a two reflector antenna [26]. We first define an Ore algebra with a differential operator $D t$ w.r.t. time $t$ and a constant time-delay operator $\delta$. Note also that the constants $K 1, K 2, T e, K p, K c$ have to be declared in the definition of the Ore algebra.

```
> Alg:=DefineOreAlgebra(diff=[Dt,t], dual_shift=[delta,s],
> polynom=[t,s], comm=[K1,K2,Te,Kp,Kc], shift_action=[delta,t]):
```

Enter the matrix $R$ of the differential time-delay linear system:

$$
\begin{aligned}
& >\mathrm{R}:=\mathrm{evalm}([[\mathrm{Dt},-\mathrm{K} 1,0,0,0,0,0,0,0] \text {, } \\
& >[0, \mathrm{Dt}+\mathrm{K} 2 / \mathrm{Te}, 0,0,0,0,-\mathrm{Kp} / \mathrm{Te} * d e l \mathrm{ta},-\mathrm{Kc} / \mathrm{Te} * \operatorname{delta},-\mathrm{Kc} / \mathrm{Te} * d e l \mathrm{ta}] \text {, } \\
& >\quad[0,0, D t,-K 1,0,0,0,0,0] \text {, } \\
& >\quad[0,0,0, \mathrm{Dt}+\mathrm{K} 2 / \mathrm{Te}, 0,0,-\mathrm{Kc} / \mathrm{Te} * \operatorname{delta},-\mathrm{Kp} / \mathrm{Te} * \operatorname{delta},-\mathrm{Kc} / \mathrm{Te} * \operatorname{delta}] \text {, } \\
& >\quad[0,0,0,0, \mathrm{Dt},-\mathrm{K} 1,0,0,0] \text {, } \\
& >\quad[0,0,0,0,0, \mathrm{Dt}+\mathrm{K} 2 / \mathrm{Te},-\mathrm{Kc} / \mathrm{Te} * \operatorname{del} \mathrm{ta},-\mathrm{Kc} / \mathrm{Te} * \operatorname{delta},-\mathrm{Kp} / \mathrm{Te} * \operatorname{delta}] \mathrm{)} \text {; } \\
& R:=\left[\begin{array}{ccccccccc}
D t & -K 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & D t+\frac{K 2}{T e} & 0 & 0 & 0 & 0 & -\frac{K p \delta}{T e} & -\frac{K c \delta}{T e} & -\frac{K c \delta}{T e} \\
0 & 0 & D t & -K 1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & D t+\frac{K 2}{T e} & 0 & 0 & -\frac{K c \delta}{T e} & -\frac{K p \delta}{T e} & -\frac{K c \delta}{T e} \\
0 & 0 & 0 & 0 & D t & -K 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & D t+\frac{K 2}{T e} & -\frac{K c \delta}{T e} & -\frac{K c \delta}{T e} & -\frac{K p \delta}{T e}
\end{array}\right]
\end{aligned}
$$

Then, we use an involution $\theta$ of $A l g$ in order to obtain $R \_a d j=\theta(R)$ :

$$
>\quad \text { R_adj:=Involution }(R, A l g):
$$

By means of the next command, we compute the torsion-free part (if Ext1[1] is not the identity matrix, then the torsion submodule is generated by the rows of Ext1 [2] modulo the module generated by the rows of $R$ ) and a parametrization of the torsion-free part in Ext1[3]. Equivalently, we check whether or not the two reflector antenna is controllable:

```
> st:=time(): Ext1:=Exti(R_adj, Alg, 1): time() - st;
    0.920
> Ext1[1];
```

$$
\left[\begin{array}{llllll}
1 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 1
\end{array}\right]
$$

