### Parametric polynomial systems and linkages

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Supelec, February 18, 2015

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# Linkages



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 $\theta_1 <$ 

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Parallel <u>P</u>R-<u>P</u>RR

#### Actuator variables

- *r*<sub>1</sub>, *r*<sub>2</sub>
- Pose variables

• *x*, *y* 

Passive variables

•  $\theta_1, \theta_2$ 



Equations

$$(F) \begin{cases} x = \cos(\frac{2\pi}{3})r_1 + \cos(\theta_1) \\ x = 1 + \cos(\frac{\pi}{3})r_2 + \cos(\theta_2) \\ y = \sin(\frac{2\pi}{3})r_1 + \sin(\theta_1) \\ y = 1 + \sin(\frac{\pi}{3})r_2 + \sin(\theta_2) \end{cases}$$

$$S : \begin{cases} f_1(\underline{T}, \underline{X}) = 0 \\ \vdots & \text{and} \\ f_k(\underline{T}, \underline{X}) = 0 \end{cases} \begin{cases} g_1(\underline{T}, \underline{X}) \neq 0 \\ \vdots \\ g_r(\underline{T}, \underline{X}) \neq 0 \end{cases}$$
$$f_i, g_j \in \mathbb{Q}[\underbrace{T_1, \cdots, T_s}_{parameters}, \underbrace{X_1, \cdots, X_n}_{unknowns}]$$

- Parametric system S
- Solutions:  $\mathcal{C} \subset \mathbb{C}^{s} \times \mathbb{C}^{n}$

### Parametric system

$$S_{\underline{t_0}}: \begin{cases} f_1(\underline{t_0}, \underline{X}) = 0 \\ \vdots & \text{and} \\ f_k(\underline{t_0}, \underline{X}) = 0 \end{cases} \begin{cases} g_1(\underline{t_0}, \underline{X}) \neq 0 \\ \vdots \\ g_r(\underline{t_0}, \underline{X}) \neq 0 \end{cases}$$
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### Parametric system



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In the applications we are interested in  $\mathcal{C}_{\mathbb{R}} \subset \mathbb{R}^{s} \times \mathbb{R}^{n}$ 

Robotics: Parallel robots

Vision: Camera calibration

Academic: Haas systems

[McAree, Daniel, Wenger, Chablat, ...]

[Gao, Tang, Yang, ...]

[Dickenstein, Rojas, Rusek, Shih]

General problem: classification of the parameters' space

- Number of solutions of  $S_{t_0}$  depends on  $t_0$
- $\Rightarrow$  Classification of the parameters

# State of the art (non exhaustive)

- Collins (1970): Cylindrical Algebraic Decomposition
  - Implementations (QEPCAD, Redlog, Mathematica, ...) , Efficient in practice for less than 3 variables
  - Worst case doubly exponential in the number of variables
- Weispfenning (1992): Comprehensive Gröbner bases
  - Implementations (Singular, Maple, Risa/Asir, ...)
  - Time complexity not well understood
- Grigoriev, Vorobjov (1999): Maps of vector of multiplicities
  - Time complexity analysis
  - Difficult to implement efficiently
- Lazard, Rouillier (2004): Minimal discriminant variety
  - Computed with Gröbner bases and CAD
  - Relatively efficient in practice and in theory under some assumptions
  - General case: combinatorial factors spoiled practical efficiency

### Discriminant variety and classification

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$$S : \begin{cases} f_1(\underline{T}, \underline{X}) = 0 \\ \vdots \\ f_k(\underline{T}, \underline{X}) = 0 \end{cases} \begin{cases} g_1(\underline{T}, \underline{X}) \neq 0 \\ \vdots \\ g_r(\underline{T}, \underline{X}) \neq 0 \end{cases}$$
$$f_i, g_j \in \mathbb{Q}[\underline{T_1, \cdots, T_s}, \underbrace{X_1, \cdots, X_n}] \\ parameters & unknowns \end{cases}$$
$$\bullet \pi : \ \mathcal{C} = V(S) \rightarrow \mathbb{C}^s \text{ canonical projection} \\ (\underline{t}, \underline{x}) \mapsto \underline{t} \end{cases}$$
Definition: covering map  
Given a connected open set  $U \subset \mathbb{C}^s$ , we say that  $(\pi, U)$  is a covering map if  
 $\bullet \pi^{-1}(U) = \mathcal{C}_1 \cup \cdots \cup \mathcal{C}_m$ 

• 
$$\pi_{|\mathcal{C}_i}: \mathcal{C}_i \to U$$
 is a diffeomorpism

• 
$$C_i \cup C_j = \emptyset$$

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$$\pi$$
 :  $\mathcal{C} = V(S) \rightarrow \mathbb{C}^s$  canonical projection  
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#### Definition: Discriminant variety

 $D(\mathcal{C}) \subset \mathbb{C}^{s}$  s.t. for all connected open set  $U \subset \mathbb{C}^{s} \setminus D(\mathcal{C})$ 

 $(\pi, U)$  is a covering map

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Property of the complex discriminant variety in the real