We conclude that the first extension module $\operatorname{ext}^{1}{ }_{A l g}(\widetilde{N}, A l g)$ of the $A l g$-module $\widetilde{N}=A l g^{1 \times 6} /\left(A^{1 \times 9} \theta(R)\right)$ associated with $R_{\text {_ }} a d j=\theta(R)$ is the zero module. Hence, the module defined by $R$ is torsion-free. Equivalently, $R$ is parametrizable and Ext1[3] gives a parametrization of $R$ involving three free parameters:
$>$ Ext1[3];
$\left[\begin{array}{ccc}K 1 \delta K c & K 1 \delta K c & K p K 1 \delta \\ D t \delta K c & D t \delta K c & K p \delta D t \\ K 1 \delta K c & K p K 1 \delta & K 1 \delta K c \\ D t \delta K c & K p \delta D t & D t \delta K c \\ K p K 1 \delta & K 1 \delta K c & K 1 \delta K c \\ K p \delta D t & D t \delta K c & D t \delta K c \\ 0 & 0 & D t^{2} T e+D t K 2 \\ 0 & D t^{2} T e+D t K 2 & 0 \\ D t^{2} T e+D t K 2 & 0 & 0\end{array}\right]$

The same parametrization can be obtained by using Parametrization. The result involves three free functions $\xi_{1}, \xi_{2}, \xi_{3}$ :

$$
\begin{aligned}
& >\text { Parametrization(R, Alg); } \\
& \qquad\left[\begin{array}{c}
K 1 K c \xi_{1}(t-1)+K 1 K c \xi_{2}(t-1)+K p K 1 \xi_{3}(t-1) \\
K c \mathrm{D}\left(\xi_{1}\right)(t-1)+K c \mathrm{D}\left(\xi_{2}\right)(t-1)+K p \mathrm{D}\left(\xi_{3}\right)(t-1) \\
K 1 K c \xi_{1}(t-1)+K p K 1 \xi_{2}(t-1)+K 1 K c \xi_{3}(t-1) \\
K c \mathrm{D}\left(\xi_{1}\right)(t-1)+K p \mathrm{D}\left(\xi_{2}\right)(t-1)+K c \mathrm{D}\left(\xi_{3}\right)(t-1) \\
K p K 1 \xi_{1}(t-1)+K 1 K c \xi_{2}(t-1)+K 1 K c \xi_{3}(t-1) \\
K p \mathrm{D}\left(\xi_{1}\right)(t-1)+K c \mathrm{D}\left(\xi_{2}\right)(t-1)+K c \mathrm{D}\left(\xi_{3}\right)(t-1) \\
T e\left(\mathrm{D}^{(2)}\right)\left(\xi_{3}\right)(t)+K 2 \mathrm{D}\left(\xi_{3}\right)(t) \\
T e\left(\mathrm{D}^{(2)}\right)\left(\xi_{2}\right)(t)+K 2 \mathrm{D}\left(\xi_{2}\right)(t) \\
T e\left(\mathrm{D}^{(2)}\right)\left(\xi_{1}\right)(t)+K 2 \mathrm{D}\left(\xi_{1}\right)(t)
\end{array}\right]
\end{aligned}
$$

The two reflector antenna is not a flat system [15, 26] because $\operatorname{ext}^{2}{ }_{A l g}(\tilde{N}, A l g)$ of the Alg-module $\widetilde{N}$ is different from zero as it is shown next:

```
> st:=time(): Ext2:=Exti(R_adj, Alg, 2): time() - st;
    0.750
> Ext2[1];
\(\left[\begin{array}{ccc}\delta & 0 & 0 \\ D t^{2} T e+D t K 2 & 0 & 0 \\ 0 & \delta & 0 \\ 0 & D t^{2} T e+D t K 2 & 0 \\ 0 & 0 & \delta \\ 0 & 0 & D t^{2} T e+D t K 2\end{array}\right]\)
```

Since the torsion-free degree $i(M)$ of $M=A l g^{1 \times 9} /\left(A l g^{1 \times 6} R\right)$ is equal to 1 (i.e., $M$ is a torsion-free but not a projective Alg-module [11, 47]), we can find a polynomial $\pi$ in the variable $\delta$ such that the system is $\pi$-free [15, 26]:

```
> PiPolynomial(R, Alg, [delta]);
```

We obtain $\pi=\delta$. By definition of the $\pi$-polynomial [15, 26], this means that if we can permit the time-advance operator $\delta^{-1}$, then the system of the two reflector antenna becomes flat, i.e., the new $D=A l g\left[\delta^{-1}\right]$-module $P=D^{1 \times 9} /\left(D^{1 \times 6} R\right)$ associated with the system is free. We shall find a basis for the $D$-module $P$ below.

We note that the fact that the two reflector antenna is not a flat system (without the advance operator $\delta^{-1}$ ) is coherent with the fact that the full rowrank matrix $R$ does not admit a right-inverse. Indeed, we can prove that a full row-rank matrix $R$ admits a right-inverse if and only if the Alg-module $M=A l g^{1 \times 9} /\left(A l g^{1 \times 6} R\right)$ is projective [11. By the Quillen-Suslin theorem (see 3 of Theorem (1), projective modules over commutative polynomial rings are free. This remark applies to our situation as we have:

```
> SyzygyModule(R, Alg); RightInverse(R, Alg);
    INJ(6)
```

    []
    The fact that the system is not flat is also coherent with the fact that its parametrization Ext1 [3] does not admit a left-inverse. Indeed, a linear system is flat if and only if it admits a left-invertible parametrization [11].

```
> LeftInverse(Ext1[3], Alg);
```

We finish by computing a basis of the free $D=A l g\left[\delta^{-1}\right]$-module $P$. In the terminology of control, such a basis is called a flat output. We apply LocalLeftINVERSE to the parametrization Ext1[3] by allowing to invert $\delta$ :

$$
\begin{aligned}
& \text { > S:=LocalLeftInverse(Ext1[3], [delta], Alg); }
\end{aligned}
$$

$$
\begin{aligned}
& \% 1:=K p^{2}-2 K c^{2}+K p K c
\end{aligned}
$$

By construction, the matrix $S$ is a left-inverse of $\operatorname{Ext1}[3]$ :

```
> Mult(S, Ext1[3], Alg);
```

Therefore, $\left(z_{1}, z_{2}, z_{3}\right)^{T}=S\left(x_{1}, \ldots, x_{6}, u_{1}, u_{2}, u_{3}\right)^{T}$ is a basis of the $D$ module $P$ associated with $R$, and thus, a flat output of the two reflector antenna. Therefore, a flat output $\left(z_{1}, \quad z_{2}, \quad z_{3}\right)^{T}$ of the system is defined by:

$$
\begin{aligned}
& >\operatorname{evalm}([\operatorname{seq}([z[\mathrm{i}](\mathrm{t})], \mathrm{i}=1 . .3)])= \\
& >\operatorname{ApplyMatrix}(\mathrm{S},[\mathrm{seq}(\mathrm{x}[\mathrm{i}](\mathrm{t}), \mathrm{i}=1 . .6), \text { seq }(\mathrm{u}[\mathrm{i}](\mathrm{t}), \mathrm{i}=1 . .3)], \mathrm{Alg}) ; \\
& \\
& {\left[\begin{array}{l}
z_{1}(t) \\
z_{2}(t) \\
z_{3}(t)
\end{array}\right]=\left[\begin{array}{l}
-\frac{K c x_{1}(t+1)}{K 1 \% 1}-\frac{K c x_{3}(t+1)}{K 1 \% 1}+\frac{(K c+K p) x_{5}(t+1)}{K 1 \% 1} \\
-\frac{K c x_{1}(t+1)}{K 1 \% 1}+\frac{(K c+K p) x_{3}(t+1)}{K 1 \% 1}-\frac{K c x_{5}(t+1)}{K 1 \% 1} \\
\frac{(K c+K p) x_{1}(t+1)}{K 1 \% 1}-\frac{K c x_{3}(t+1)}{K 1 \% 1}-\frac{K c x_{5}(t+1)}{K 1 \% 1}
\end{array}\right]} \\
& \\
& \% 1:=K p K c-2 K c^{2}+K p^{2}
\end{aligned}
$$