For all connected open set  $U \subset \mathbb{R}^s \setminus D(\mathcal{C})$ 

 $(\pi_{\mathbb{R}}, U)$  is a covering map

Number of real roots of  $S_p^{\mathbb{R}}$  constant for all  $p \in U$ 



#### Definition: Minimal discriminant variety

The intersection of all the discriminant varieties of S.

$$D_{min}(\mathcal{C}) = V(D_1(\underline{T}), \dots, D_m(\underline{T}))$$
$$D_i \in \mathbb{Q}[T_1, \dots, T_s]$$

$$D_{min}(\mathcal{C}) = \begin{cases} D_{ineq}(\mathcal{C}):\\ D_{\infty}(\mathcal{C}):\\ D_{mult}(\mathcal{C}):\\ D_{sd}(\mathcal{C}): \end{cases}$$



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projection of  $\overline{\mathcal{C}} \cap \bigcup_i V(g_i(\underline{T}, \underline{X}))$ divergence of the solutions projection of the multiple solutions components of dimension < s



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#### Describing the real roots with the discriminant variety





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#### Describing the real roots with the discriminant variety



# Example

- 3-RPR: a 9-bar linkage
  - Parallel robot
  - r1 fixed
  - Parameter space Q: r<sub>2</sub>, r<sub>3</sub>
  - Workspace W: $B_{1x}, B_{1y}, \alpha_x, \alpha_y$
  - Constraint equations:

$$f_1 = f_2 = f_3 = f_4 = 0$$

### • Discriminant variety and partition of Q





# Cuspidal points

System (S):  $I: \begin{cases} f_1 = 0 \\ f_2 = 0 \\ f_3 = 0 \\ f_4 = 0 \end{cases}$ 

$$\mathcal{J}(I): \underline{j_0}:= \det(\vec{df_1}, \vec{df_2}, \vec{df_3}, \vec{df_4}) = 0$$

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$$\mathcal{J}(I+\mathcal{J}(I)): \begin{cases} j_0 := \det(\vec{df_1}, \vec{df_2}, \vec{df_3}, \vec{df_4}) = 0\\ j_1 := \det(\vec{df_1}, \vec{df_2}, \vec{df_3}, \vec{df_9}) = 0\\ j_2 := \det(\vec{df_1}, \vec{df_2}, \vec{df_9}, \vec{df_4}) = 0\\ j_3 := \det(\vec{df_1}, \vec{df_2}, \vec{df_3}, \vec{df_4}) = 0\\ j_4 := \det(\vec{df_9}, \vec{df_2}, \vec{df_3}, \vec{df_4}) = 0 \end{cases}$$

# Cuspidal points

System (S):  $J(I) : j_{0} := \det(\vec{df_{1}}, \vec{df_{2}}, \vec{df_{3}}, \vec{df_{4}}) = 0$   $f_{2} = 0$   $f_{3} = 0$   $f_{4} = 0$   $J(I + J(I)) : \begin{cases} j_{0} := \det(\vec{df_{1}}, \vec{df_{2}}, \vec{df_{3}}, \vec{df_{4}}) = 0 \\ j_{1} := \det(\vec{df_{1}}, \vec{df_{2}}, \vec{df_{3}}, \vec{df_{4}}) = 0 \\ j_{2} := \det(\vec{df_{1}}, \vec{df_{2}}, \vec{df_{3}}, \vec{df_{4}}) = 0 \\ j_{3} := \det(\vec{df_{1}}, \vec{df_{2}}, \vec{df_{3}}, \vec{df_{4}}) = 0 \\ j_{4} := \det(\vec{df_{1}}, \vec{df_{2}}, \vec{df_{3}}, \vec{df_{4}}) = 0 \end{cases}$ 

- Curve in  $\mathbb{C}^7$  (determinantal ideal)
- Description:
  - r<sub>1</sub> : parameter
  - $r_2, r_3, t_x, t_y, u_x, u_y$  : unknowns
  - $N: x \mapsto #\{\text{real solutions of } (S) \text{ for } r_1 = x\}$



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# 10 cuspidal points



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### 11-bar linkage

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# Planar rigid linkage





- Several assembly modes
- Number depends on *c<sub>ij</sub>*
- Max number of assembly modes?

# Properties of minimally rigid linkages

• Construction steps



Henneberg steps:  $H_1$  and  $H_2$ 

• 3-bar rigid linkage



• Construction steps



Henneberg steps:  $H_1$  and  $H_2$ 

• 5-bar rigid linkage



• Construction steps



Henneberg steps:  $H_1$  and  $H_2$ 

• 7-bar rigid linkage



• Construction steps



Henneberg steps:  $H_1$  and  $H_2$ 

• 9-bar rigid linkage



• Construction steps



Henneberg steps:  $H_1$  and  $H_2$ 

• 11-bar rigid linkage



## Properties known before [Emiris and M. 11]

#### Maximal number of assembly modes

bars	3	5	7	9	11	13	15	17
upper	2	4	8	24	64	128	512	2048
lower	2	4	8	24	48	96	288	576