Finally, if we substitute $\left(z_{1}, \quad z_{2}, \quad z_{3}\right)^{T}$ into the parametrization Ext1[3] of the system, we obtain $\left(x_{1}, \ldots, x_{6}, u_{1}, u_{2}, u_{3}\right)^{T}=T\left(x_{1}, \ldots, x_{6}, u_{1}, u_{2}, u_{3}\right)^{T}$, where the matrix $T$ is defined by:

$$
\left.\begin{array}{l}
>\mathrm{T}:=\operatorname{Mult}(\operatorname{Ext} 1[3], \mathrm{S}, \mathrm{Alg}) ; \\
\\
T:=\left[\begin{array}{c}
1,0,0,0,0,0,0,0,0 \\
\frac{D t}{K 1}, 0,0,0,0,0,0,0,0 \\
0,0,1,0,0,0,0,0,0 \\
0,0, \frac{D t}{K 1}, 0,0,0,0,0,0 \\
0,0,0,0,1,0,0,0,0 \\
0,0,0,0, \frac{D t}{K 1}, 0,0,0,0
\end{array}\right. \\
\frac{D t(D t T e+K 2)(K p+K c)}{\delta K 1 \% 1}, 0, \% 2,0, \% 2,0,0,0,0 \\
\% 2,0, \frac{D t(D t T e+K 2)(K p+K c)}{\delta K 1 \% 1}, 0, \% 2,0,0,0,0 \\
\% 2,0, \% 2,0, \frac{D t(D t T e+K 2)(K p+K c)}{\delta K 1 \% 1}, 0,0,0,0
\end{array}\right]
$$

We note that $\left(x_{2}, x_{4}, x_{6}, u_{1}, u_{2}, u_{3}\right)^{T}$ is expressed in terms of $x_{1}, x_{3}$ and $x_{5}$ only. Thus, $\left(x_{1}, x_{3}, x_{5}\right)$ is also a basis of the $D=\operatorname{Alg}\left[\delta^{-1}\right]$-module $P$ (compare with [26]). More precisely, we have:

```
> evalm([seq([x[i] (t)=ApplyMatrix(T,[seq(x[j] (t),j=1..6),
> seq(u[j](t),j=1..3)],Alg)[i,1]],i=1..6)]);
```

$$
\begin{aligned}
& {\left[\begin{array}{c}
x_{1}(t)=x_{1}(t) \\
x_{2}(t)=\frac{\mathrm{D}\left(x_{1}\right)(t)}{K 1} \\
x_{3}(t)=x_{3}(t) \\
x_{4}(t)=\frac{\mathrm{D}\left(x_{3}\right)(t)}{K 1} \\
x_{5}(t)=x_{5}(t) \\
x_{6}(t)=\frac{\mathrm{D}\left(x_{5}\right)(t)}{K 1}
\end{array}\right]} \\
& >\text { evalm([seq([u[i] (t)=ApplyMatrix (T, [seq(x[j](t),j=1..6), } \\
& >\operatorname{seq}(u[j](t), j=1 . .3)], A l g)[6+i, 1]], i=1 . .3)]) \text {; } \\
& {\left[u_{1}(t)=\frac{K 2(K c+K p) \mathrm{D}\left(x_{1}\right)(t+1)}{K 1 \% 1}+\frac{T e(K c+K p)\left(\mathrm{D}^{(2)}\right)\left(x_{1}\right)(t+1)}{K 1 \% 1}\right.} \\
& -\frac{K 2 K c \mathrm{D}\left(x_{3}\right)(t+1)}{K 1 \% 1}-\frac{T e K c\left(\mathrm{D}^{(2)}\right)\left(x_{3}\right)(t+1)}{K 1 \% 1}-\frac{K 2 K c \mathrm{D}\left(x_{5}\right)(t+1)}{K 1 \% 1} \\
& \left.-\frac{T e K c\left(\mathrm{D}^{(2)}\right)\left(x_{5}\right)(t+1)}{K 1 \% 1}\right] \\
& {\left[u_{2}(t)=-\frac{K 2 K c \mathrm{D}\left(x_{1}\right)(t+1)}{K 1 \% 1}-\frac{T e K c\left(\mathrm{D}^{(2)}\right)\left(x_{1}\right)(t+1)}{K 1 \% 1}+\frac{K 2(K c+K p) \mathrm{D}\left(x_{3}\right)(t+1)}{K 1 \% 1}\right.} \\
& \left.+\frac{T e(K c+K p)\left(\mathrm{D}^{(2)}\right)\left(x_{3}\right)(t+1)}{K 1 \% 1}-\frac{K 2 K c \mathrm{D}\left(x_{5}\right)(t+1)}{K 1 \% 1}-\frac{T e K c\left(\mathrm{D}^{(2)}\right)\left(x_{5}\right)(t+1)}{K 1 \% 1}\right] \\
& {\left[u_{3}(t)=-\frac{K 2 K c \mathrm{D}\left(x_{1}\right)(t+1)}{K 1 \% 1}-\frac{T e K c\left(\mathrm{D}^{(2)}\right)\left(x_{1}\right)(t+1)}{K 1 \% 1}-\frac{K 2 K c \mathrm{D}\left(x_{3}\right)(t+1)}{K 1 \% 1}\right.} \\
& -\frac{T e K c\left(\mathrm{D}^{(2)}\right)\left(x_{3}\right)(t+1)}{K 1 \% 1}+\frac{K 2(K c+K p) \mathrm{D}\left(x_{5}\right)(t+1)}{K 1 \% 1} \\
& \left.+\frac{T e(K c+K p)\left(\mathrm{D}^{(2)}\right)\left(x_{5}\right)(t+1)}{K 1 \% 1}\right] \\
& \% 1:=K p K c \quad 2 K c^{2}+K p^{2}
\end{aligned}
$$

We refer to [26] for applications of the previous results to the motion planning and tracking problems. See also 8 for more details.

Example 10. We consider the differential time-delay system of a vibrating string with an interior mass 27. We define the Ore algebra $A l g$, where D is the differential operator w.r.t. $t$ and $\sigma_{1}$ and $\sigma_{2}$ are two non-commensurate time-delay operators. The constant parameters $\eta_{1}, \eta_{2}$ (composition of the mass, tensions and densities) of the system must be declared in the algebra Alg:

```
> Alg:=DefineOreAlgebra(diff=[D,t], dual_shift=[sigma1,y1],
> dual_shift=[sigma2,y2], polynom=[t,y1,y2], comm=[eta1,eta2]):
```

We only study the case of position control on both boundaries [27]. For the case of a single control, we refer to [8]. We enter the system matrix $R$ :

```
> R:=evalm([[1,1,-1,-1,0,0],[D+eta1,D-eta1,-eta2,eta2,0,0],
> [sigma1^2,1,0,0,-sigma1,0],[0,0,1,sigma2^2,0,-sigma2]]);
\[
R:=\left[\begin{array}{cccccc}
1 & 1 & -1 & -1 & 0 & 0 \\
\mathrm{D}+\eta 1 & \mathrm{D}-\eta 1 & -\eta 2 & \eta 2 & 0 & 0 \\
\sigma 1^{2} & 1 & 0 & 0 & -\sigma 1 & 0 \\
0 & 0 & 1 & \sigma 2^{2} & 0 & -\sigma 2
\end{array}\right]
\]
```

We use an involution $\theta$ of $A l g$ in order to obtain $R_{-} a d j=\theta(R)$ :

```
> R_adj:=Involution(R, Alg):
```

We check controllability of the system by applying Exti to R_adj:

```
> st:=time(): Ext1:=Exti(R_adj, Alg, 1): time()-st; Ext1[1];
```

    1.191
    $\left[\begin{array}{llll}1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1\end{array}\right]$

Since Ext1[1] is the identity matrix, then we obtain that the Alg-module $M=A l g^{1 \times 6} /\left(A l g^{1 \times 4} R\right)$ associated with the system is torsion-free. This means that the vibrating string with interior mass is controllable and, equivalently, parametrizable. A parametrization of the system is then given by Ext1[3]:

```
> Ext1[3];
[c}c=[\begin{array}{c}{2\sigma2\eta2,-\sigma2\sigma1\eta2,-\eta2\sigma1+\sigma1\eta1-\sigma1\textrm{D}}\\{0,\sigma2\sigma1\eta2,\eta2\sigma1+\sigma1\textrm{D}+\sigma1\eta1}\\{\sigma2\textrm{D}+\sigma2\eta2+\sigma2\eta1,-\sigma2\sigma1\eta1,0}\end{array}
```