#### Theorem

A linkage is minimally rigid  $\Leftrightarrow$  It is constructed with  $H_1$  and  $H_2$ 

### Corollary

$$#Links = 2#Joints - 3$$

#### Outline

- Upper Bound
  - Algebraic Modeling
  - Mixed Volume
- 2 Lower Bound
  - Adaptive Sampling

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### Algebraic Modeling I



*c<sub>ij</sub>*: 10 parameters *x<sub>i</sub>*, *y<sub>i</sub>*: 14 variables

$$x_1 = 0, y_1 = 0 x_2 = 1, y_2 = 0$$

$$\begin{cases} x_3^2 + y_3^2 = c_{13} \\ (x_3 - 1)^2 + y_3^2 = c_{23} \\ (x_5 - 1)^2 + y_5^2 = c_{25} \\ (x_6 - x_3)^2 + (y_6 - y_3)^2 = c_{36} \\ x_4^2 + y_4^2 = c_{14} \end{cases} \begin{cases} x_7^2 + y_7^2 = c_{17} \\ (x_6 - x_4)^2 + (y_6 - y_4)^2 = c_{46} \\ (x_5 - x_6)^2 + (y_5 - y_6)^2 = c_{56} \\ (x_7 - x_5)^2 + (y_7 - y_5)^2 = c_{57} \\ (x_4 - x_7)^2 + (y_4 - y_7)^2 = c_{47} \end{cases}$$

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### Number of solutions

• Mixed Volume: *n*! *Volume*(*Support*)



#### • Our system: 2<sup>10</sup>

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# Algebraic Modeling II

#### Distance matrix

		$v_1$	<i>v</i> <sub>2</sub>	V3	V4	<i>V</i> 5	V <sub>6</sub>	V7	
	Γ0	1	1	1	1	1	1	1	1
$v_1$	1	0	<i>c</i> <sub>12</sub>	<i>c</i> <sub>13</sub>	<i>c</i> <sub>14</sub>	<i>x</i> <sub>15</sub>	<i>x</i> <sub>16</sub>	<i>c</i> <sub>17</sub>	
<i>v</i> <sub>2</sub>	1	<i>c</i> <sub>12</sub>	0	<i>c</i> <sub>23</sub>	<i>x</i> <sub>24</sub>	<i>C</i> 25	<i>x</i> <sub>26</sub>	<i>x</i> 27	
V3	1	<i>c</i> <sub>13</sub>	<i>c</i> <sub>23</sub>	0	<i>x</i> 34	<i>x</i> 35	<i>c</i> <sub>36</sub>	X37	
V4	1	<i>C</i> <sub>14</sub>	<i>x</i> <sub>24</sub>	<i>x</i> 34	0	X45	<b>C</b> 46	<b>C</b> 47	
<i>V</i> 5	1	<i>x</i> <sub>15</sub>	<i>C</i> 25	<i>x</i> 35	X45	0	<i>C</i> 56	<i>C</i> 57	
v <sub>6</sub>	1	<i>x</i> <sub>16</sub>	x <sub>26</sub>	<i>c</i> <sub>36</sub>	<i>c</i> 46	<i>c</i> <sub>56</sub>	0	x <sub>67</sub>	
V7	[ 1	<i>c</i> <sub>17</sub>	x <sub>27</sub>	x <sub>37</sub>	C47	<i>C</i> 57	x <sub>67</sub>	0	

#### Theorem

The distance matrix has rank 4.

#### Corollary

All the 5x5 minors vanish.

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### Algebraic Modeling II





- Upper Bound
  - Mixed volume: 56
- Lower Bound?

# Adaptive Sampling

- Uniform sampling
  - No linkage found with 56 assembly modes
- Adaptive sampling
  - Simulated annealing
  - Cross-Entropy Method



Results

• Random simulations for different sampling methods

Uniform	Simulated annealing	Cross-entropy
44 (572)	52 (17)	52 (199)
42 (196)	54 (247)	54 (132)
48 (27)	48 (362)	52 (186)
44 (200)	52 (14)	54 (130)
42 (200)	54 (547)	<mark>56</mark> (497)
44 (424)	54 (315)	<mark>56</mark> (328)
46 (48)	<mark>56</mark> (425)	<mark>56</mark> (454)
42 (170)	50 (585)	54 (375)
42 (18)	54 (26)	<mark>56</mark> (552)
46 (366)	52 (474)	<mark>56</mark> (355)
42 (18) 46 (366)	54 (26) 52 (474)	56 (552) 56 (355)

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Results



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# Conclusion

#### • 9-bar linkage

- Discriminant variety can be computed:
  - on the equation constraints
  - on the cuspidal equation constraints
- Classification of the parameter space
- 11-bar linkage
  - No complete classification of the parameter space
  - Distance matrices and mixed volume:
    - at most 56 assembly modes
  - simulated annealing and cross entropy method:
    - a 11-bar linkage with exactly 56 assembly modes
- *n* vertices linkage
  - State-of-the-art:  $\Omega(2.89^n)$  and  $O(4^n)$  possible embeddings
  - New lower bound:  $\Omega(2.3^n)$

# Merci !

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