Therefore, the system can be parametrized by means of three free functions. We now want to check whether this parametrization is a minimal one [11, 33. In order to do that, let us compute the rank of the Alg-module $M$.

```
> OreRank(R, Alg);
```

$$
2
$$

Hence, we know that there exist some parametrizations of the system with only two arbitrary functions [11, 33]. We find some minimal parametrizations:

```
> P:=MinimalParametrizations(R, Alg);
```

$$
\begin{aligned}
& P:=\left[\begin{array}{cc}
2 \sigma 2 \eta 2 & -\sigma 2 \sigma 1 \eta 2 \\
0 & \sigma 2 \sigma 1 \eta 2 \\
\sigma 2 \mathrm{D}+\sigma 2 \eta 2+\sigma 2 \eta 1 & -\sigma 2 \sigma 1 \eta 1 \\
-\sigma 2 \mathrm{D}+\sigma 2 \eta 2-\sigma 2 \eta 1 & \sigma 2 \sigma 1 \eta 1 \\
2 \sigma 2 \sigma 1 \eta 2 & \sigma 2 \eta 2-\sigma 2 \eta 2 \sigma 1^{2} \\
\mathrm{D}-\mathrm{D} \sigma 2^{2}+\eta 2 \sigma 2^{2}-\eta 1 \sigma 2^{2}+\eta 2+\eta 1-\sigma 1 \eta 1+\sigma 1 \eta 1 \sigma 2^{2}
\end{array}\right], \\
& {\left[\begin{array}{c}
2 \sigma 2 \eta 2,-\eta 2 \sigma 1+\sigma 1 \eta 1-\sigma 1 \mathrm{D} \\
0, \eta 2 \sigma 1+\sigma 1 \mathrm{D}+\sigma 1 \eta 1 \\
\sigma 2 \mathrm{D}+\sigma 2 \eta 2+\sigma 2 \eta 1,0 \\
-\sigma 2 \mathrm{D}+\sigma 2 \eta 2-\sigma 2 \eta 1,2 \sigma 1 \eta 1 \\
2 \sigma 2 \sigma 1 \eta 2,-\eta 2 \sigma 1^{2}+\eta 2+\eta 1 \sigma 1^{2}-\sigma 1^{2} \mathrm{D}+\mathrm{D}+\eta 1 \\
\mathrm{D}-\mathrm{D} \sigma 2^{2}+\eta 2 \sigma 2^{2}-\eta 1 \sigma 2^{2}+\eta 2+\eta 1,2 \sigma 2 \sigma 1 \eta 1
\end{array}\right],} \\
& \left.\left[\begin{array}{cc}
-\sigma 2 \sigma 1 \eta 2 & -\eta 2 \sigma 1+\sigma 1 \eta 1-\sigma 1 \mathrm{D} \\
\sigma 2 \sigma 1 \eta 2 & \eta 2 \sigma 1+\sigma 1 \mathrm{D}+\sigma 1 \eta 1 \\
-\sigma 2 \sigma 1 \eta 1 & 0 \\
\sigma 2 \sigma 1 \eta 1 & 2 \sigma 1 \eta 1 \\
\sigma 2 \eta 2-\sigma 2 \eta 2 \sigma 1^{2} & -\eta 2 \sigma 1^{2}+\eta 2+\eta 1 \sigma 1^{2}-\sigma 1^{2} \mathrm{D}+\mathrm{D}+\eta 1 \\
-\sigma 1 \eta 1+\sigma 1 \eta 1 \sigma 2^{2} & 2 \sigma 2 \sigma 1 \eta 1
\end{array}\right]\right]
\end{aligned}
$$

Since $R$ has full row rank (this fact can be checked by computing $\operatorname{SyzygyModule}(R, A l g)$ ), we know that $M$ is projective, and thus, free if and only if $R$ admits a right-inverse (see [11, 33] for more details).
> RightInverse(R, Alg);

Hence, $M$ is not projective, which implies that $M$ is not free, i.e., the vibrating string with interior mass is not a flat system [27]. Another way to verify this fact is to compute $\operatorname{ext}_{A l g}^{2}(\widetilde{N}, A l g)$ and $\operatorname{ext}_{A l g}^{3}(\widetilde{N}, A l g)$ of the $A l g$-module $\widetilde{N}=$ $A l g^{1 \times 4} /\left(A l g^{1 \times 6}\right.$ R_adj):
> Exti(R_adj, Alg, 2);

$$
\left[\left[\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right],\left[\begin{array}{ccc}
\sigma 2 \eta 2 & 0 & \sigma 1 \eta 1 \\
\eta 2+\eta 1+\mathrm{D}-\sigma 1 \eta 1 & 0 \\
0 & \sigma 2 \eta 2 & \eta 2+\eta 1+\mathrm{D}
\end{array}\right],\left[\begin{array}{c}
-\sigma 1 \eta 1 \\
-\mathrm{D}-\eta 2-\eta 1 \\
\sigma 2 \eta 2
\end{array}\right]\right]
$$

> Exti(R_adj, Alg, 3);

$$
\left[\left[\begin{array}{c}
\sigma 2 \\
\sigma 1 \\
\eta 2+\eta 1+\mathrm{D}
\end{array}\right],[1], \operatorname{SURJ}(1)\right]
$$

We see that $\operatorname{ext}^{2}{ }_{A l g}(\tilde{N}, A l g)$ equals zero but $\operatorname{ext}^{3}{ }_{A l g}(\widetilde{N}, A l g)$ is different from zero. Therefore, $M$ is a reflexive but not a projective $A l g$-module. Indeed, we recall that $M$ is reflexive (resp., projective) iff $\operatorname{ext}_{A l g}^{i}(\tilde{N}, A l g)$ equals zero for $i=1,2$ (resp., $i=1,2,3$ ). Let us find a polynomial $\pi$ in the variable $\sigma_{1}$ such that the system is $\pi$-free [15, 26, 27.

```
> PiPolynomial(R, Alg, [sigma1]);
    [\sigma1]
```

Let us find a polynomial $\pi$ in the variable $\sigma_{2}$ such that the system is $\pi$-free.

```
> PiPolynomial(R, Alg, [sigma2]);
    [\sigma2]
```

Hence, if we invert $\sigma_{1}$ or $\sigma_{2}$, i.e., we allow ourselves to have a time-advance operator, then, by definition of the $\pi$-polynomial, the system becomes flat. A flat output for this system can be computed from a left-inverse of the minimal parametrization $P$, where we allow $\sigma_{1}$ or $\sigma_{2}$ to appear in the denominators.

We compute the annihilator of the $\operatorname{Alg}$-module $M_{1}=A l g^{1 \times 2} /\left(A l g^{1 \times 6} P[1]\right)$ of the minimal parametrization $P[1]$.

```
> Ann1:=AnnExti(linalg[transpose](P[1]), Alg, 1);
    Ann1 := [\sigma2]
```

Let us compute a left-inverse of the minimal parametrization $P[1]$ by allowing $\sigma_{2}$ to appear in the denominators.

```
> L1:=LocalLeftInverse(P[1], Ann1, Alg);
```

$$
L 1:=\left[\begin{array}{cccccc}
0 & 0 & \frac{1}{2 \sigma 2 \eta 2} & \frac{1}{2 \sigma 2 \eta 2} & 0 & 0 \\
0 & \frac{\sigma 1}{\sigma 2 \eta 2} & -\frac{\sigma 1}{\sigma 2 \eta 2} & -\frac{\sigma 1}{\sigma 2 \eta 2} & \frac{1}{\sigma 2 \eta 2} & 0
\end{array}\right]
$$

We easily check that $L 1$ is a left-inverse of $P[1]$.
$>$ Mult(L1, P[1], Alg);

$$
\left[\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right]
$$

If we use $\sigma_{2}^{-1}$, then we obtain that a flat output of the system is defined by

$$
\left(\xi_{1}, \xi_{2}\right)^{T}=L 1\left(\phi_{1}, \psi_{1}, \phi_{2}, \psi_{2}, u, v\right)^{T}
$$

where $\phi_{1}, \psi_{1}, \phi_{2}, \psi_{2}, u, v$ are the system variables [27]. Let us point out that any multiplication of $\left(\xi_{1}, \xi_{2}\right)^{T}$ by a unimodular matrix over the commutative ring $\mathbb{Q}\left(\eta_{1}, \eta_{2}\right)\left[\frac{d}{d t}, \sigma_{1}, \sigma_{2}, \sigma_{2}^{-1}\right]$ gives a new flat output of the system. For instance, we obtain the following flat output of the system [27]:

$$
\xi_{1}^{\prime}=2 \eta_{2} \sigma_{2} \xi_{1}=\phi_{2}+\psi_{2}, \quad \xi_{2}^{\prime}=\eta_{2} \sigma_{2}\left(\xi_{2}+2 \sigma_{1} \xi_{1}\right)=\sigma_{1} \psi_{1}+u
$$

We can repeat the same procedure for $P[2]$ and $P[3]$.

$$
\begin{gathered}
>\text { Ann2:=AnnExti (linalg[transpose] (P[2]), Alg, 1); } \\
>\text { Ann3:=AnnExti(linalg[transpose] }(\mathrm{P}[3]), \mathrm{Alg}, 1) ; \\
\text { Ann2 }:=[\eta 2+\eta 1+\mathrm{D}]
\end{gathered}
$$

$$
\text { Ann3 }:=[\sigma 1]
$$

The annihilator of $P[3]$ only contains $\sigma_{1}$. Let us compute a flat output by allowing the time-advance operator $\sigma_{1}^{-1}$ to appear in the basis.
> L3:=LocalLeftInverse(P[3], Ann3, Alg);

$$
L 3:=\left[\begin{array}{cccc}
0 & 0 & 0 & \frac{\sigma 2}{\sigma 1 \eta 1} \\
0 & 0-\frac{1}{\sigma 1 \eta 1} \\
0 & 0 & \frac{1}{2 \sigma 1 \eta 1} & \frac{1}{2 \sigma 1 \eta 1}
\end{array} 0\right.
$$

$L 3$ is a left-inverse of $P[3]$ over $\mathbb{Q}\left(\eta_{1}, \eta_{2}\right)\left[\frac{d}{d t}, \sigma_{1}, \sigma_{2}, \sigma_{1}^{-1}\right]$ as we can check:

```
> Mult(L3, P[3], Alg);
```

$$
\left[\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right]
$$

Therefore, if we use the time-advance operator $\sigma_{1}^{-1}$, we obtain the flat output of the system $\left(\xi_{1}, \xi_{2}\right)^{T}=L 3\left(\phi_{1}, \psi_{1}, \phi_{2}, \psi_{2}, u, v\right)^{T}$. Using trivial linear combinations of $\xi_{1}$ and $\xi_{2}$ over the ring $\mathbb{Q}\left(\eta_{1}, \eta_{2}\right)\left[\frac{d}{d t}, \sigma_{1}, \sigma_{2}, \sigma_{1}^{-1}\right]$, we then obtain that $\left(\xi_{1}^{\prime}=\sigma_{2} \psi_{2}-v, \quad \xi_{2}^{\prime}=\phi_{2}+\psi_{2}\right)$ is another flat output of the system over $\mathbb{Q}\left(\eta_{1}, \eta_{2}\right)\left[\frac{d}{d t}, \sigma_{1}, \sigma_{2}, \sigma_{1}^{-1}\right]$.

We refer to [27] for applications of the previous results to the motion planning and tracking problems [26. See also [8 for more details and examples.

## 8 Conclusion

We hope to have convinced the reader of the main interest of the package OreModules for the study of the structural properties of multidimensional linear systems over Ore algebras. To our knowledge, OreModules is the first implementation of homological methods with regard to applications in control theory. We hope that OreModules will become in the future a platform for the implementation of different algorithms obtained in the literature of multidimensional linear systems (see e.g., [1, 4, 13, 16, 18, 19, 20, 31, 32, 33, 34, 35, 37, 40, 43, 47, 49 and the references therein).

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